## decsustenno

## MACRO ASSEMBLER REFERENCE MANUAL

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AA-C780C-TB
```

April 1978

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## PREFACE

```
This manual is a reference for the programmer with some knowledge of
assemblers and assembly languages.
Using the MACRO assembler effectively involves using other
DECsystem-l0 facilities: the monitor (TOPS-l0), the LINK program, the
CREF program, a debugging program, a text editor (SOS or TECO), and
machine language. Therefore the following DECsystem-l0 documents will
prove useful:
Operating System Commands
    AA-0916C-TB
Monitor Calls
    AA-0974C-TB
LINK Reference Manual
    AA-0988C-TB
SOS User's Guide
    DEC-10-USOSA-A-D
Introduction to TECO
    DEC-10-UTECA-A-D
DDT Dynamic Debugging Techigue
    DEC-10-UDDTA-A-D
GALAXY Batch Reference Manual
    DEC-10-OGBRA-A-DN1
Hardware Reference Manual
    EK-10/20-HR-001
```


## CHAPTER 1

## INTRODUCTION TO MACRO

MACRO is the symbolic assembler program for the DECsystem-lo. The assembler reads a file of MACRO statements and composes relocatable binary machine instruction code suitable for loading by LINK, the system's linking loader.

MACRO is a statement-oriented language; statements are in free format and are processed in two passes. In processing statements, the assembler:

1. Interprets machine instruction mnemonics
2. Accepts symbol definitions
3. Interprets symbols
4. Interprets pseudo-ops
5. Accepts macro definitions
6. Expands macros on call
7. Assigns memory addresses
8. Generates a relocatable binary program file (.REL file) for input to LINK
9. Generates a program listing file showing source statements, the corresponding binary code, and any errors found
10. Generates a UNIVERSAL file that can be searched by other assemblies

In addition to translating machine instruction mnemonics and special-purpose operators called pseudo-ops, MACRO allows you to create your own language elements, called macros. In this way you can tailor the assembler's functions for each program.

Since the assembler is device independent, you can use any peripheral devices for input and output files. For example, you can use a terminal for your source program input, a line printer for your program listing output, and a disk for your binary program output.

MACRO programs must use the monitor for device-independent input/output services. (See the Monitor Calls manual.)

## INTRODUCTION TO MACRO

## NOTES

The following conventions are used
throughout this manual:

1. All numbers in the examples are
octal unless otherwise indicated.
2. All numbers in the text are decimal
unless otherwise indicated.
3. The name of the assembler, mACRO,

| appears in uppercase letters; |
| :--- |
| references to user-defined macros |

appear in lowercase letters.
4. Examples sometimes show the code
generated as it appears in the
program listing file. This file is
described in Section 6.1 .

### 1.1 HOW THE ASSEMBLER OPERATES

MACRO is a 2-pass assembler; it reads your source program twice. On Pass l, some symbolic addresses will not be resolved, if they refer to parts of the program not yet read. These symbolic references are entered in the symbol table and will be resolved on Pass 2.

The main purpose of Pass $l$ is to build symbol tables and to make a rudimentary assembly of each source statement.

The first task of Pass 1 is initializing all impure data areas that MACRO uses (internally) for assembly. This area includes all dynamic storage areas and all buffer areas.

MACRO then reads a command string into memory. This command string contains specifications for the files to be used during assembly. After scanning the command string for proper syntax, MACRO initializes the specified output files.

As assembly begins, MACRO initiates a routine that retrieves source lines from the proper input file. If no such file is currently open, MACRO opens the next input file specified in the command string. Source lines are assembled as they are retrieved from input files.

Assembly Pass 2 performs the same steps as Pass l. However, during Pass 2 MACRO writes the object code to the binary (and usually relocatable) output file; it also generates the program listing file, followed by the symbol table listing for the program.

MACRO can also generate a cross-referenced symbol table. (See Chapter 6.)

During Pass 2 MACRO also flags erroneous source statements with single-character error codes. (See Chapter 7.) These error codes appear in the program listing file.

The relocatable binary object file created during Pass 2 contains all binary code generated; this code is in a form suitable for loading by the LINK program. (See the LINK Reference Manual.)

MACRO processes relocation counters on both passes. If a labeled statement has a different relocation value on the second pass, MACRO generates a phase error.

### 1.2 ADDRESSES AND MEMORY

The address space of a DECsystem-l0 program consists of 256 K (lK = 1024 words), each word having 36 bits. Since the total number of storage locations is 2 to the 18 th power, the address of a location can be expressed in 18 bits, or one halfword.

The left halfword of a storage location is bits 0 to lif the right halfword is bits 18 to 35 .

### 1.3 RELOCATABLE ADDRESSES

Normally the binary program generated by MACRO is relocatable. This means that when the program is loaded for execution, it can be loaded anywhere in physical memory. (The address for loading is selected at load time, and depends on what has already been loaded.)

Unless you specify otherwise, MACRO assembles your binary program beginning with address 0 ( 400000 for high-segment code). References to addresses within your program are therefore relative to 0 (400000 for the high segment), and must be changed at loading time. LINK does this by adding the load address to all such relative addresses, resolving them to absolute addresses.

For programs assembled with multiple PSECT counters, each PSECT begins with the relative address 0. At load time, each PSECT has its own relocation constant; PSECT origins must be selected carefully to avoid overlapping of PSECTs in memory.

ELEMENTS OF MACRO

The character set recognized in MACRO statements includes all ASCII alphanumeric characters and 28 special characters (ASCII 040 through 137). Lowercase letters (ASCII l4l through l72) are treated internally as uppercase letters (ASCII lol through l32).

MACRO also recognizes seven ASCII control codes: horizontal tab (011), linefeed (012), vertical tab (0l3), formfeed (0l4), carriagereturn (015), CTRL/underscore (037), and CTRL/Z (032).

MACRO accepts any ASCII character in quoted text, or as text arguments to the ASCII and ASCIZ pseudo-ops.

NOTES

```
1. The line-continuation character (CTRL/-) is always effective.
2. Delimiters for certain pseudo-ops (such as ASCII, ASCIZ, and COMMENT) can be any nonblank, nontab ASCII character.
```

```
Characters and their codes are listed in Appendix A.
A MACRO program consists of statements made up of MACRO language
elements. Separated into general types, these are:
    1. Special characters
    2. Numbers
    3. Literals
    4. Symbols
    5. Expressions
    6. MACRO-defined mnemonics
    7. Pseudo-ops
    8. Macros
The format of a MACRO statement is discussed in Chapter 4.
```


### 2.1 SPECIAL CHARACTERS

Characters and combinations that have special interpretations in MACRO are listed in Appendix $B$. These interpretations apply only in the contexts described. In particular, they do not apply within comment fields or text strings.

## 2. 2 NUMBERS

The two properties of numbers that are important to MACRO are:

1. In what radix (base) the number is given
2. How the number should be placed in memory

You can control the interpretation of these properties by using pseudo-ops or special characters to indicate your choices.

### 2.2.1 Integers

MACRO stores an integer in its binary form, right justified in bits l to 35 of its storage word. If you use a sign, place it immediately before the integer. (If you omit the sign, the integer is assumed positive.) For a negative integer, MACRO first forms its absolute value in bits lo 35 , then takes its two's complement. Therefore a positive integer is stored with 0 in bit 0 , while a negative integer has 1 in bit 0 .

The largest integer that MACRO can store is 377777777777 (octal); the smallest (most negative) is 400000000000 (octal).

### 2.2.2 Radix

The initial implicit radix for a MACRO program is octal (base 8). The integers you use in your program will be interpreted as octal unless you indicate otherwise.

You can change the radix to any base from 2 to 10 by using the RADIX pseudo-op. (See the pseudo-op RADIX in Chapter 3.) The new radix will remain in effect until you change it.

Without changing the prevailing radix, you can write a particular expression in binary, octal, or decimal. To do this, prefix the integer with ${ }^{\wedge} B$ for binary, ${ }^{\circ}$ O for octal, or ${ }^{\wedge} D$ for decimal. The indicated radix applies only to the single integer immediately following it.

## ELEMENTS OF MACRO

NOTES

1. A single-digit number is always interpreted as radix l0. Thus 9 is seen as decimal 9 , even if the current radix is 8.
2. In the notations for " $B,{ }^{\text { }} \mathrm{D}$, and ${ }^{\text {" } O \text {, }}$ the up-arrow in the text indicates the up-arrow character, not the CONTROL character.
```
For example, suppose the current radix is 8. Then you can write the
decimal number 23 as:
2 7 ~ o c t a l ~ ( c u r r e n t ~ r a d i x )
^D23 decimal
*Bl0l11 binary
```

But you cannot write decimal 23 as "D45-22 since the "D applies only to the first number, 45 ; the 22 is octal. However, you can write decimal 23 as " $D<45-22\rangle$.

### 2.2.3 Adding Zeros to Integers in Source Code

You can add zeros to an integer (multiply it by a constant) in your program by suffixing $K, M$, or $G$ to it.

```
K adds 3 zeros (K = "kilo-", thousands)
M adds 6 zeros (M = "mega-", millions)
G adds 9 zeros (G = "giga-", billions)
```

These zeros are suffixed before any conversion, so that in radix 10 , 5 K means 5000 decimal; in radix 8 , 5 K means 5000 octal, or 2560 decimal.

### 2.2.4 Fixed-Point Decimal Numbers

To indicate a fixed-point decimal number, prefix it with ${ }^{\wedge} F$, include a decimal point wherever you wish, and suffix $B n$ to show that you want to place the "assumed point" after bit $n$ in the storage word. If you omit the decimal point, MACRO assumes that it follows the last digit. If you omit the Bn, MACRO assumes B35.

To handle the number, MACRO forms the integer part in a fullword register, and the fractional part in another fullword register. It then places the integer part (right justified) in bits lo $n$ ( $n$ is from your $B n$ ) of a binary word, and the fractional part (left justified) in the remaining bits. The integer part is truncated at the left, and the fractional part at the right. Bit 0 shows the sign of the number.

For example, ${ }^{\text {a }}$ Fl23.45B8 is formed in two registers as
000000000173 (integer part, right justified)
346314631462 (fractional part, left justified)
Since the $B n$ operator sets the assumed point after bit 8, the integer part is placed in bits 1 to 8 , and the fractional part in bits 9 to 35. (The sign bit 0 is 0 , showing a positive number.) Truncation is on the left and right, respectively, giving


You can show a fixed-point decimal number as negative by placing a minus sign before the "F. The absolute value of the negative number is formed in two registers as a positive number, then two's complemented. This complementing sets bit 0 to $l$, showing that the number is negative.

NOTE

> The binary number resulting from ${ }^{\wedge} F$ does not show where the assumed point should be. You must keep track of this through your own programming conventions.

Examples:

| 000000 | 000173 | $-F 123.45$ |
| :--- | :--- | :--- |
| 000173 | 346314 | $-F 123.45 B 17$ |
| 346314 | 631462 | $M F 123.45 B-1$ |
| 777777 | 777604 | $-M F 123.45$ |
| 777604 | 431463 | $--M F 123.45 B 17$ |
| 431463 | 146316 | $-M F 123.45 B-1$ |

### 2.2.5 Floating-Point Decimal Numbers

In your program, a floating-point decimal number is a string of digits with a leading, trailing, or embedded decimal point and an optional leading sign. MACRO recognizes this as a mixed number in radix 10.

MACRO forms a floating-point decimal number with the sign in bit 0 , a binary exponent in bits 1 to 8 , and a normalized binary fraction in bits 9 to 35.

The normalized fraction can be viewed as follows: its numerator is the binary number in bits 9 to 35 , whose value is less than 2 to the 28th power, but greater than or equal to 2 to the 27 th power. Its denominator is 2 to the 28 th power, so that the value of the fraction is always less than l, but greater than or equal to 0 . (This value is 0 only if the entire stored number is 0.)

The binary exponent is incremented by. 128 so that exponents from -128 to 127 are represented as 0 to 255.

For a negative floating-point decimal number, MACRO first forms its absolute value as a positive number, then takes the two's complement of the entire word.

Examples:
The floating-point number 17. generates the binary
$0 \quad 10000101 \quad 100 \quad 010000 \quad 000 \quad 000 \quad 000 \quad 000 \quad 000 \quad 000$
where bit 0 shows the positive sign, bits 1 to 8 show the binary exponent, and bits 9 to 35 show the proper binary fraction. The binary exponent is 133 (decimal), which after subtracting the added 128 gives 5. The fraction is equal to 0.53125 decimal. And 0.53125 times 2 to the 5 th power is 17 , which is the number given.

Similarly, 153. generates
$0 \quad 10001000 \quad 100110010 \quad 000 \quad 000 \quad 000 \quad 000 \quad 000 \quad 000$
while -153. generates
101110 lll 011001110000000000000000000
These two examples show that a negative number is two's complemented. Notice that since the binary fraction for a negative number always has some nonzero bits, the exponent field (taken by itself) appears to be one's complemented.

As in FORTRAN, you can write a floating-point decimal number with a suffixed $E \pm n$, and the number will be multiplied by 10 to the $\pm n t h$ power. If the sign is missing, $n$ is assumed positive.

Examples:

| 2840000. | can be written | $284 . \mathrm{E}+4$ |
| ---: | :--- | :--- |
| 2840000. | can be written | .284 E 7 |
| .0000284 | can be written | $.284 \mathrm{E}-4$ |
| .0000284 | can be written | $284 . \mathrm{E}-7$ |

Using this $E$ notation with an integer (no decimal point) is not allowed, and causes an error. Therefore you can use 284.E4, but 284E4 is illegal.

NOTE
MACRO's algorithm for handling numbers
given with the E notation is not
identical to FORTRAN's. The binary
values generated by the two translators
may differ in the lowest order bits.

## 2．2．6 Binary Shifting

Binary shifting of a number with Bn sets the location of the rightmost bit at bit $n$ in the storage word，where $n$ is a decimal integer．The shift takes place after the binary number is formed．Any bits shifted outside the range（bits 0 to 35 ）of the storage word are lost．

For example，here are some numbers with their binary representations given in octal：

| 300000 | 000000 | D3E2 |
| :---: | :---: | :---: |
| 000000 | 042000 | －17825 |
| 000001. | 0000000 | 1． B 17 |
| 400000 | 000000 | 180 |
| 777777 | 777777 | －1836 |
| 000000 | 000001 | $1 \mathrm{B35}$ |
| 000000 | フワ7ワ77 | $\cdots 183 \%$ |

## 2．2．7 Underscore Shifting

You can also shift a number by using the underscore operator．（On some terminals this is a left－arrow．）If $V$ is an expression with value $n$ ，suffixing $\quad-V$ to a number shifts it $n$ bits to the left．（If $n$ is negative，the shift is to the right．）

In an expression of the form $W \_V, W$ and $V$ can be any expressions including symbols．The binary value of $W$ is formed in a register，$V$ is evaluated，and the binary of $W$ is shifted $V$ bits when placed in storage．

NOTE
An expression such as－3．75E4＿～D18 is legal，but the shift occurs after conversion to floating－point decimal storage format．Therefore the sign， exponent，and fraction fields are all shifted away from their usual locations． This is true also for other storage formats．

## 2．2．8 Querying the Position of a Bit Pattern

You can query the position of a bit pattern by prefixing ${ }^{n}$ L（up－arrow L）to an expression．This generates the number of leading zeros in the binary value of the expression．（＂L0 generates 36 decimal．）

## ELEMENTS OF MACRO

```
For example, suppose the current radix is l0. Then
    ^Ll53 generates 35 (29 decimal)
    ^Ll53. generates 1
    ^L-153 generates 0
    `L-l53. generates 0
```

In the first example, ${ }^{\text {N }}$ Ll53 generates 29 (decimal) because the binary
representation of 153 decimal has its leftmost 1 in bit 28:
000000000000000000000000000010011001
But in the second example, the binary form of 153. is in
floating-point format (see Section 2.2.5),
010001000100110010000000000000000000
and its leftmost $l$ is in bit l.
In both of the last two examples, ${ }^{\wedge}$ L-153 and " L-153. generate 0. This
is because a negative number in any format sets bit 0 to 1 .

### 2.3 LITERALS

A literal is a character string within square brackets inserted in your source code. MACRO stores the code generated by the enclosed string in a literal pool beginning with the first available literal storage location, and places the address of this location in place of the literal. The literal pool is normally at the end of the binary program. (See the pseudo-op LIT in Chapter 3.)

The statements


```
    2206000000137
are equivalent to
    13501 0 00002020' LIMS TL,FLLACE:
    22O6 O OO OOO137 FINACE: FOTNT 6, JBUFERy%
A literal can also be used to generate a constant:
    FUSH 17,EOM ;Gemerate zero fullworg
    MOUE L,[3,y,] {Generate a word with 3 in
    ; lefthalf and l4 in rishtmalf
```

    2-7
    
## ELEMENTS OF MACRO

Multiline literals are also allowed:

```
GETCHF: ILIIR T2,T1 ;Get a character
    CAIN T2,0 fis it a mull?
    JFST [MOUE TI,TXTFTR gYes, retrieve fointer
        ILIB ra,Tt yGet a mew character
        CAIN r2,"?" &IS it a question mark?
        JFST FMOUE T1,TXTFT1 ;Yes, set altermate mointer
                                    ILME reyrt ;Get the message character
                                    JFST GETHLFJ कGO to nelw rowtime
        FOF.d F,] {Not asessiom mark, returm
    FOF.S F, gNot a mull, returm
```

The text of a literal continues until a matching closing square bracket is found (unquoted and not in comment field).

A literal can include any term, symbol, expression, or statement, but it must generate at least one but no more than 99 words of data. A statement that does not generate data (such as a direct-assignment statement or a RADIX pseudo-op) can be included in a literal, but the literal must not consist entirely of such statements.

You can nest literals up to 18 levels. You can include any number of labels in a literal, but a forward reference to a label in a literal is illegal.

If you use a dot (.) in a literal to retrieve the location counter, remember that the counter is pointing at the statement containing the literal, not at the literal itself.

In nested literals, a dot location counter references a statement outside the outermost literal.

In the sequence

JFST EHREZ ACI, U
CATE ACT,OF
JRST + + 1
JFST FUTSTS.
SKTFE C
the expression . +1 generates the address of SKIPE $C$, not JRST EVTSTS.
Literals having the same value are collapsed in MACRO's literal pool. Thus for the statements:

FOSH FP:COT
FUSH FP, [O]
MOUET ACI,[ASCIZ/TEST1/]
the same address is shared by the two literals [0], and by the null word generated at the end of [ASCIZ /TESTl/]. Literal collapsing is suppressed for those literals that contain errors, undefined expressions, or EXTERNAL symbols.

## ELEMENTS OF MACRO

### 2.4 SYMBOLS

MACRO symbols include:

1. MACRO-defined pseudo-ops (discussed in Chapter 3)
2. MACRO-defined mnemonics (discussed in Section 2.6)
3. User-defined macros (discussed in Chapter 5)
4. User-defined opdefs (discussed at OPDEF in Chapter 3)
5. User-defined labels (discussed in this section)
6. Direct-assignment symbols (discussed in Section 2.4.2.2)
7. Dummy-arguments for macros (discussed in Chapter 5)

MACRO stores symbols in three symbol tables:

1. Op-code table: machine instruction mnemonics and pseudo-ops
2. Macro table: macros, user-defined OPDEFs, and synonyms (See the SYN pseudo-op in Chapter 3.)
3. User symbol table: labels and direct-assignment symbols

An entry in one of these tables shows the symbol, its type, and its value.

Symbols are helpful in your programs because:

1. Defining a symbol as a label gives a name to an address. You can use the label in debugging or as a target for program control statements.
2. In revising a program, you can change a value throughout your program by changing a symbol definition.
3. You can give names to values to make computations clearer.
4. You can make values available to other programs.

### 2.4.1 Selecting Valid Symbols

Follow these rules in selecting symbols:

1. Use only letters, numerals, dots (.), dollar signs (\$), and percent signs (\%). MACRO will consider any other character (including a blank) as a delimiter.
2. Do not begin a symbol with a numeral.
3. If you use a dot for the first character, do not use a numeral for the second. Do not use dots for the first two characters; doing so can interfere with MACRO's created symbols. (See Section 5.5.2.)
4. Make the first six characters unique among your symbols. You can use more than six characters, but MACRO will use only the first six.

## ELEMENTS OF MACRO

```
Examples:
    VELOCITY (legal, only VELOCI is meaningful to MACRO)
    CHG.VEL (legal, only CHG.VE is meaningful to MACRO)
    CHG VEL (illegal, looks like two symbols to MACRO)
    1STNUM (illegal, begins with a numeral)
    NUMl (legal)
    .llll (illegal, begins with dot-numeral)
    ..llll (unwise, could interfere with created symbols)
```


### 2.4.2 Defining Symbols

You can define a symbol by making it a label or by giving its value in a direct-assignment statement. Labels cannot be redefined, but direct-assignment symbols can be redefined anywhere in your program.

You can also define special-purpose symbols called OPDEFs and macros using the pseudo-op OPDEF and the pseudo-op DEFINE. (See Chapter 3.)

```
2.4.2.l Defining Labels - A label is always a symbol with a suffixed
colon. A label is in the first (leftmost) field of a MACRO statement
and is one of the forms:
    ERRFOUND: (MACRO uses only ERRFOU)
    CASEl: (legal label)
    OK:CONTIN: (legal; you can use more than one label
    at a location)
    CASE2:: (double colon declares label INTERNAL;
        see Section 2.4.5.2)
    (colon and exclamation point suppresses
    output by debugger)
    (double colon and exclamation point
declares label INTERNAL and suppresses
output by debugger)
```

When MACRO processes the label, the symbol and the current value of the location counter are entered in the user symbol table. A reference to the symbol addresses the code at the label.

You cannot redefine a label to have a value different from its original value. A label is relocatable if the address it represents is relocatable; otherwise it is absolute.
2.4.2.2 Direct Assignments - You define a direct-assignment symbol by associating it with an expression. (See Section 2.5 for a discussion of expressions.) A direct assignment is in one of the forms:

| symbol=expression | (symbol and value of expression are <br> entered in user symbol table) |
| :--- | :--- |
| symbol==expression | (symbol and value of expression are <br> entered in user symbol table, output by <br> debugger is suppressed) |
| symbol=:expression | (symbol and value of expression are <br> entered in user symbol table, symbol is <br> declared INTERNAL; see Section 2.4.5.2) |
| symbol==:expression | (symbol and value of expression are <br> entered in user symbol table, symbol is <br> declared INTERNAL, output by debugger is |
|  |  |

You can redefine a direct-assignment symbol at any time; the new direct assignment simply replaces the old definition.

## NOTE

If you assign a multiword value using direct assignment, only the first word of the value is assigned to the symbol. For example, $A=A S C I Z / A B C D E F G H / ~ i s$ equivalent to $A=A S C I Z / A B C D E /$, since only the first five characters in the string correspond to code in the first word.

### 2.4.3 Variable Symbols

You can specify a symbol as a variable by suffixing it with a number sign (\#). A variable symbol needs no explicit storage allocation. On finding your END statement, MACRO assembles variables into locations following the literal pool.

You can assemble variables anywhere in your program by using the VAR pseudo-op. This pseudo-op causes all variables found so far to be assembled immediately. (Variables found after the VAR statement are assembled at the end of the program or at the next VAR statement.)

## ELEMENTS OF MACRO

### 2.4.4 Using Symbols

When you use a symbol in your program, MACRO looks it up in the symbol tables. Normally MACRO searches the macro table first, then the op-code table, and finally the user symbol table. However, if MACRO has already found an operator in the current statement and is expecting operands, then it searches the user symbol table first.

You can control the order of search for symbol tables by using the pseudo-op .DIRECTIVE MACPRF.

### 2.4.5 Symbol Attributes

The value of a symbol is either relocatable or absolute. The relocatability of a label is determined by the relocatability of the address assigned to it. You can define either an absolute or a relocatable value for a direct-assignment symbol.

In addition, each symbol in your program has one of the following attributes: local, INTERNAL global, or EXTERNAL global. This attribute is determined when the symbol is defined.
2.4.5.1 Local Symbols - A local symbol is defined for the use of the current program only. You can define the same symbol to have different values in separately assembled programs. A symbol is local unless you indicate otherwise.
2.4.5.2 Global Symbols - A global symbol is defined in one program, but is also available for use in other programs. Its table entry is visible to all programs in which the symbol is declared global.

A global symbol must be declared INTERNAL in the program where it is defined; it can be defined in only one program. In other programs sharing the global symbol, it must be declared EXTERNAL; it can be EXTERNAL in any number of programs.

To declare a symbol as INTERNAL global, you can:

1. Use the INTERN pseudo-op.

INTEFN FILAG.
2. Insert a colon after $=$ in a direct-assignment statement.

FLAG2=: 200
FLAG3 $==: 200$
3. Use an extra colon with a label.

Flag $4:$ :
4. For subroutine entry points, use the ENTRY pseudo-op. (This pseudo-op does more than declare the symbol INTERNAL. See Chapter 3.)

ENTFY FIAGG

## ELEMENTS OF MACRO

To declare a symbol as an EXTERNAL global, you can:

1. Use the EXTERN pseudo-op.

EXTEFN FL.AGG
2. Suffix \#\# to the symbol at any of its uses. (Doing this once is sufficient, but you can use \#\# with all references to the symbol.)

FLAG7**

### 2.5 EXPRESSIONS

You can combine numbers and defined symbols with arithmetic and logical operators to form expressions. You can nest expressions by using angle brackets. MACRO evaluates each expression (innermost nesting levels first), and either resolves it to a fullword value, or generates a Polish expression to pass to LINK. (See Sections 2.5.3 and 2.5.4.)

### 2.5.1 Arithmetic Expressions

An arithmetic expression can include any number or defined symbol, and any of the following operators:

+ addition
- subtraction
* multiplication
/ division
These examples assume that WORDS, $X, Y$, and $Z$ have been defined elsewhere:

MOUEI 3,WORTIS/S
AMII 12, $X+Y-Z$
ALIMI 12, ©WORDG/5+1.*5

### 2.5.2 Logical Expressions

A logical expression can include any number or defined symbol whose value is absolute, and any of the following operators:
\& AND
! OR (inclusive OR)
^! XOR (exclusive OR)

-     - NOT

The unary operation "-A generates the fullword one's complement of the value of $A$.

Each of the binary operations \&, !, and "! generates a fullword by performing the indicated operation over corresponding bits of the two operands. For example, A\&B generates a fullword whose bit 0 is the result of A's bit 0 ANDed with B's bit 0 , and so forth for all 36 bits.

### 2.5.3 Polish (Complex) Expressions

MACRO cannot evaluate certain expressions containing relocatable values or EXTERNAL symbols. Instead MACRO generates special expressions called Polish expressions, which tell LINK how to resolve the values at load time. MACRO also generates Polish expressions to resolve inter-PSECT references.

For example, assume that $A$ and $B$ are externally defined symbols. Then MACRO cannot perform the operations $A+B-3$, but instead generates a special Polish block containing an expression to pass to LINK; the expression is equivalent to - +AB3. (See REL Block Type 11 in the LINK Reference Manual.) At load time, the values of $A$ and $B$ are available to LINK, and the expression is resolved.

NOTE
If you have used reverse Polish notation with a calculator, you should notice that although MACRO's Polish expressions are similar, they are not reversed. (These notations are called Polish because they were invented by the Polish logician Jan Lukasiewicz.)

### 2.5.4 Evaluating Expressions

2.5.4.l Hierarchy of Operations - MACRO has a hierarchy of operations in evaluating expressions. In an expression without nests (angle brackets), or within a nested expression, MACRO performs its operations in this effective order:

1. All unary operations and shifts: +, -, ^-, ${ }^{\wedge} \mathrm{D},{ }^{\wedge} \mathrm{O},{ }^{\wedge} \mathrm{B}, \mathrm{B}$ (binary shift), - (underscore shift), ${ }^{\wedge} \mathrm{F},{ }^{\mathrm{N}} \mathrm{L}, \mathrm{E}, \mathrm{K}, \mathrm{M}, \mathrm{G}$. Zeros are added for $K, M$, and $G$ before any other operation is performed.
2. Logical binary operations (from left to right): ! (OR), ^! (XOR), \& (AND).
3. Multiplication and division (from left to right): *, /.
4. Addition and subtraction (binary operations): +, -

You can override this hierarchy by using angle brackets to show what you want done first. For example, suppose you want to calculate the sum of $A$ and $B$, divided by $C$. You cannot do this with $A+B / C$ because MACRO will perform the division $B / C$ first, then add the result to $A$. With angle brackets you can write the expression $\langle A+B\rangle / C$, telling MACRO to add $A$ and $B$ first, then divide the result by $C$.

Expressions can be nested to any level. The innermost nest is evaluated first; the outermost, last. Some examples of legal expressions (assuming that Al, Bl, and C are defined symbols) are:
$A 1+B 1 / 5$
$A 1+B 1>/ 5$

- A1\&BI"!C
- B101M - mase


## NOTE'

```
An expression given in halfword notation
(that is, lefthalf,,righthalf) has each
half evaluated separately in a 36-bit
register. Then the 18 low-order bits of
each half are joined to form a fullword.
For example, the expression \langle4,,6\rangle/2
generates the value 000002 000003.
```

2.5.4.2 Evaluating Expressions with Relocatable Values - The value of an expression is usually either absolute or relocatable. Recall that relocatable values in your binary code will have the relocation constant added at load time by LINK.

Assume that $A$ and $B$ are relocatable symbols, and that $X$ and $Y$ are absolute symbols, and that the relocation constant is $k$. Let $a+k$ and $b+k$ be the values of $A$ and $B$ after relocation. Then $A+X$ makes sense (to LINK) because it means $\langle a+k\rangle+X$, which is the same as $\langle a+X\rangle+k$, clearly relocatable.

Since $X$ and $Y$ are both absolute, any operation combining them gives an absolute result.

Now look at the expression $A+B$. This means $\langle a+k\rangle+\langle b+k\rangle$, which is the same as <a+b>+2k, neither absolute nor relocatable. Similarly, $A * B$ means $\langle a+k\rangle *\langle b+k\rangle$, or $\langle a * b\rangle+\langle a+b\rangle * k+k * k$, again neither absolute nor relocatable. Such expressions cannot be evaluated by MACRO and are passed as Polish expressions to LINK.

More generally, you can see if an expression is absolute or relocatable by substituting relocated forms as above (for example, $a+k)$, and separating it (if possible) into the form
absolute+n*k
where absolute is an absolute expression. If $n=0$, the expression is absolute; if $n=1$, it is relocatable. If $n$ is neither 0 nor $l$, or if the expression cannot be put into the form above, then the expression is neither absolute nor relocatable. (Nevertheless, LINK will correctly evaluate the expression at load time.)

## ELEMENTS OF MACRO

## 2. 6 MACRO-DEFINED MNEMONICS

MACRO-defined mnemonics are words that MACRO recognizes and can translate to binary code. These mnemonics include:

1. Machine instruction mnemonics
2. I/O instruction mnemonics
3. I/O device code mnemonics
4. KLlO EXTEND instruction mnemonics
5. JRST and JFCL mnemonics
6. DECsystem-10 monitor call mnemonics
7. DECsystem-10 CALLI mnemonics
8. DECsystem-10 TTCALL mnemonics
9. DECsystem-10 MTAPE mnemonics
10. F40-switch-dependent mnemonics

Each type of mnemonic is discussed and tabulated in Appendix $C$. These mnemonics, together with MACRO's pseudo-ops and special characters, form the MACRO language.

## CHAPTER 3

PSEUDO-OPS

```
A pseudo-op is a statement that directs the assembler to generate code
or set switches to control assembly and listing of your program. For
example, the pseudo-op RADIX does not generate code, but it tells
MACRO how to interpret numbers in your program. The pseudo-op EXP
generates one word of code for each argument given with it.
To use a pseudo-op in your program, follow it with a space or tab, and
any required or optional arguments or parameters. The program
examples in Appendix D show pseudo-ops used in context.
This chapter describes the use and functions of each pseudo-op
(alphabetically). The headings included for each description, if
applicable, are:
1. FORMAT
2. FUNCTION
3. EXAMPLES
4. OPTIONAL NOTATIONS
5. RELATED PSEUDO-OPS
6. COMMON ERRORS
Some entries under COMMON ERRORS cite single-character error codes (for example, Merror). These codes are discussed in Section 8.2.
Many of the examples show some parts of the code assembled. The format and meaning of assembled code is discussed in Section 6.l.
```

ARRAY

| FORMAT | ARRAY sym[expression] |
| :---: | :---: |
|  | expression $=$ an integer value in the current radix, indicating the number of words to be allocated; the expression cannot be EXTERNAL, relocatable, or a floating-point decimal number, and its value must not be negative. |
| FUNCTION | Reserves a block of storage whose length is the value of the expression, and whose location is identified by the symbol. Storage is allocated along with other variable symbols in the program. |
|  | If the pseudo-op TWOSEG is used, ARRAY storage must be in the low segment. (See the VAR pseudo-op.) |
|  | The allocated storage is not necessarily zeroed. |
|  | If you use ARRAY in a PSECT, storage is allocated within that PSECT. |
|  | NOTE |
|  | Though the expression portion of an OPDEF must be in square brackets, this use of the brackets is completely unrelated to literals or literal handling. |
| EXAMPLES | AFERAY STAFTE200] |
|  | AFFAY FLLACE[1000] |
|  | AFFAY EFFSE 2000.7 |
| OPTIONAL | ARRAY syml,sym 2 [expression] |
|  | Both syml and sym 2 have a length equal to the value of the expression. |
| RELATED PSEUDO-OPS | BLOCK, . COMMON, INTEGER, VAR |
| COMMON ERRORS | Using an EXTERNAL symbol for name or size of the array (E error). |

## ASCII

| FORMAT | ASCII dtextd |
| :---: | :---: |
|  | d = delimiter; first nonblank character, whose second appearance terminates the text. |
|  | text $=$ string of text characters to be entered. |
| FUNCTION | Enters ASCII text in the binary code. Each character uses seven bits. Characters are left justified in storage, five per word, with bit 35 in each word set to 0 , and any unused bits in the last word set to 0 . |
| EXAMPLES | 105122122117122 ASCII /EFROR MESSAGE/ |
|  | 040115105123123 |
|  | 10.1107105000000 |
|  | $123124101 \quad 122124$ ASCIT !STAFTTNG AGAIN! |
|  | 111116107040101 |
|  | $107 \quad 101111.116000$ |
|  | 105116104123040 ASCII PENTIS WITH ZEFOS? |
|  | 1271111241110040 |
|  | $\begin{array}{lllllllllllllll}132 & 105 & 122 & 117 & 123\end{array}$ |
| OPTIONAL NOTATIONS | Omit the space or tab after ASCII. This is not allowed |
|  | if the delimiter is a letter, number, dot, dollar sign, or percent sign (that is, a possible symbol |
|  | constituent), or if the ASCII value of the delimiter character is less than 040 or greater than 172. |
|  | Right justified ASCII can be entered by using double quotes to surround up to five characters; for example, |
|  | 20101000000101 MOUEI AC1, "A" |
| RELATED PSEUDO-OPS | ASCIZ, . DIRECTIVE FLBLST, RADIX50, SIXBIT |
| COMMON ERRORS | Using the delimiter character in the text string. |
|  | Missing the end delimiter (that is, attempting to use a carriage return as a delimiter). |
|  | Using more than 5 characters in a right-justified ASCII string, or more than 2 characters if in the address field (Q error). |
|  | Giving direct assignment of a long ASCII string value to a symbol (for example A=ASCII /ABCDEFGH/). Only the first word (five characters, left justified) is assigned. |
|  | Using ASCII when ASCIZ is required. |

## ASCIZ

FORMAT

FUNCTION

EXAMPLES

OPTIONAL NOTATIONS

RELATED PSEUDO-OPS

## COMMON

 ERRORSASCIZ dtextd

```
    d = delimiter; first nonblank character, whose second
        appearance terminates the text.
    text = string of text characters to be entered.
    Enters ASCII text exactly as in the pseudo-op ASCII,
    except that a trailing null character is guaranteed.
    That is, if the number of characters in text is a
    multiple of five, a fullword of zeros is generated.
    105 122 122 117 122 ASCIZ /ERFOR MESSAGE/
040 115 105 123 123
101 107 105 OQO 000
123 124 101 122 124 ASCIZ !STARTING AGAIN!
111 116 107 040 101
107 101 111 116 000
105 116 104 123 040 ASCIZ PENLIS WITH ZEFOS?
127 1111124 110 040
132 105 122 117 123
000000 000 000 000
```

    Omit the space or tab after ASCIZ. This is not allowed
    if the delimiter is a letter, number, dot, dollar sign,
    or percent sign (that is, a possible symbol
    constituent), or if the ASCII value of the delimiter
    character is less than 040 or greater than 172 .
    ASCII, .DIRECTIVE FLBLST, RADIX50, SIXBIT
Using the delimiter character in the text string.
Missing the end delimiter (that is, attempting to use a
carriage return as a delimiter).
Giving direct assignment of a long ASCII string value
to a symbol (for example A=ASCII /ABCDEFGH/). Only the
first word (five characters, left justified) is
assigned.
In a macro, using a delimiter character that interferes
with recognition of a dummy-argument. For example, in
the macro
IUEFINE FOO(X)
ASCIZ. $X$.
$X$ is not seen as a dummy-argument because. $X$. is itself
a valid symbol.

## PSEUDO-OPS

## ASCIZ (Cont.)

```
In the macro
    IUEFINE FOO(X)<
    ASCIZ /X/
    >
X is seen as a dummy-argument because the slash (/) is
not valid in a symbol.
The macro
    IIEFINE FOO(X)<
    ASCIZ .'X'.
uses the concatenation operator (') to assure
recognition of X as a dummy-argument. (See Section 5.4
for a discussion on concatenating arguments.)
```

| FORMAT | . ASSIGN syml,sym 2 ,increment |
| :---: | :---: |
|  | syml and sym2 = global symbols. |
|  | increment $=$ expression with integer value. |
| FUNCTION | MACRO generates a REL Block Type 100. (See the LINK |
|  | Reference Manual.) At the time the program is loaded |
|  | into memory, assigns the value of sym2 to syml, and adds increment to sym 2 . |
|  | The . ASSIGN pseudo-op is useful for assigning a block of storage in one module and providing another module with the symbols needed to reference that block. |
| EXAMPLES | $\cdot A S S T G N A y F O=$ <br> A Assimbs the value of $F C$ to $A$, f then redefines the value of ; $F C$ to be $F C+E$. |
|  |  <br> ; EkRt, then rederimes EFRS to <br> ; be ERFR when the current <br> a value of ERNO. |
| OPTIONAL NOTATIONS | . ASSIGN syml, sym 2 |
|  | If the increment is missing, its value is 1 . |
| COMMON | Syml or sym 2 not global. |
| ERRORS |  |
|  | Increment not defined at assembly time. |

PSEUDO-OPS
$\square$

| FORMAT | ASUPPRESS |
| :--- | :--- |
| FUNCTION | Causes all local or INTERNAL symbols that are not <br>  <br> referenced after the ASUPPRESS to be deleted from |
|  | MACRO's symbol table at the end of Pass 2. These |
|  | symbols will not be output to LINK, will not be |
|  | available to the debugger, and will not appear in the |
|  | symbol table in the program listing file. |

BLOCK


|  | PSEUDO-OPS |
| :---: | :---: |
|  | BYTE |
| FORMAT | BYTE bytedef ... bytedef |
|  | bytedef $=(\mathrm{n})$ expression, ... expression |
|  | $\mathrm{n}=$ byte size in bits; n is a decimal expression in the range l to 36 . |
|  | expression = value to be stored. |
| FUNCTION | Stores values of expressions in $n$-bit bytes, starting at bit 0 of the storage word. The first value is stored in bits 0 to $n-1$; the second in bits $n$ to $2 n-1$; and so forth for each given value. |
|  | If a byte will not fit in the remaining bits of a word, the bits are zeroed and the byte begins in bit 0 of the next word. If a value is too large for the byte, it is truncated on the left. |
|  | If the byte size is 0 or is missing (empty parentheses), a zero word is generated. |
| EXAMPLES | 000002 UELOCY $=2$ |
|  | 050000.010502 EYTE (6)5,0,1101,5yUELOCY |
|  | generates the storage value 050000 010502. The two commas indicate a null argument; the lol (octal) is too large for the byte size and is left truncated. |
|  | ```070001007000 BYTE (6)7,0,1(9)7,0,1,"A" 001101000000``` |
|  | Notice that the code for "A" (101) is right justified in its 9-bit byte. |
| COMMON | Byte size too big (A error). |
|  | Missing left or right parenthesis (A error). |
|  | Extraneous comma before left parenthesis; the comma generates a null byte. |
|  | Using an EXTERNAL symbol or EXTERNAL complex expression for $n$ or expression. |


| FORMAT | COMMENT dtextd |
| :---: | :---: |
|  | $d=$ delimiter; the first nonblank character, whose second appearance terminates the text. |
|  | text $=$ text to be entered as a comment. |
| FUNCTION | Treats the text between the delimiters as a comment. The text can include a CR-LF to facilitate multiline comments, as shown below. |
| EXAMPLES | COMMENT /THIS IS A COMMENT THAT IS MORE THAN 1 LINE LONG/ |
| OPTIONAL | Omit the space or tab after COMMENT. This is not |
| NOTATIONS | allowed if the delimiter is a letter, number, dot |
|  | dollar sign, or percent sign (that is, a possible symbol constituent), or if the ASCII value of the |
|  | delimiter character is less than 040 or greater than |
|  | 172. |
|  | Use a semicolon (;) to make the rest of the line into a comment. |
| RELATED | REMARK |
| PSEUDO-OPS |  |
| COMMON | Using the delimiter character in the text string. |
| ERRORS |  |
|  | Missing the end delimiter (that is, attempting to use carriage return as a delimiter). |

## PSEUDO-OPS

|  | . COMMON |
| :---: | :---: |
| FORMAT | . COMMON symbol [expression] |
|  | symbol $=$ name of a FORTRAN COMMON block. |
|  | expression $=$ an expression having a positive integer value; this value defines the length of the COMMON block. |
| FUNCTION | Defines a FORTRAN or FORTRAN-compatible COMMON block. Causes the equivalent action of a FORTRAN labeled COMMON. (See the FORTRAN Programmer's Reference Manual.) |
|  | You can use . COMMON to define blank COMMON; to do this, use the symbol . COMM. as the name of the COMMON block. (Both FORTRAN and LINK recognize this as the name of blank COMMON.) |
|  | To define a COMMON block, MACRO generates a REL Block Type 20. (See the LINK Reference Manual.) |
|  | If used, the . COMMON pseudo-op must precede any MACRO statement that generates binary code, and must precede any other reference to the symbol name. |
| EXAMPLES | . COMMON LAATA1[50] |
| OPTIONAL NOTATIONS | . COMMON symbol,...,symbol[expression] |
|  | defines a COMMON array for each symbol given. Each array has a length equal to the value of the expression. |
| RELATED PSEUDO-OPS | ARRAY, BLOCK, EXTERN, INTEGER |
| COMMON ERRORS | Missing left or right square bracket (A error). |
|  | Using a relocatable value or EXTERNAL symbol in expression. |

## PSEUDO-OPS

. CREF

| FORMAT | . CREF |
| :---: | :---: |
| FUNCTION | Resumes output of cross-referencing that was suspended by the . XCREF pseudo-op. |
| OPTIONAL | Can apply to specific symbols to cancel a previous |
| NOTATIONS | . XCREF on those symbols, as in |
|  | . CREF symbol,...,symbol |
| RELATED | . XCREF |
| PSEUDO-OPG |  |
| COMMON | Specifying a nonexistent symbol (A error). |
| ERRORS |  |




## PSEUDO-OPS



```
.DIRECTIVE
```

| FORMAT | DIRECTIVE directive,..., directive |
| :---: | :---: |
| FUNCTION | Sets switches to enable or disable MACRO features. a directive has a logical opposite, you can use NO as a prefix to reverse the directive. The directives are: |
|  | .ITABM - include spaces and tabs as part. of passed arguments in macro call. |
|  | . XTABM - strip leading and trailing spaces and tabs from passed arguments in macro call. .XTABM is the default setting. |
|  | MACMPD - match paired delimiters in macro call. MACMPD is the default for assembly. It implies . XTABM and disables. ITABM. Using .DIRECTIVE NO MACMPD disables all quoting characters except angle brackets in macro arguments, and offers you a choice of . ITABM or . XTABM. |
|  | LITLST - list all binary code for literals in-line. |
|  | FLBLST - list only first line of binary code for multiline text. NO FLBLST is the default. |
|  | . OKOVL - allow overflow for arithmetic and for the pseudo-ops DEC, EXP, and OCT. |
|  | .EROVL - give an $N$ error for arithmetic overflow. .EROVL is the default. |
|  | MACPRF - prefer macro definition of symbol over other definitions of the same symbol. This does not affect the searching of .UNV files. |
|  | SFCOND - suppress source listing for failing conditional assembly. The lines containing the opening and closing angle brackets are not suppressed. |
|  | . NOBIN - do not generate binary (.REL) file. |
|  | KAlO - enter KAlO as CPU type in header block of binary file. |
|  | KI10 - enter KIlO as CPU type in header block of binary file. |
|  | KLlO - enter KLlO as CPU type in header block of binary file. |
| EXAMPLES | - HFEECTIVE: MACMFI, NOEIN |
| COMMON ERRORS | Using NO with a directive that does not have a logical opposite. |


| FORMAT | END expression |
| :---: | :---: |
|  | expression $=$ an optional operand that specifies the address of the first instruction to be executed; can be EXTERNAL. |
| FUNCTION | Must be the last statement in a MACRO program. |
|  | Statements after END are ignored. The starting address |
|  | is optional and normally is given only in the main |
|  | program. (Since subprograms are called from the main program, they need not specify a starting address.) |
|  | When the assembler first encounters an END statemen |
|  | it terminates Pass 1 and begins Pass 2. The END |
|  | terminates Pass 2 on the second encounter, after which |
|  | the assembler simulates XLISTed LIT and VAR statements |
|  | beginning at the current location. (In a PSECTed |
|  | program, the LIT and VAR statements are simulated for each PSECT.) |
| EXAMPLES | ENI STAFT |
|  | START is a label at the starting address. |
| OPTIONAL | Use the END statement to specify a transfer word in |
| NOTATIONS | some output file formats. (See pseudo-ops RIM, RIMIO, and RIMIOB in Appendix E.) |
| RELATED | PRGEND |
| PSEUDO-OPS |  |
| COMMON | Failing to end a text string or literal with a |
| ERRORS | closing delimiter; MACRO cannot see the END statement. |
|  | Including an END statement in a source file when it is |
|  | not the last file in a group of files you want |
|  | assembled as a single program. |
|  | Closing the input file immediately after the characters "END" with no following carriage return. |



## ENTRY

| FORMAT | ```ENTRY symbol,...,symbol symbol = name of an entry point in a library subroutine.``` |
| :---: | :---: |
| FUNCTION | Defines each symbol in the list following the ENTRY pseudo-op as an INTERNAL symbol and places them in a REL Block Type 4 at the beginning of the . REL output file. If this . REL file is later included in an indexed library of subroutines, then the symbol will also be included in a REL Block Type 14 at the beginning of the library. (Except for this, ENTRY is equivalent to INTERN.) |
|  | If LINK is in library search mode, a subroutine will be loaded if the program to be executed contains an undefined global symbol that matches a name in the library entry list for that program. |
|  | Since library subroutines are external to programs using them, the calling program must list them in EXTERN statements. |
| EXAMPLES | If the MATRIX subroutine is a library subroutine, it must contain the statement |
|  | ENTFY MATEIX |
|  | in order to make the symbol MATRIX available to other programs. In addition, it must define the symbol MATRIX as a label at the address where execution of the call is to begin: |
|  | MATFIX: |
| RELATED PSEUDO-OPS | INTERN, EXTERN |
| COMMON ERRORS | Not defining the symbol in the program. |
|  | Purging an ENTRY symbol in Pass 2 only. The ENTRY symbol is normally output at the beginning of Pass 2; a PURGE of an ENTRY symbol must occur in Pass 1 to be effective. |

$\square$ ExP


## EXTERN

| FORMAT | EXTERN symbol,...,symbol |
| :---: | :---: |
| FUNCTION | Identifies symbols as being defined in other programs. EXTERNAL symbols cannot be defined within the current program. |
|  | At load time, the value of an EXTERNAL symbol is resolved by LINK if you load a module that defines the symbol as an INTERNAL symbol. (If you do not load such a module, LINK gives an error message for the undefined EXTERNAL symbol.) |
|  | An EXTERNAL symbol cannot be used for any program values affecting address assignment (such as arguments to LOC or RELOC). |
|  | For a discussion of global symbols and their resolution by LINK, see Section 2.4.5.2. |
| EXAMPLES | EXTEFN SQFT, CUBE, TYFE |
| OPTIONAL NOTATIONS | Suffix \#\# to the symbol. This declares the symbol |
|  | EXTERNAL, and eliminates the need for the EXTERN |
|  | pseudo-op. Most programmers who use the \#\# notation do |
|  | site that the symbol is EXTERNAL. |
|  | For example, the two statements |
|  | EXTEFN A |
|  | ATWO $=$ A*2 |
|  | can be simplified to |
|  | ATWO: A \# \# ${ }_{2}{ }^{2}$ |
| RELATED | INTERN, ENTRY, UNIVERSAL |
| PSEUDO-OPS |  |
| COMMON <br> ERRORS | Attempting to declare a symbol as EXTERNAL after its |
|  | first use has made it local (by default) or INTERNAL (by declaration). |
|  | Declaring a symbol as EXTERNAL in a program that searches a UNIVERSAL file that gives a conflicting definition. |



| FORMAT | . IF expression, qualifier, <code> |
| :---: | :---: |
| FUNCTION | Gives criterion and code for conditional assembly. The code is assembled if: |
|  | qualifier is AND expression is |
|  | ABSOLUTE absolute |
|  | ASSIGNMENT a direct-assignment symbol |
|  | ENTRY a symbol given in ENTRY pseudo-op |
|  | EXTERNAL an EXTERNAL symbol |
|  | INTERNAL an INTERNAL or ENTRY symbol |
|  | GLOBAL a global symbol |
|  | LABEL a label |
|  | LOCAL a local symbol |
|  | LRELOCATABLE a lefthalf relocatable symbol |
|  | MACRO a macro name |
|  | NEEDED an undefined but referenced symbol |
|  | NUMERIC numeric |
|  | OPCODE an opcode |
|  | OPDEF a symbol defined by OPDEF pseudo-op |
|  | REFERENCED a symbol already in the symbol table |
|  | RELOCATABLE a relocatable symbol |
|  | RRELOCATABLE a righthalf relocatable symbol |
|  | SYMBOL a symbol (instead of a number) |
|  | SYNONYM a symbol defined by SYN pseudo-op |
|  | NOTE |
|  | If the expression has different properties in Pass 1 and Pass 2, the number of words of code generated may be different for each pass. |
| EXAMPLES | - IF FOO, MACRO, FFOO, |
| OPTIONAL | Abbreviate qualifier up to unique initial letters. |
| NOTATIONS | For example, you can abbreviate OPCODE to OPC, but not to OP, since OPDEF has the same first two letters. |
|  | Omit the comma preceding the left angle bracket. |
| $\begin{aligned} & \text { RELATED } \\ & \text { PSEUDO-OPS } \end{aligned}$ | . DIRECTIVE SFCOND, .IFN, IFx group |
| $\begin{aligned} & \text { COMMON } \\ & \text { ERRORS } \end{aligned}$ | Omitting the comma between expression and qualifier. |
|  | Mismatching angle brackets. |
|  | Misplacing the .IF statement in such $a$ way that the property given by the qualifier is different in Pass l and Pass 2. For example, the following code generates phase errors in Pass 2: |
|  | -IF FOO, OFMEF , MFCL OFDEF FOO[JFST] |

NXTLAB: ENI

```
.IFN
```

| FORMAT | . IFN expression,qualifier, <code> |
| :---: | :---: |
| FUNCTION | Gives criterion and code for conditional assembly. The code is assembled if: |
|  | qualifier is AND expression IS NOT |
|  | ABSOLUTE absolute |
|  | ASSIGNMENT a direct-assignment symbol |
|  | ENTRY a symbol given in ENTRY pseudo-op |
|  | EXTERNAL an EXTERNAL symbol |
|  | INTERNAL an INTERNAL or ENTRY symbol |
|  | GLOBAL a global symbol |
|  | LABEL a label |
|  | LOCAL a local symbol |
|  | LRELOCATABLE a lefthalf relocatable symbol |
|  | MACRO a macro name |
|  | NEEDED an undefined but referenced symbol |
|  | NUMERIC numeric |
|  | OPCODE an opcode |
|  | OPDEF a symbol defined by OPDEF pseudo-op |
|  | REFERENCED a symbol already in the symbol table |
|  | RELOCATABLE a relocatable symbol |
|  | RRELOCATABLE a righthalf relocatable symbol |
|  | SYMBOL a symbol (instead of a number) |
|  | SYNONYM a symbol defined by SYN pseudo-op |
|  | NOTE |
|  | If the expression has different properties in Pass 1 and Pass 2, the number of words of code generated may be different for each pass. |
| EXAMPLES | , IFN FOO, OFIEF, GOFDEF FOO[270日8]: |
| OPTIONAL NOTATIONS | Abbreviate qualifier up to unique initial letters. For example, OPCODE can be abbreviated to OPC, but not to |
|  | OP, since OPDEF has the same first two letters. |
|  | Omit the comma preceding the left angle bracket. |
| RELATED PSEUDO-OPS | . DIRECTIVE SFCOND, . IF, IFx group |
| COMMON ERRORS | Omitting the comma between expression and qualifier. |
|  | Mismatching angle brackets. |
|  | Misplacing the. IFN statement in such a way that the property given by the qualifier is different in Pass l and Pass 2. For example, the following code generates phase errors in Pass 2: |
|  | -IFN FOO,OFDEF, UFCL OFIEF FOOLJFSTI |

```
IFx group
```

FUNCTION

```
Gives criterion and code for conditional assembly. A
symbol or expression used to define the conditions for
assembly must be defined before MACRO reaches the
conditional statement. If the value of such a symbol
or expression is not the same on both assembly passes,
a different number of words of code may be generated,
and a phase error can occur.
The forms of the IF pseudo-op are listed below; in the
first six forms, n is the value of the given
expression.
IFE expression,<code> - assemble code if n=0.
IFN expression,<code> - assemble code if n}=0\mathrm{ .
IFG expression,<code> - assemble code if n>0.
IFGE expression,<code> - assemble code if n\geqslant0.
IFL expression,<code> - assemble code if n<0.
IFLE expression,<code> - assemble code if n\leqslant0.
IFl <code> - assemble code on Pass l.
IF2 <code> - assemble code on Pass 2.
IFDEF symbol,<code> - assemble code if the symbol is
    defined as user-defined, an opcode, or a
    pseudo-op.
IFNDEF symbol,<code> - assemble code if the symbol is
    not defined as user-defined, an opcode, or a
    pseudo-op. Code is also assembled if the symbol
    has been referenced, but is not yet defined. This
    can occur during pass 1.
IFIDN <stringl><string2>,<code> - assemble code if the
    strings are identical.
IFDIF <stringl><string2>,<code> - assemble code if the
    strings are different.
```


## NOTES

1. For IFIDN and IFDIF, the assembler compares the two strings (interpreted as ASCII) character by character.
2. The IFIDN and IFDIF pseudo-ops usually appear in macro definitions, where one or both strings are dummy-arguments.
(Continued on next page)

|  | PSEUDO-OPS |
| :---: | :---: |
|  | IFx group (Cont.) |
|  | IFB <string>,<code> - assemble code if the string contains only blanks and tabs. |
|  | IFNB <string>, <code> - assemble code if the string does not contain only blanks and tabs. |
| EXAMPLES | $\% \% C C==\% \% C C+1$ FIncrement character count <br> IFG $\% \% C C-5, \% \% C C=0$ ; Word overflowed? <br> $\% W C=\% \% W C+1 \%$ ;Yes, to newt word |
| OPTIONAL NOTATIONS | Omit angle brackets enclosing code for single-line conditionals. |
|  | Omit the comma preceding the code if the code is enclosed in angle brackets. |
|  | For IFIDN, IFDIF, IFB, and IFNB only: use a nonblank, nontab character other than < as the initial and terminal delimiters for a string (as in pseudo-ops ASCII and ASCIZ). You can then include angle brackets in the string. |
| RELATED PSEUDO-OPS | . DIRECTIVE SFCOND, .IF, .IFN |
| COMMON | Comparison string too large (A error). |
|  | Mismatched angle brackets. |
|  | EXTERNAL symbol used for comparison (E error). |
|  | String not properly delimited. |
|  | Missing comma with single-line conditional. |



| FORMAT | INTERN symbol,..., symbol |
| :---: | :---: |
| FUNCTION | Declares each given symbol to be INTERNAL global; |
|  | therefore its definition, which must be in the current |
|  | program, is available to other programs at load time. |
|  | Each such symbol must be defined as a label, a variable, or a direct-assignment symbol. |
|  | MACRO builds a list of symbol definitions that will be available to other programs at load time. |
|  | OPDEF symbols can be declared INTERNAL, and thus be made available to other programs at load time. |
|  | However, if the current program has another symbol (besides the OPDEF symbol) of the same name, the |
|  | INTERNAL declaration will apply to that symbol rather than to the OPDEF symbol. |
| EXAMPLES | INTERN SQUARE, CEFOOT, TYFES |
| OPTIONAL | TAG: $\quad$; INTERNAL label |
| NOTATIONS | VALUE=:expression ;INTERNAL direct assignment |
| RELATED | EXTERN, ENTRY |
| PSEUDO-OPS |  |
| COMMON ERRORS | Failing to define an INTERNAL symbol in the curren |
|  | program. |
|  | Using INTERN for a library entry point (when ENTRY is required). |

```
FORMAT IOWD expl,exp2
    exp1, exp2 = expressions.
FUNCTION Generates one I/O transfer word in a special format for
    use in BLKI and BLKO and all five pushdown instructions
    (ADJSP, PUSH, POP, PUSHJ, POPJ). The left half of the
    assembled word contains the 2's complement of the value
    of expl, and the right half contains the value exp2-l.
EXAMPLES
OPTIONAL XWD - expl,exp2-1
NOTATIONS
COMMON
ERRORS
    The following line shows how IOWD 6,^D256 places -6
    (octal 777772) in the left halfword and 256 (octal 377)
    in the right halfword:
    777772 000377 IOWN 6,"n256
    The following lines show IOWD STL,STK used in a
        literal. The LIT pseudo-op then shows the code
        generated in the literal pool.
        000017 F===17
        000001 AC1==1
        000100 STL ==100
        STK: ELOCK STL
        20017000001053' MOUE F,[IOWM STL,STK.T
        261 17000 000001 FUSH F,AC1
        25400000001054, JFST ENN
        . . . . . .
        LIT
7 7 7 7 0 0 ~ 0 0 0 0 0 1
10400000 000170 EN[I: HAL.TF
-expl,,exp2-1
Using a relocatable expression for expl (R error).
```



| FORMAT | ```IRPC darg,<code> darg = one of the dummy-arguments of the enclosing macro definition. (IRPC can only be used in the body of a macro definition.)``` |
| :---: | :---: |
| FUNCTION | Generates one expansion of code for each character of the string that replaces darg. Each occurrence of darg within the expansion is replaced by the character currently controlling the expansion. (See Section 5.6.) |
|  | Concatenation and line continuation are not allowed across end-of-IRPC, since a carriage return and linefeed are appended to each expansion. See the example below. |
| EXAMPLES | IEFINE $A(B)$ IRFC E,GASCTZ $\backslash B \backslash$ A(STRING) MFFC |
|  | 123000000000000 ASCTZ \S |
|  | 124000000000000 ASCIZ \T |
|  | 122000000000000 ASCIZ \FV |
|  | 111000000000000 ASCIZ \IT |
|  | 116000000000000 ASCIZ \N |
|  | 107000000000000 ASCIZ \G\ |
| $\begin{aligned} & \text { RELATED } \\ & \text { PSEUOD-OPS } \end{aligned}$ | IRP, STOPI |
| COMMON | IRPC NOT IN A MACRO (A ERROR). |
|  | Argument is not a dummy symbol (A error). |
|  | Argument is a created symbol (A error). |
|  | Mismatched angle brackets. |


| FORMAT | LALL |
| :---: | :---: |
| FUNCTION | Causes the assembler to print in the program listing |
|  | file everything that is processed, including all text |
|  | and macro expansions. Since XALL is the default, you |
|  | must use LALL if you want full macro expansions listed. |
|  | This can be helpful in debugging a program. |
|  | LALL does not produce comments in a macro expansion if |
|  | the comments are preceded by double semicolons (;i). |
|  | This is because such comments are not stored. |
| OPTIONAL | Use the /E switch described in Table 7-1. |
| NOTATIONS |  |
| RELATED | LIST, SALL, XALL, XLIST |
| PSEUDO-OPS |  |


| FORMAT | . LINK chain-number, store-address,chain-address <br> chain-number $=$ a positive integer expression that associates the link with others having the same number. <br> store-address $=$ a symbol giving the store address for this entry in the chain. <br> chain-address $=$ an optional integer expression giving the address of this entry in the chain. If you omit the chain-address, MACRO generates a 0 and LINK uses the store-address as the chain-address. |
| :---: | :---: |
| FUNCTION | Generates static chains at load time. MACRO generates a REL Block Type 12. (See the LINK Reference Manual for a full discussion of LINK's handling of these chains.) |
| EXAMPLES | See the LINK Reference Manual (REL Block Type 12) for extensive examples of using . LINK and .LNKEND. |
| RELATED PSEUDO-OPS | . LNKEND |
| COMMON | Chain-number not absolute (A error). |
| ERRORS | EXTERNAL expression for store-address or chain-address (E error). |


| FORMAT | LIST |
| :--- | :--- |
| FUNCTION | Resumes listing following an XLIST statement. The LIST <br> function is implicitly contained in the END statement. |
| OPTIONAL <br> NOTATIONS | Use the /L switch described in Table 7-1. |
| RELATED <br> PSEUDO-OPS | LALL, SALL, XALL, XLIST |

```
LIT
```

FORMAT
LIT
FUNCTION
Assembles literals beginning at the current address. The literals assembled are those found since the previous LIT, or since the beginning of the program, whichever is later. The location counter is incremented by 1 for each word assembled.

In a PSECTed program, LIT assembles only literals in the current PSECT.

A literal found after the LIT is not affected. It will be assembled at the next following LIT, or at the END statement, whichever is earlier.

At the END statement, unassembled literals are placed in open-ended storage after the end-of-program. If data is also to be entered in open-ended storage, literals stored there may be overwritten. (See Appendix $F$ for a discussion of storage allocation.) This possibility is avoided by using LIT before the END statement.

Assembling literals with LIT also produces a listing of their binary code. Literals unassembled at the END are XLISTed.

Literals having the same value are collapsed in MACRO's literal pool. Thus for the statements:

FUSH Fig[O]
FUSH F:[0]
MOUEI AC1,[ASCIZ/TESTI/]
the same address is shared by the two literals [0], and by the null word generated at the end of [ASCIZ /TESTl/]. Literal collapsing is suppressed for those literals that contain errors, undefined expressions, or EXTERNAL symbols.

NOTES

1. If the code immediately preceding a LIT does not cause a transfer of execution control to some other location, execution will "fall into" the literal pool, producing unpredictable results.
2. In a file containing PRGEND pseudo-ops, only one LIT is permitted in each module before the last one. The last module (containing the END statement), or any file without PRGENDs, can contain multiple LITs.
(Continued on next page)


```
LOC
```

| FORMAT | LOC expression |
| :---: | :---: |
|  | expression $=$ an optional operand whose value gives the address at which sequential address assignment is to continue. |
| FUNCTION | Sets the location counter to the value of the expression and begins assigning absolute addresses to the instructions and data following the LOC instruction. |
|  | If no address is specified, the location counter is restored to its value previous to the last LOC pseudo-op or RELOC-RELOC sequence. (See example below.) If no previous LOC pseudo-op was encountered, the assumed address is 0 . |
|  | To switch to relocatable address mode, use the pseudo-op RELOC. If no argument is specified, RELOC (in this context) restores the location counter to its value previous to the LOC pseudo-op or LOC-LOC sequence. (An implicit RELOC 0 begins each program.) |
|  | If an entire program is to be assigned absolute locations, a LOC statement must precede all instructions and data. |
|  | Note that, unlike RELOC-RELOC sequences, typically used to switch between segments in a two-segment program, LOC-LOC sequences cannot be successfully interrupted and then resumed. This is demonstrated in the example below. |
| EXAMPLES | $400000^{\prime}$ TWOSEG 400000 ;set uF hises  <br> $000000^{\prime}$ FELOC Back to lowses |
|  | 000010 LOC 10 iset lif liol - Loc |
|  | 000010 000000 000001 NEC 1,2 |
|  | 000011000000000002 |
|  | 000100 LOC 100 |
|  | 000100000000000003 LIEC 3.4 |
|  | 000101000000000004 |
|  | 000012 L.OC |
|  | 000102 L.OC |
|  |  |
|  | FBut we can't resume LOC-LOC <br> 400000' FELOC |
|  | 000102 LOC |
|  | 000102 LOC |
|  | 400000, FELOC $\quad$;But FELOC- |
|  | 000102 LOC |
| RELATED <br> PSEUDO-OPS | RELOC, .ORG, TWOSEG |
| COMMON ERRORS | Using an EXTERNAL expression for the address expression (E error). |


| FORMAT | .MFRMT |
| :--- | :--- |
| FUNCTION | Causes multiformat listing of binary code. The type of <br> instruction assembled determines this format. (See <br> Section 6.l.) .MFRMT is the default setting. |
| OPTIONAL <br> NOTATIONS | Use the $/ F$ switch described in Table 7-1. |
| RELATED <br> PSEUDO-OPS | .HFRMT |



| MLON |
| :---: |


| FORMAT | MLON |
| :--- | :--- |
| FUNCTION | Suspends the effect of an earlier MLOFF pseudo-op, <br> thereby enabling the use of multiline literals. MLON <br> is the default setting. |
| RELATED <br> PSEUDO-OPS | MLOFF |


| FORMAT | .NODDT symbol, ..., symbol |
| :---: | :---: |
| FUNCTION | Suppresses debugger output of each given symbol. Each symbol must have been previously defined. Symbols suppressed with . NODDT can include OPDEF symbols. |
| EXAMPLES | . NOHLIT CALL, F.JFigT,Fo |
| OPTIONAL NOTATIONS | Use $==$ for direct-assignment symbols. (See Section |
|  | 2.4.2.2.) |
|  | Use : ! for label symbols. (See Section 2.4.2.1.) |
| RELATED PSEUDO-OPS | PURGE |
| COMMON ERRORS | Using . NODDT with an undefined symbol argument. |

$\square$

| FORMAT | NOSYM |
| :--- | :--- |
| FUNCTION | Suppresses listing of the symbol table in the program <br> listing file. |
|  | Suppressing the listing of symbol tables is useful for <br> a library file containing many PRGENDs. |


| FORMAT | OCT expression, ..., expression |
| :--- | :--- |
| FUNCTION | Defines the local radix for the line as octal; the <br> value of each expression is entered in a fullword of <br> code. The location counter is incremented by l for |
| each expression. |  |

OPDEF

| FORMAT | OPDEF symbol [expression] |
| :---: | :---: |
| FUNCTION | Defines the symbol as an operator equivalent |
|  | expression, giving the symbol a fullword value. When |
|  | the operator is later used with operands, the |
|  | accumulator fields are added, the indirect bits are |
|  | ORed, the memory addresses are added, and the index |
|  | register addresses are added. |
|  | An OPDEF can be declared INTERNAL, using the INTERN |
|  | pseudo-op. However, if a symbol of the same name |
|  | exists, the INTERNAL declaration will apply only to |
|  | that symbol, and not to the OPDEF. |

NOTES

1. If you use a relocatable symbol in defining an OPDEF, the value of the symbol may not be the same for all references to the OPDEF.
2. Though the expression portion of an OPDEF must be in square brackets, this use of the brackets is completely unrelated to literals or literal handling.

EXAMPLES

RELATED
PSEUDO-OPS
COMMON
ERRORS

200062 OOOO10 OFDEF CAL [MOUE 1, (2SYM(2)]
20002104000014 CAL 1ッBOL (2)
The CAL statement is equivalent to:
20002104000014 MOVE $2,05 Y M+B O L(4)$
DEFINE, SYN

OPDEF of macroname or SYN symbol (A error).
No code generated by statement in square brackets (A error).

Missing square brackets (A error).

```
PSEUDO-OPS
```

.ORG

PAGE

| FORMAT | PAGE |
| :--- | :--- |
| FUNCTION | Causes the assembler to list the current line and then <br> skip to the top of the next listing page. The subpage <br> number is incremented, but the page number is not. |
| OPTIONAL | A formfeed character (CTRL/L) in the input text |
| NOTATIONS | has a similar effect, but increments the page number <br> and resets the subpage number. |


| FORMAT | PASS 2 |
| :---: | :---: |
| FUNCTION | Switches the assembler to Pass 2 processing for the remaining code. All code preceding this statement will have been processed by Pass l only; all following code by Pass 2 only. |
|  | You can use PASS2 to reduce assembly time during debugging; you can also use PASS2 to omit the second pass for a UNIVERSAL file containing only symbol definitions (OPDEFs, macros, and direct assignments). |
| EXAMPLES | Testing a macro defined in the Pass 1 portion: |
|  | IFE NON, <br> FRINTX PHORFIELE EFROR FASS2 <br> ENII |
|  | $\geqslant$ |
|  | stops assembly if NON $=0$. |



| FORMAT | ```PHASE address address = an integer expression; cannot be an EXTERNAL symbol.``` |
| :---: | :---: |
| FUNCTION | Assembles part of a program so that it can be moved other locations for execution. To use this feature, the subroutine is assembled at sequential relocatable or absolute addresses along with the rest of the program, but the first statement before the subroutine is PHASE, followed by the address of the first location of the block into which the subroutine is to be moved prior to execution. All address assignments in the subroutine are in relation to the address argument. The subroutine is terminated by DEPHASE, which restores the location counter. |
| EXAMPLES | In the following example, which is the central loop in a matrix inversion, a block transfer instruction moves the subroutine LOOP into accumulators 11 to 15 for execution. (This results in faster execution on KAlO and KIlO processors.) |
|  | 002000' $20000000402002^{\prime \prime}$ MAIN: MOUE [XWI LOOFX,LOOF] |
|  | 002001, 25100000000015 ELT LOOF+4 |
|  | 002002' 25400000000011 JFST LOOF |
|  | 000011 LODFX: FHASE 11 |
|  | 00001121002003000002 LOOF: MOUN AC, $\mathrm{A}(\mathrm{X})$ |
|  | 00001216002000000100 FMF AC,MFYF |
|  | 00001314202004000002 FALM AC, A (Y) |
|  | 00001436503000000011 SOJGE X, . 3 |
|  | $00001525400000002000^{\circ}$ JTSST MAIN |
|  | 002010' DEFHASE |
|  | The label LOOP represents accumulator 11 , and the . -3 in the SOJGE instruction represents accumulator 11. |
|  | Note that the code inside the PHASE-to-DEPHASE program segment is loaded into the address following the previous relocatable code; all labels inside the |
|  | segment, however, have the address corresponding to the phase address. Thus the phased code, if it contains |
|  | control transfers other than skips, cannot be executed |
|  | until it has been moved (for example, by a BLT |
|  | instruction) to the address for which it was assembled. |
| RELATED | DEPHASE |
| PSEUDO-OPS |  |
| COMMON | Using an EXTERNAL symbol or complex EXTERNAL expression |
| ERRORS | as the address (E |


| FORMAT | POINT bytesize,address,bitplace |
| :---: | :---: |
| FUNCTION | Generates a byte pointer word for use with the machine language mnemonics ADJBP, LDB, IBP, ILDB, and IDBP. |
|  | Bytesize gives the decimal number of bits in the byte, and is assembled in bits 6 to 11 of the storage word. |
|  | Address gives the location of the byte word, and is assembled in bits 13 to 35. Bitplace gives the position (in decimal) of the rightmost bit of the byte. |
|  | MACRO places the value 35 minus bitplace in bits 0 to 5 of the storage word. |
|  | If the address is indirect, bit 13 is set. If the address is indexed, the index is placed in bits 14 to |
|  | 17. The default bytesize is 0. The default bitplace |
|  | is -l, so that the byte increment instructions IBP, |
|  | ILDB, and IDBP will begin at the left of the address word. |
| EXAMPLES | 3606000000000 FOINT 6.0,5 |
|  | 4406000000100 FOINT 6,100 |
| COMMON <br> ERRORS | Bytesize or bitplace not given in decimal. |
|  | Bytesize or bitplace not absolute. |
|  | Bytesize or bitplace EXTERNAL. |

## PRGEND

FORMAT
FUNCTION

OPTIONAL NOTATIONS

RELATED
PSEUDO-OPS
COMMON ERRORS

PRGEND
Replaces the END statement for all except the last program of a multiprogram assembly. PRGEND closes the local symbol table for the current module.

You can use PRGEND to place several small programs into one file to save space and disk accesses. The resulting binary file can be loaded in search mode. (See the LINK Reference Manual.)

Using PRGEND requires extra memory for assembly, since the tables for each program must be saved for Pass 2. Functionally, however, PRGEND is identical to END, except that PRGEND does not end the current assembly pass.

NOTE

1. PRGEND is not allowed in macros or PSECTS.
2. PRGEND clears the TWOSEG pseudo-op.
3. Like END, PRGEND causes assembly of all unassembled literals and variable symbols.
4. In a file containing PRGENDs, using more than one LIT pseudo-op in any but the last program produces unpredictable results.

Give an argument with PRGEND, specifying the start address for the program. See the END pseudo-op for a discussion of this argument and its meaning.

END, LIT, VAR

Failing to end a text string, REPEAT, conditional code, DEFINE, or literal with a closing delimiter; MACRO cannot see any following PRGEND or END.

Confusing multiprogram and multifile assemblies. A multiprogram assembly involves multiple programs separated by PRGENDs. A multifile assembly always involves multiple files separated by end-of-file. The two types of assemblies are not mutually exclusive.

PSEUDO-OPS

PRINTX

| FORMAT | PRINTX text |
| :---: | :---: |
| FUNCTION | Causes text to be output during assembly. On Pass 1 the text is output to the terminal and the listing device. On Pass 2 the text is output to the terminal, but only if the terminal is not the listing device. |
|  | PRINTX is frequently used to output conditional information and, in very long assemblies, to report progress of the assembler through Pass 1. |
| EXAMPLES | FRINTX ASSEMELER HAS REACHEL FOINT NOWGO |
|  | IFGE .-1000, FFINTX COLE MORE THAN 1 P\% |

```
FORMAT .PSECT name/attribute,origin
    name = a valid symbol giving the name of the PSECT.
    attribute = either CONCATENATE or OVERLAID.
    origin = an expression giving an address for the PSECT
        origin.
    Specifies the relocation counter to be used for the
    code following. MACRO generates a REL Block Type 23.
    (See the LINK Reference Manual.)
    Do not use PRGEND and .PSECT in the same file. MACRO
    will treat the first PRGEND as an END statement and
    ignore any following source code.
    For a complete discussion of PSECTS and their handling,
    see Section 9.l.3.
EXAMPLES
OPTIONAL
NOTATIONS
RELATED
PSEUDO-OPS
Using TWOSEG and .PSECT in the same module.
Using HISEG and .PSECT in the same module.
```

| FORMAT | ```PURGE symbol,...,symbol symbol = an assigned symbol, a label, an operator, or a macro name.``` |
| :---: | :---: |
| FUNCTION | Deletes symbols from the symbol tables. Normally used at the end of a program to conserve storage and to delete symbols for the debugger. Purged symbol table space is reused by the assembler. |
|  | If you use the same symbol for both a macro name or OPDEF and a label, a PURGE statement deletes the macro name or OPDEF. Repeating the instruction then purges the label. |
|  | Purging a symbol that is EXTERNAL or undefined suppresses any error messages associated with it. |
| EXAMPLES | 000040 O00001. $\quad$ LAEEL: 1,1 |
| $\begin{aligned} & \text { RELATED } \\ & \text { PSEUDO-OPS } \end{aligned}$ | . NODDT, XPUNGE |

## RADIX



## RADIX50

| FORMAT | RADIX50 code, symbol |
| :---: | :---: |
| FUNCTION | Packs the symbol into bits 4 to 35 of the storage word, with the code in bits 0 to 3. |
|  | The "50" in RADIX50 is octal, so that the radix in |
|  | decimal is 40. The 40 characters permitted in symbols |
|  | are the "digits" of the RADIX50 symbol expression. |
|  | Thus a symbol is seen by RADIX50 as a "6-digit" number |
|  | in base 40 , converted to binary, and placed in bits 4 to 35 in storage. |
|  | The code expression for RADIX50 is a number in the range 0 to 74 octal. Its binary equivalent should end |
|  | with two zeros (that is, the octal should end with 0 or 4), since the two low-order bits will not be stored. |
|  | The four high-order bits are placed in bits 0 to 3 in storage. |
|  | See Appendix A for the octal values of RADIX50 characters. |
| EXAMPLES | 126633472376 FALIIX50 10, SYMEOL |
|  | 466633472376 FAAIIX 5044 , SYMEOL. |
| OPTIONAL NOTATIONS | The mnemonic SQUOZE can be used in place of RADIX50. |
|  |  |
|  | RADIX50 , symbol (code is taken as zero). |
| RELATED | SQUOZE |
| PSEUDO-OPS |  |
| COMMON | RADIX50 code not absolute (A error). |
| ERRORS |  |
|  | RADIX50 code does not end with 0 or 4 (Q error). |

## RELOC

| FORMAT | RELOC expression |  |  |
| :---: | :---: | :---: | :---: |
|  | expression $=$ an optional operand that specifies the address at which sequential address assignment is to continue. |  |  |
| FUNCTION | Sets the location counter to and begins assigning rel instructions and data that | to the value ocatable follow. | e of expression, addresses to the |
|  | In a PSECTed program, RELOC for the current PSECT. | sets the | location counter |
|  | If no address is specified, the location counter is restored to its value before the last RELOC, or before the last LOC-LOC sequence, whichever is later. (See the first example below.) If no previous RELOC or LOC-LOC sequence was encountered, the location counter is set to 0 . |  |  |
|  | An implicit RELOC 0 begins switch to absolute address | every MAC mode, use | CRO program. To the pseudo-op LOC. |
|  | Note that RELOC-RELOC sequ switch between segments in be interrupted and then resu in the first example below. | quences (ty in a two-seg umed. This | ypically used to gment program) can s is demonstrated |
| EXAMPLES | $400000^{\prime}$ | TWOSEG 400 | 0000 ; Set uf hises |
|  | 000000' | FELOC | ; Back to lowses |
|  | 000000'000000 000001 | LEC 1.2 |  |
|  | 000001'000000 000002 |  |  |
|  | $400000^{\prime}$ | FEELOC | ; Back to hises |
|  | $400000^{\prime}$ 25500 000000000 | JFCL |  |
|  | 000137 | LOC 137 | flemosit version <br> ; ir absolute 137 |
|  | $000137 \quad 000100 \quad 000001$ | XWN 100,1 |  |
|  | $400001^{\prime}$ | FELOC: | ;Back to hises <br> ; where left off |
|  | 400001, $25400000400000^{\prime}$ $000002^{\prime}$ | $\underset{\text { FIELOC }}{\text { JFST }}-1$ | F Back to lowses |
| RELATED PSEUDO-OPS | LOC, . ORG, TWOSEG |  |  |

PSEUDO-OPS


| FORMAT | REMARK text |
| :--- | :--- |
| FUNCTION | Text is a comment. |
| EXAMPLES | REMAFK I CAN SAY ANYTHING HERE. |
| OPTIONAL <br> NOTATIONS | A comment line can also begin with a semicolon. |
| RELATED <br> PSEUDO-OPS | COMMENT |
| COMMON <br> ERRORS | Continuing REMARK text to next line without using the <br> continuation character (CTRL/underscore). |

```
REPEAT
```

| FORMAT | REPEAT expression, <code> |
| :--- | :--- |
|  | expression the repeat index, which gives the number |
| of times to repeat assembly of the code given; |  |
| the repeat index can be any expression having a |  |
| nonnegative integer value. |  |


| FORMAT | . REQUEST filespec |
| :---: | :---: |
| FUNCTION | Causes the specified file to be loaded only to satisfy a global request; that is, the file is loaded in library search mode. (See Chapter 7 for a discussion of files.) |
|  | The filespec must not include a file extension. If you specify a path, only the project-programmer number is allowed; SFDs are not allowed. |
|  | MACRO generates a REL Block Type 17. (See the LINK Reference Manual.) |
| EXAMPLES | - FEQUEST KSK:MACROS <br> -FEQUEST MACFOOS |
| OPTIONAL NOTATIONS | DSK: is the default device. |
|  | Your default path at load time is the default path. |
| RELATED | . REQUIRE, .TEXT |
| PSEUDO-OPS |  |

## PSEUDO-OPS

.REQUIRE

```
FORMAT .REQUIRE filespec
FUNCTION Causes the specified file to be loaded automatically,
    independent of any global requests. (See Chapter 7 for
    discussion of files.)
    The filespec must not include a file extension. If you
    specify a path, only the project-programmer number is
    allowed; SFDs are not allowed.
    MACRO generates a REL Block Type 16. (See the LINK
    Reference Manual.)
EXAMPLES .FEQUIFE IISK:MACROS
    . REQUIFE MACROS
    - FEQUIFE SYG:MACREL
OPTIONAL DSK: is the default device.
NOTATIONS
RELATED
    Your default path at load time is the default path.
PSEUDO-OPS
    .REQUEST, .TEXT
```

| FORMAT | SALL |
| :--- | :--- |
| FUNCTION | Causes suppression of all macro and repeat expansions <br> and their text; only the input file and the binary <br> generated will be listed. SALL can be nullified by <br> either XALL or LALL. Using SALL generally produces the <br> tidiest listing file. |
| OPTIONAL <br> NOTATIONS | Use the /M switch described in Table 7-1. |
| RELATED <br> PSEUDO-OPS | LALL, LIST, XALL, XLIST |


|  | SEARCH |
| :---: | :---: |
| FORMAT | SEARCH tablename(filename),...,tablename (filename) |
| FUNCTION | Defines a list of symbol tables for MACRO to search if a symbol is not found in the current symbol table. A maximum of ten tables can be specified. Tables are searched in the order specified. |
|  | When the SEARCH pseudo-op is seen, MACRO checks its internal UNIVERSAL table for a memory-resident UNIVERSAL of the specified name. (See the UNIVERSAL pseudo-op for further discussion of memory-resident UNIVERSAL tables and use of the /U switch.) |
|  | If no such entry is found in the UNIVERSAL table, MACRO reads in the symbol table using the given file specification. If no file specification is given, MACRO reads tablename. UNV from the default path. If no such file is found, MACRO then tries UNV:tablename.UNV and SYS:tablename.UNV, in that order. |
|  | When all the specified files are found, MACRO builds <br> table for the search sequence. If MACRO cannot find a given symbol in the current symbol table, the UNIVERSAL tables are searched in the order specified. When the symbol is found, it is moved into the current symbol table. This procedure saves time (at the expense of core) on future references to the same symbol. |
|  | A UNIVERSAL file can search other UNIVERSAL files, provided all names in the search list have been assembled. |
|  | The internal table of UNIVERSAL names is cleared on each run (.R MACRO) or START command, but is not cleared when MACRO responds with an asterisk. |
|  | In a PSECTed program, all UNIVERSAL symbols belong to the blank PSECT. |
| EXAMPLES | SEAFCH MONSYM, MACSYM |
| OPTIONAL NOTATIONS | Omit the filename and its enclosing parentheses. MACRO then looks on DSK:, UNV:, and SYS: (in that order) for tablename.UNV. |
| RELATED PSEUDO-OPS | UNIVERSAL |
| COMMON | Not purging a macro that redefines itself (P error). |
| ERRORS | If a macro is found in a universal file, the definition is copied into the current macro table and the auxiliary table is not searched on Pass 2. Thus, a macro that redefines itself can cause $P$ errors similar to enclosing the macro by IFl. Such macros should be purged before Pass 2. |


| FORMAT | SIXBIT dtextd |
| :---: | :---: |
|  | ```d = delimiter; first nonblank character, whose second appearance terminates the text.``` |
| FUNCTION | Enters strings of text characters in 6-bit format. Six characters per word are left justified in sequential storage words. Any unused bits are set to zero. |
|  | Lowercase letters in SIXBIT text strings are treated as uppercase. Otherwise, only the SIXBIT character set is allowed. (See Appendix A for SIXBIT characters and their octal codes.) |
| EXAMPLES | 644570640063 SIXETT \TEXT STRING\} |
|  |  |
|  | 644570 640000 EXF SIXBIT /TEXT/ |
| OPTIONAL | Omit the space or tab after SIXBIT. This is not |
| NOTATIONS | allowed if the delimiter is a letter, number, dot, |
|  | dollar sign, or percent sign (that is, a possible symbol constituent), or if the ASCII value of the |
|  | delimiter character is less than 040 or greater than |
|  | 172. |
|  | Right-justified SIXBIT can be entered by using single quotes to surround up to six characters; for example, |
|  | 006251475064 'FIGHT' |
| RELATED | ASCII, ASCIZ, . DIRECTIVE FLBLST |
| COMMON | Using the delimiter character in the text string. |
|  | Missing the end delimiter (that is, attempting to use a carriage return as a delimiter). |
|  | Using more than six characters in a right-justified |
|  | SIXBIT string, or more than three characters if in the address field (Q error). |
|  | Using non-SIXBIT characters in the text string. |


| FORMAT | SQUOZE code,symbol |
| :---: | :---: |
| FUNCTION | SQUOZE is a mnemonic for RADIX50. |
| EXAMPLES | 126633472376 RAIIX50 10,SYMEOL <br> 126633472376 SQUOZE 10,SYMEOL |
| OPTIONAL NOTATIONS | RADIX50 code,symbol <br> SQUOZE , symbol (code is taken as 0). |
| RELATED PSEUDO-OPS | RADIX50 |
| COMMON Errors | Code not absolute (A error). <br> Code does not end with 0 or 4 (Q error). |

## STOPI

```
FORMAT
FUNCTION
EXAMPLES
RELATED IRP, IRPC
PSEUDO-OPS
COMMON
ERRORS
```

STOPI

Ends an IRP or IRPC before all subarguments or characters are used. The current expansion is completed, but no new expansions are started. STOPI can be used with conditionals inside IRP or IRPC to end the repeat if the given condition is met.

LALL
प्TEFINE ONETWO(A)
IFF A\&IFILNAACONE, ©STOFI
EXF $1 \gg$
IFF A, SIFIINACTWOQ, STOFI
EXF $2 \mathbf{2}$
$>$
ONETWO $A, B, D \subset$
IRF
IFIDNAACONES,STOFI

IF ILINGOONE, ©STOFI

IF I IN OMONE , STOFI

IFF
IFIDN©ACTWO, ©STOFI

IFILINESTWO\%,STOFI

IFTDN؟ITWO\%, STOFI
EXF $2>$
$\cdots$
ONETWO GA,ONE,B,ONE,TWO:IFF
IFIIN A A ONE , ©STOFI
EXF 13
IFIDNCONE OONE YSTOFI

IFF
IFIIN世A) TWO:, STOFI
EXF 2\%
IFIDNONESTWO\%, GTOFI
EXF 2 -
IFILINBETWO\%,STOFI
EXF 2 :
IFILNOONESTWO, GTOFI
EXF $2 \boldsymbol{y}$
IFIDNSTWO®TWO\%, GTOFI
000000 EXF 2)

IRP, IRPC

STOPI not inside IRP or IRPC.

```
FORMAT SUBTTL subtitle
FUNCTION Defines a subtitle (of up to 80 characters) to be
    printed at the top of each page of the listing file
    until the end-of-listing or until another SUBTTL
statement is found.
    The initial SUBTTL usually appears on the second line
    of the first page of the input file, immediately
    following the TITLE statement.
    For subsequent SUBTTL statements, the following rule
    applies: if the new SUBTTL is on the first line of a
new page, then the new subtitle appears on that page;
if not, the new subtitle appears on the next page.
```

NOTE
The statements

- ••

FFigenil
TITLE FOO
GUBTTL EAF
do not cause BAR to appear as the subtitle on the first page of the listing of FOO.

SUBTTL affects only the listing file, and subtitles can be changed as often as desired.

SUBTTL THIS SECTION CONTAINS IEUTCE MEFENLENT FOUTINES
TITLE

## SUPPRESS

| FORMAT | SUPPRESS symbol,..., symbol |
| :---: | :---: |
| FUNCTION | Turns on a suppress bit in the symbol table for the |
|  | specified symbols. The suppress bit will be turned off for any symbol later referenced in the program. |
|  | Symbols whose suppress bits are on at the end of |
|  | assembly are not listed in the symbol table, but will |
|  | be listed in any tables built by CREF unless they are XCREFed. |
|  | When an appended parameter file (as opposed to a |
|  | UNIVERSAL file) is used in an assembly, many symbols |
|  | may be defined but never used. These take up space in |
|  | the binary file and complicate listing of the file. |
|  | Unused and unwanted symbols can be removed from tables |
|  | by SUPPRESS or ASUPPRESS. These pseudo-ops control the |
|  | suppress bit in each entry of the symbol table; if the |
|  | bit is on, the symbol in that location is not output. |
| RELATED | ASUPPRESS |
| PSEUDO-OPS |  |
| COMMON | Attempting to suppress an undefined symbol. |
| ERRORS |  |

## SYN

```
FORMAT SYN syml,sym2
    sym1 = a defined symbol.
    sym2 = a symbol to be defined as synonymous with syml.
    Defines sym2 as synonymous with syml.
    If syml is defined as both a label and an operator,
    sym2 assumes the label definition.
    The following are legal SYN statements:
    SYN X,K゙
    GYN FAII,AIMD
    SYN ENI,XENI
    To turn XLIST into a null operator,
    IUEFINE,XL < >
    SYN ,XL,XLIST
    To restore its operation,
    FURGE XLIST
RELATED
PSEUDO-OPS
COMMON
ERRORS
DEFINE, OPDEF
Missing symbol (A error).
Unknown symbol - first operand not defined (A error).
Missing comma (A error).
Using a variable as one of the symbol arguments (A
error).
```

PSEUDO-OPS

TAPE

| FORMAT | TAPE |
| :---: | :---: |
| FUNCTION | Causes the assembler to begin assembling the program contained in the next source file in the MACRO command string. |
| EXAMPLES | (Interactive) |
|  | - F MACRO |
|  | *ISK : EINAME, LFTT: =TTY: , IISK: MORE FAF:AM=6 |
|  | TAFE |
|  | ;THIS COMMENT WIIL EE IGNOFED |
|  | $\cdots$ |
|  | This sets PARAM to 6 and assembles the remainder of the program from the source file DSK:MORE. Since MACRO is a two-pass assembler, the TTY: file must be repeated for Pass 2. |
|  | [MCREF1 ENI OF PASS 1] |
|  | F'ARAM $=6$ |
|  | TAFE |
|  |  |
|  | Note that all text after the TAPE pseudo-op is ignored. |


| FORMAT | . TEXT dtextd |
| :---: | :---: |
|  | ```d = delimiter; first nonblank character, whose second appearance terminates the text.``` |
| FUNCTION | Generates an ASCIZ REL Block Type for LINK and inserts the text string directly into the . REL file output as a separate block. (See the LINK Reference Manual.) |
|  | The text inserted in the . REL file is interpreted as a command string for LINK. Therefore a MACRO program loaded by user commands to LINK can contain additional LINK commands, carried out when the MACRO program is loaded. |
| EXAMPLES | .TEXT / SET:.HIGH.:500000' |
| OPTIONAL | Omit the space or tab after . TExT. This is not allowed |
| NOTATIONS | if the delimiter is a letter, number, dot, dollar sign, or percent sign (that is, a possible symbol constituent), or if the ASCII value of the delimiter character is less than 040 or greater than 172. |
| RELATED | . REQUEST, . REQUIRE |
| PSEUDO-OPS |  |
| COMMONERRORS | Using the delimiter character in the text string. |
|  | Missing the end delimiter (that is, attempting to use a carriage return as a delimiter). |


| FORMAT | TITLE title |
| :--- | :--- |
| FUNCTION | Gives the program name and a title to be printed at the <br> top of each page of the program listing. |
|  | The first characters (up to six characters, or up to |
|  | the first non-RADIX50 character) are the program name. |
|  | This name is used when debugging with DDT to gain |


| FORMAT | TWOSEG expression <br> expression $=$ any expression giving a nonnegative value <br> as the beginning of the program high segment; |
| :--- | :--- |
| cannot be EXTERNAL. |  |

## UNIVERSAL

FORMAT

FUNCTION

UNIVERSAL tablename

Declares the symbol table of the current program available to other programs, and stores the given tablename in MACRO's internal UNIVERSAL table. The tablename is also taken as the program name, and appears in the heading of each page of the listing file.

When an END or PRGEND statement is found, the symbol table is placed immediately after the assembler's pushdown stacks and buffers. In addition to this memory-resident copy of the UNIVERSAL symbol table, the file tablename. UNV is generated. (This file can be suppressed by the /U switch described in Table 7-1.)

UNIVERSAL files can be used to generate data, but are more commonly used to generate symbols, macros, and OPDEFs. The symbols and OPDEFs generated in a UNIVERSAL program need not be declared INTERNAL, since its local symbols are available to accessing programs. (See the SEARCH pseudo-op.)

Memory-resident UNIVERSAL symbol tables are cleared on each run (.R MACRO) or START, but are not cleared when MACRO responds with an asterisk. This saves redundant lookups when many programs search a common set of UNIVERSALs.

Note that if a sequence of programs (or even one program) searches more than ten UNIVERSAL symbol tables, a SEARCH table overflow occurs. This overflow forces reinitialization of the assembler by a run (.R MACRO) or START command.

For a UNIVERSAL program that does not generate data (that is, it has only symbol, macro, and OPDEF definitions), you can save time by using l-pass assembly. However, such a file must not contain forward references to symbol definitions.

A UNIVERSAL file cannot contain PSECTs.
(Continued on next page)

```
PSEUDO-OPS
```

```
UNIVERSAL (Cont.)
```

NOTES

1. For COMPILE-class commands, the existence of the file tablename.REL may prevent recompilation of the UNIVERSAL file tablename. MAC. To avoid this, force compilation of the . MAC file by including /COMPIL in the command string.
2. Generally, a UNIVERSAL file need not be reassembled when referencing programs are assembled with newer versions of MACRO. However, if the UNIVERSAL's assembler version is newer than the program's, you may get the MCRUVS message, indicating skewed UNIVERSAL versions. In this case, reassembly or one or both files is required (using the same assembler version).
```
EXAMPLES UNIUEFSAL. S1
    START=765
    AC1=1
    F=O
    ENII
RELATED SEARCH, TITLE
PSEUDO-OPS
COMMON Using TITLE and UNIVERSAL in the same module (M error).
ERRORS
```




FORMAT
FUNCTION

RELATED
PSEUDO-OPS
OPTIONAL NOTATIONS

XALL
Resumes standard listing after previous LALL or SALL. (XALL is the default among these three.)

XALL suppresses all lines of the program listing file that do not generate binary code.

XALL does not suppress REPEAT expansions.

NOTE
Under XALL only one listing line is output for each source line generating binary code in a macro expansion. Occasionally, a single line of a macro definition expands into several lines of listing text. When this occurs, part of a binary-generating source line may not be listed.

You can avoid this by temporarily setting the listing mode to LALL (list all) or SALL (suppress all) around such lines.

LALL, LIST, SALL, XLIST

Use the /X switch described in Table 7-1.

## . XCREF

| FORMAT | . XCREF symbol,..., symbol |
| :--- | :--- |
| FUNCTION | Suspends output of cross-referencing for the specified <br> symbols. References to these symbols between this |
| statement and the next .CREF or the end of the program |  |
| will not appear in the cross-reference listing. |  |

XLIST

| FORMAT | XLIST |
| :---: | :---: |
| FUNCTION | Suspends output to the program listing file. This output occurs only in Pass 2; XLIST does not affect Pass 1. To resume output, use the pseudo-op LIST. |
| EXAMPLES | The following sequence of code shows an XLIST pseudo-op suppressing listing of literals: |
|  | EXIT FEnd of erosram XLIGT FHon't list literals LIT LIST ENI |
|  | This sequence of code lists as: |
|  |  |
|  | Note that the high-segment break will be greater than 401023' because the literals are assembled after the HALTF. |
| RELATED <br> PSEUDO-OPS | LALL, LIST, SALL, XALL |
| OPTIONAL NOTATIONS | Use the /S switch described in Table 7-1. |

## PSEUDO-OPS

XPUNGE

| FORMAT | XPUNGE |
| :--- | :--- |
| FUNCTION | Deletes all local symbols during Pass 2. This reduces <br> the size of the . REL file and speeds up loading. <br> XPUNGE should immediately precede the END statement. |
| RELATED <br> PSEUDO-OPS | PURGE |


| FORMAT | XWD lefthalf,righthalf |
| :---: | :---: |
| FUNCTION | Enters two halfwords in a single storage word. Each half is formed in a 36-bit register, and the low-order 18 bits are placed in the halfword. The high-order bits are ignored. |
|  | XWD statements are used to set up pointer words for block transfer instructions. Block transfer pointer words contain two l8-bit addresses; the left half is the starting location of the block to be moved, and the right half is the first location of the destination. |
| EXAMPLES | $402017^{\prime}$ 200 02 0 00 $403040^{\prime}$ MOVE 2,[XWI FFOM1, TO1] <br> $402020^{\prime}$ 251 02 0 00 403035 ELT 2,TOEND1 |
|  |  |
| OPTIONAL NOTATIONS | lefthalf, righthalf |
|  | BYTE (18) lefthalf,righthalf |
| COMMON ERRORS | Using halfword with absolute value larger than 18 bits (Q error). |
|  | Using two commas between the arguinents to XWD. For example, XWD A,3 is correct; XWD A, 3 is incorrect. |

PSEUDO-OPS



## CHAPTER 4

MACRO STATEMENTS AND STATEMENT PROCESSING

A MACRO statement has one or more of the following: a label, an operator, one or more operands, and a comment. The general form of a MACRO statement is:
label: operator operand,operand ;comment
A carriage return ends the statement.

NOTES

1. Direct-assignment statements receive special handling. (See Section 2.4.2.2.)
2. Processing of macros is not discussed here because a macro call produces a text substitution. After substitution, the text is processed as described in this chapter. Macros are discussed in Chapter 5.

### 4.1 LABELS

A label is always a symbol with a suffixed colon. (See Section 2.4.2.1.) The assembler recognizes a label by finding the colon. If a statement has labels (you can use more than one), they must be the first elements in the statement.

A label can be defined only once; its value is the address of the first word of code generated after it.

Since a label gives an address, the label can be either absolute or relocatable. A label is a local symbol by default. You can declare a label INTERNAL global or EXTERNAL global. (See Section 2.4.5.)

### 4.2 OPERATORS

After processing any labels, the assembler views the following nonblank, nontab characters as a possible operator. An operator is one of the following:

1. A MACRO-defined mnemonic. All mnemonics are listed in Appendix $C$, and are discussed in the Hardware Reference Manual.
2. A user-defined operator. (See the pseudo-op OPDEF in Chapter 3.)
3. A pseudo-op. (See Chapter 3.)

If the characters found do not form one of the above, then MACRO views them as an expression.

An operator is ended by the first non-RADIX50 character: if it is ended by a blank or tab, operands may follow; if it is ended by a semicolon, there are no operands and the comment field begins; if it is ended by a carriage return, the statement ends and there are no operands or comments.

### 4.3 OPERANDS

After processing labels and the operator, if any, the assembler views as operands all characters up to the first unquoted semicolon or carriage return. Commas delimit the operands.

The operator in a statement determines the number (none, one, two or more) and kinds of permitted or required operands. Any expected operand not found is interpreted as null. An operand can be any expression or symbol appropriate for the operator.

### 4.4 COMMENTS

The first unquoted semicolon in a statement begins the comment field. You can use any ASCII characters in a comment; however, angle brackets in a comment may produce unpredictable results. You can continue a comment to the next line by typing CTRL/ , followed by a carriage return.

If the first nonblank, nontab character in a line is a semicolon, the entire line is a comment. You can also enter a full line of comment with the pseudo-op REMARK, or a multiline comment with the pseudo-op COMMENT. (See Chapter 3.)

Comments do not affect binary program output.

### 4.5 STATEMENT PROCESSING

MACRO processes your program as a linear stream of data. During Pass 1. MACRO may find references to symbols not yet defined. These symbols are entered in the user symbol table. Whenever a symbol is defined, it is entered in the table with its value, so that on Pass 2 all definitions can be found in the table. The values then replace the symbols in the binary code generated.

## NOTE

Delayed definition is allowed only for labels and direct-assignment symbols. A symbol that contributes to code generation (for example, an OPDEF, a macro, or a REPEAT index) must be defined before any reference to it.

Statement processing proceeds as follows:

1. Labels are found and entered in the user symbol table.
2. The next characters up to the first unquoted semicolon, blank, tab, comma, or equal sign are processed.
a. Equal sign: the characters form a symbol, and the following characters form an expression. The symbol and the value of the expression are entered in the user symbol table.
b. Other delimiter: the characters form an expression or an operator. If an operator, it is found in a table and assembled. If an expression, its value is assembled.
3. If the operator takes operands, the next characters up to the first unquoted semicolon or carriage return form operands. Unquoted commas delimit operands. For each operand, leading and trailing blanks and tabs are ignored. Operands are evaluated and assembled for the given operator.
4. The first unquoted semicolon ends processing of the line. Any further characters up to the first carriage return are comment.
5. The first unquoted carriage return ends the statement. Any following characters begin a new statement.

### 4.6 ASSIGNING ADDRESSES

MACRO normally (and by default) assembles statements with relocatable addresses. Assembly begins with the zero storage word and proceeds sequentially. Each time MACRO assembles a word of binary code, it increments its location counter by $l$.

A mnemonic operator generates one word of binary code. Directassignment statements and some pseudo-ops do not generate code. Some pseudo-ops generate more than one word of code.

You can control address assignment by setting the assembler's location counter using the pseudo-ops LOC and RELOC. (See Section 9.1.)

You can also reference addresses relative to the location counter by using the dot symbol (.). For example, the expression. -1 used as an address refers to the location immediately preceding the current location.

In revising MACRO programs, you can cause an incorrect address to be assembled by adding or removing statements within the range of a . +n expression. For example, in the sequence

```
000000' 33200001 000000 SKIFE O(AC)
000001' 254 00 0 00 001020' .JFST GOTONE
000002' 34401 0 00000000, AO.JA AC,...2
```

the expression . -2 gives the address of the SKIPE statement. If you revise this sequence by inserting a statement, you should change the expression to . -3 so that it still refers to the correct statement.

| $000000^{\prime}$ | 332 | 00 | 0 | 01 | 000000 | SKIFE O(AC) |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 000001' | 254 | 00 | 0 | 00 | 001020' | JFiST GOTONE |  |
| 000002' | 350 | 00 | 0 | 00 | 000014 | AOS NULCNT | \#Added line |
| 000003 ' | 344 | 01 | 0 | 00 | 000000' | AOJA AC, - 3 | ¢ Chansed line |

For this reason, use great care with such expressions other than . +1 and .-l. Using labels avoids this problem entirely.

### 4.7 MACHINE INSTRUCTION MNEMONICS AND FORMATS

There are two $k$ inds of machine instruction mnemonics: primary and input/output. Primary instructions generate binary code in primary instruction format; input/output instructions generate binary code in input/output instruction format.

### 4.7.1 Primary Instructions

A primary instruction is in one of the forms
mnemonic accumulator, address
mnemonic accumulator, mnemonic address
where mnemonic is a machine instruction mnemonic, accumulator is an accumulator register address, and address is a memory address. The memory address can be modified by indexing, indirect addressing, or both.

## MACRO STATEMENTS AND STATEMENT PROCESSING

A complete list of machine instruction mnemonics and their octal codes is given in Appendix $C$, and these mnemonics are discussed in the Hardware Reference Manual.

The accumulator address gives the address of a register, and can be any expression or symbol whose value is an integer in the range 0 to 17 octal.

The memory address gives a location in memory, and can be any expression or symbol whose value is an integer in the range 0 to octal 777777 .

You can modify the memory address by indirect addressing, indexed addressing, or both. For indirect addressing, prefix an at sign (@) to the memory address in your program. For indexed addressing, suffix an index register address in parentheses to the memory address in your program. This address can be any expression or symbol whose value is an integer in the range 1 to octal 17.

NOTE

> To assemble the index, MACRO places the index register address in a fullword of storage, swaps its halfwords, and then adds the swapped word to the instruction word.

For an example of a primary instruction (assuming that ACl7, TEMP, and XR have the octal values 17,100 , and 3 , respectively), the statement

ADK AC17, DTEMF (XF)
generates the binary code

|  | ruction code. | indirect bit |  |  |  |  |  | memory <br> address |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 010 | 111000 | 1111 | 1 | 0 | 011 | 000 | 000 | 000 | 001 | 000 | 000 |
|  |  | umulat | $\begin{aligned} & \text { index } \\ & \text { register } \end{aligned}$ |  |  |  |  |  |  |  |  |

which appears in the program listing as

```
27017103000100 ALM AC17,(OTEMF'(XF)
```

The mnemonic ADD has the octal code 270 , and this is assembled into bits 0 to 8. The accumulator goes into bits 9 to 12 . Since the @ appears with the memory address, bit 13 is set to l. The index register goes into bits 14 to l7. Finally, the memory address is assembled into bits 18 to 35 .

If any element is missing from a primary instruction, zeros are assembled in its instruction word field.

## MACRO STATEMENTS AND STATEMENT PROCESSING

### 4.7.2 Mnemonics With Implicit Accumulators

A few mnemonics set bits in the accumulator field as well as in the instruction field. Therefore these mnemonics do not take accumulator operands, and are of the form
mnemonic address
These mnemonics and their octal codes are listed in Table $C-5$ in Appendix $C$.

For example, the mnemonic JFOV gives the octal code 25504 ; JFCL gives 255. Therefore both give the opcode 255 in bits 0 to 8 , but JFOV also sets the accumulator bits (9 to 12) to binary 0001. This makes JFOV 100 equivalent to JFCL 1,100 :

```
25501000000100 JFOU 100
25501000000100 JFFCL 1,100
```


### 4.7.3 Input/Output Instructions

An input/output statement in your program resembles a primary instruction statement except that the first operand gives a device number instead of an accumulator. The general format is:
mnemonic device,address
In an input/output instruction, the indirect, index, and address fields (bits 13 to 35 inclusive) are assembled exactly as in a primary instruction.

Unlike a primary instruction word, however, an input/output word has a split instruction code in bits 0 to 2 (always set to 111 binary) and 10 to l2, and a device code in bits 3 to 9 . The device code can be any expression or symbol giving a valid device code for your system.
(MACRO-defined I/O instruction mnemonics and device code mnemonics are listed in Tables $\mathrm{C}-2$ and $\mathrm{C}-3$ in Appendix C. )

For example (assuming that NVR has the octal value 1037), the statement

DATAI CDR,@NVR(4)
generates the binary code

which appears in the listing as
711404104001037 [IATAI CLFPONUF (4)

The octal code for the mnemonic DATAI is 70004 , which is written in bits 0 to 14 . The octal device code 114 (for card reader) is then overwritten in bits 3 to 9. The @ in the statement sets bit l3 to l. The index register and memory address are placed in bits 14 to 17 and 18 to 35 , as in a primary instruction.

### 4.7.4 Extended Instructions

The KLl0 Extended Instruction Set is a multifunction instruction set that performs character-string editing, decimal-to-binary conversion, string move with left or right justification, string move with offset or translation, and string compare.

The Extended Instruction Set consists of a single kLlo instruction (EXTEND, octal l23) and a set of 16 extended operators. (See the Supplement to the Hardware Reference Manual.)

The KLIO EXTEND instruction mnemonics are listed in Table $C-4$ in Appendix C .

## CHAPTER 5

USING MACROS

A macro is a sequence of statements defined and named in your program. When you call a macro (by invoking its name in your program), the sequence of statements from its definition is generated in line, replacing the call. A macro can have arguments.

By using macros with arguments, you can generate passages of code that are similar, but whose differences are controlled by the passed arguments. This saves repetition in building a source file.

### 5.1 DEFINING MACROS

Before you can call a macro, you must define it. You can also redefine a macro if you wish; the new definition simply replaces the old one.

To define (or redefine) a macro, use the pseudo-op DEFINE:
DEFINE macroname (darglist)<macrobody>
where macroname is the name of the macro, darglist is an optional list of dummy-arguments, and macrobody is a sequence of statements.

The macroname is a symbol; you must follow the rules for valid symbols in selecting a macroname. (See Section 2.4.1.)

The optional dummy-argument list can give one or more dummy-argument symbols through which values are passed to the sequence of statements. If a macro definition has dummy-arguments, they must be enclosed in parentheses. Use commas as delimiters between dummy-arguments. For each dummy-argument, leading and trailing spaces and tabs are ignored.

The macrobody is the sequence of statements you want to generate when you call the macro. The macrobody must be enclosed in angle brackets.

Here is an example of a macro definition:

## USING MACROS

| DEFINE | UMAG (WHERE,LENG) |  |
| :---: | :---: | :---: |
|  |  |  |
|  | MOVE O, WHERE | GGet first <br> ; comforment |
|  | FMF 0 | ; Sauare it |
|  | MOUE 1, WHEFEE 1 | ;Get second <br> ; comporient |
|  | FMF 1,1 | - Square it |
|  | FAII 1 | ; Add scuare <br> ; of second |
|  | MOUE1, WHEFET? | gGet thiros <br> ; comporient |
|  | FMF 1, 1 | FSatare it |
|  | FALI 1 | ;Add sauare <br> ; of third |
|  | FUSHJ 17,FSQFit | ;Floatiris SQRT <br> ; routirie |
|  | MOUEM l.ENG | iStore the <br> ; lerigth |
| $>$ |  |  |

NOTE
Comments in a macro use storage. If you
begin a comment with a double semicolon,
the comment is listed in the definition
but not stored for listing with
expansions.

### 5.2 CALLING MACROS

You can call a macro by putting its name in your program. Recall that you must define the macro before you can call it. You can use the macroname as a label, an operator, or an operand.

If the macro's definition has dummy-arguments, the macro call can have arguments. The arguments passed to the macro are inserted into the defined sequence of statements as it is generated. The first passed argument replaces the first dummy-argument; the second passed argument replaces the second dummy-argument; this treatment continues for each argument passed. Any missing arguments are passed as nulls (zeros) or filled in by default arguments (see Section 5.5).

NOTE
If $F O O$ is a macro with four dummyarguments, the call FOO A, C passes A and $C$ as the first and third arguments. The second argument is passed as nulls; it is not considered missing and cannot be replaced by a default argument. The fourth argument is missing and will be replaced by a default argument if one has been defined; otherwise it is passed as nulls. (See Section 5.5.1.)

## USING MACROS

After argument substitution, the defined sequence of statements replaces the macroname and argument list in the source text. For example, suppose you have defined VMAG(A,B) as shown in Section 5.l above, and VMAG appears in your program as

```
    LALL
    P7=245
    ULEN=11
    FLACE=15
    MOUEM FLACE ;FiJt it in FL.ACE
    TAG2: UMAG FLACE,VLEN
    TAG3: MOVE 1,VLEN ;Get lensth
```

    TAG1: MOUE 1,FP
    Then the code to be assembled is:


Notice that the macro definition has the dummy-arguments $A$ and $B$ in the macrobody. The call VMAG PLACE,VLEN causes PLACE to replace each appearance of $A$, and VLEN to replace each appearance of $B$.

NOTES

1. Under LALL, when the text of a macrobody is listed at call, it is enclosed in up-arrows (").
2. Under XALL, the beginning of the text of a macrobody is marked by an up-arrow; the ending is marked by an up-arrow only if the last line of the macrobody generates binary code.

## USING MACROS

### 5.2.1 Macro Call Format

In a macro call, delimit the macroname with one or more blanks or tabs.

If the macro has arguments, the first nonblank, nontab character begins the argument list. Each argument ends with a comma, a carriage return, or a semicolon. These three characters cannot be used within arguments unless enclosed by special quoting characters. (See Section 5.2.2.)

Leading and trailing spaces and tabs are stripped from each argument unless they are within special quoting characters. Embedded spaces and tabs are not stripped.

You can continue an argument to the next line by using CTRL/underscore. Otherwise an unquoted carriage return or semicolon ends the argument and the argument list. An unguoted semicolon also begins the comment field.

### 5.2.2 Quoting Characters in Arguments

The special quoting characters for macro argument handing are:
< > angle brackets
( ) parentheses
[ ] square brackets
" " quote marks

NOTE
Single quote marks (apostrophes) are not
special quoting characters.

Any character, including the semicolon (;), enclosed in special quoting characters is treated as a regular character. If one of the special quoting characters is to be passed as a regular character, it must be enclosed by different special quoting characters.

Here are the rules for macro argument handing. In the examples, FOO is assumed to be a defined macro:

1. The special quoting characters are not argument delimiters. They only tell the assembler to treat the enclosed characters as regular characters.

FOO $C\langle A, B\rangle$ has one argument: $C\langle A, B\rangle$.
FOO $C, D\langle A, B\rangle$ has two arguments: $C$ and $D\langle A, B\rangle$.
2. With the two exceptions explained below, special quoting characters are always included in passed arguments.

FOO $A,(B, C)$ has two arguments: $A$ and ( $B, C$ ).
FOO [XWD 1,Ll]-1 (AC) has one argument: [XWD l,Ll]-l(AC).
FOO "(",0 has two arguments: "(" and 0.
Exception l: If the first character of the argument list is a left parenthesis, then it and its matching right parenthesis delimit the argument list. They are not treated as special quoting characters and are not included in passed arguments. All nested quoting characters except angle brackets are disabled. After stripping the outer parentheses, angle brackets are handled as described in Exception 2 below.

FOO ( $\mathrm{A}, \mathrm{B}, \mathrm{C}$ ) has three arguments: $\mathrm{A}, \mathrm{B}$, and C .
FOO (?LENGTH >132) has one argument: ?LENGTH >132.
FOO ([A,B]) has two arguments: [A and B].
FOO ( $\langle A, B\rangle$ ) has one argument: A,B.
Exception 2: If a left angle bracket is the first character of the argument list, or the first character after an unquoted comma, then it and its matching right angle bracket are treated as special quoting characters, but are not included in passed arguments.

FOO <A, B>, C has two arguments: $A, B$ and $C$.
FOO $C,\langle A, B\rangle$ has two arguments: $C$ and $A, B$.
You can alter this argument handling by using the pseudo-op. DIRECTIVE with MACMPD, . ITABM, and . XTABM. (See Chapter 3.)

## NOTE

```
To pass special characters in a macro
call, we suggest defining the macro so
that the delimiters are part of the
passed argument. For example, use
IIEFINE T1 (A) OOUTSTF [ASCTZ AT%
rather than
IEFINE TS (A) GOUTSTF FASCIZ \A\I%
The call Tl ">>" will work, but T2 ">>"
will not.
```


### 5.2.3 Listing of Called Macros

You can control the listing of called macros by using the pseudo-ops XALL, SALL, and LALL. LALL causes macro expansions to be listed in full; XALL suppresses part of the listing; LALL suppresses all of the listing. The default among these three is XALL.

The following example shows the action of these pseudo-ops on macro listings:

```
IIEFINE FOO (N)<
IFE N,2>
IFN N,1>
>
SALLL
    XAL.L
    FOO(0)m
    LAl..L
    FOO(0)m
    IFN O,<1>
    FOO(1)m
    IFE 1,2%
```

$000000 \quad 000002 \quad \mathrm{FOO}(0)$
$000000000001 \quad F 00(1)$
000000 000002 IFE 0,2
000000000001 IFN $1,1 \%$
000000000002 TFE 0.2 S
000000000001 IFN 1,41

### 5.3 NESTING MACRO DEFINITIONS

You can nest macro definitions. That is, you can define a macro within the body of another macro definition. Notice, however, that the nested macro is not defined to the assembler until the nesting macro is called.

Here is an example:
LEFINE FEFSON (A) 《
DEFINE CHILI (E) <
MEFINE GFANICHILI (C) < EXF A,ByP
$\because$

Until the DEFINE PERSON statement is assembled, calls to PERSON, CHILD, and GRANDCHILD are illegal. These macros are not yet defined to the assembler.

When the DEFINE PERSON statement is reached and assembled, PERSON can be called, but not CHILD or GRANDCHILD. The call PERSON 1 generates the text

PERSON 1 1
IIEFINE CHILI (B) <
LIEFINE GRANHCHILI (C) < EXF 1, $\mathrm{E}, \mathrm{C}$ :
$->$
thus defining CHILD to the assembler. The following call CHILD 2 generates the text

CHILD $2^{-}$
DEFINE GRANICHILI (C) 《
EXF 1,2,C
and GRANDCHILD is defined to the assembler. Finally, a call to GRANDCHILD 3 generates

GFANIICHILI $3^{m}$
EXF $1,2,3^{m}$
$000000 \quad 000001$
$000000 \quad 000002$
$000000 \quad 000003$
Notice the result of a subsequent call to CHILD 10. The text
CHILII $10^{-}$
DEFTNE GRANNCHILI (C) ©
EXF 1,10,C
is generated, and this definition replaces the old definition of GRANDCHILD; the definitions of PERSON and CHILD are not changed. After this, the call GRANDCHILD 3 generates

GRANHCHILII $3^{\prime \prime}$

| 000000 | 000001 |
| :--- | :--- |
| 000000 | 000010 |
| 000000 | 000003 |

NOTE
Using multiple angle brackets for a passed argument preserves the argument as one unit. For example passing the argument <<A,B,C>> to nested macros causes the outer macro to pass <A,B,C> as one argument; the first nested macro passes $A, B$, and $C$ as three arguments.

## USING MACROS

### 5.4 CONCATENATING ARGUMENTS

The apostrophe (') is the concatenation operator for macro calls. If you insert an apostrophe immediately before or after a dummy-argument in the body of a macro, the assembler removes it at call. This removal joins (concatenates) the passed argument to the neighboring character in the generated text.
(One application of this concatenation is shown under COMMON ERRORS for the ASCIZ pseudo-op.)

If the apostrophe precedes the dummy-argument, the passed argument is suffixed to the preceding character; if the apostrophe follows the dummy-argument, the passed argument is prefixed to the following character.

You can use more than one apostrophe with a dummy-argument. In this case only apostrophes next to the dummy-argument will be removed (at most one from each side). Other apostrophes are treated as regular characters in the macrobody. The following example shows the treatment of apostrophes on both sides of the dummy-argument, and of double apostrophes.

```
IUEFINE O (FREFIX,MITIFIX) &
    IIEFINE OCOMF (SUFFIX) <
                                    FREFIX'O'MILIFIX''SUFFIX%
```

Now the call $O$ A, J generates
IEFINE OCOMF (SUFFIX) 《 AOJ'SUFFIX:

```
because when the assembler replaces PREFIX with A, the apostrophe
following is removed to form AO. When J replaces MIDFIX, the
preceding apostrophe and first following apostrophe are removed to
form AOJ'SUFFIX.
Now the call OCOMP LE generates
                            OCOMF LEE
    34300000000000 AO.ILEN
since the apostrophe is removed to join AOJ to LE.
5.5 DEFAULT ARGUMENTS AND CREATED SYMBOLS
Ordinarily, an argument missing from a macro call is passed as nulls.
For example, the macro defined by
    DEFINE WORIIS (A,E,C) <
        EXFF A,E,C>
```

when called by WORDS 1,1 generates three words containing 1,1, and 0 ,
respectively.

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You can, however, alter this handing by specifying default values other than nulls, or by using created symbols.

### 5.5.1 Specifying Default Values

If you want a missing argument to default to some value other than nulls, you can specify the default value in your DEFINE statement. Do this by inserting the default value in angle brackets immediately after the dummy-argument. For example, the macro defined by

DEFINE WORDS (A,R222., C 333)
EXF A,B,C)
when called by WORDS $1, l$ generates three words containing 1,1, and 333, respectively.

## WORLIS 1,1m

| 000000 | 000001 |
| :--- | :--- |
| 000000 | 000001 |
| 000000 | 000333 |

NOTE
An argument passed as nulls by consecutive commas is not considered missing and cannot invoke a default value. Therefore missing arguments can occur only at the end of the list of passed arguments.

### 5.5.2 Created Symbols

A symbol used as a label in a macrobody must be different for each call of the macro (since duplicate labels are not allowed). Therefore for each call a different symbol for the label must be passed as an argument.

If you do not refer to such a label from outside the macro, you can simply let the assembler provide a new label for each call. This label is called a created symbol, and is of the form ..nnnn where nnnn is a 4-digit number.

To use a created symbol in place of a passed argument, use the percent sign (\%) as the first character of the dummy-argument in your DEFINE statement. The assembler then creates a symbol for use in the macro expansion if that argument is missing from a call to the macro. If you provide an argument in the call, the passed argument overrides the created symbol.

## USING MACROS

NOTES

1. A null argument (indicated by two adjacent delimiters) is not treated as missing.
2. Avoid using symbols of the form ..nnnn, since they could interfere with created symbols.
```
The following example shows a macro defined with a created symbol, the
macro called using the created symbol, and the macro called overriding
the created symbol:
    IUEFINE COMF'AR (TEST,SAVE,INIEX,%HERE) ;
    %HERE: MOVE SAVE,TEST
    SETZ INIIEX,
    CAME SAUE,TABLE(INIIEX)
    JNST %HEFE
>
COMFAF T1,T2,T3M
    ..OOO1: MOUE T2,T1
    SETZ T3,
    CAME T2,TABLE(T3)
    JFST . OOO1
COMFAR T1,T2,T4,HEFE1-
HEFEE1: MOUE T2,T1.
            GETZ TA
            CAME T2,TABLE(TA)
            JFST HEFEL
```


### 5.6 INDEFINITE REPETITION

The pseudo-ops IRP, IRPC, and STOPI give a convenient way to repeat all or part of a macro; you can change arguments on each repetition if you wish, and the number of repetitions can be computed at assembly time. You can use these three pseudo-ops only within the body of a macro definition.

To see how IRP works, assume the macro definition
LIEFINE DOEACH (A) \&
IRF A,CA>
The call DOEACH <ALPHA,BETA,GAMMA> produces the code

| 000200 | ALFHA=200 |
| :--- | :--- |
| 000300 | BETA $=300$ |
| 000400 | GAMMA $=400$ |
|  | DOEACH ALFFHA, BETA, GAMMA - |
|  |  |
| 000200 | ALFFHA IFF |
| 000300 | BETA |
| 000400 | GAMMA |

because each subargument passed to IRP generates one repetition of the code. Notice that the range of IRP must be enclosed in angle brackets.

NOTE
Using angle brackets in the call to
DOEACH is critical, since they make the
string ALPHA, BETA,GAMMA a single
argument for IRP. IRP then sees the
commas as delimiting subarguments.

IRPC is similar to IRP, but an argument passed to IRPC generates one repetition for each character of the argument.

STOPI ends the action of IRP or IRPC after assembly of the current expansion. You can use STOPI with a conditional assembly to calculate a stopping point during assembly. For example:


### 5.7 ALTERNATE INTERPRETATIONS OF CHARACTERS PASSED TO MACROS

The normal argument passed by a macro call is simply the string of characters given with the call. MACRO offers three alternate interpretations of the passed argument.

## USING MACROS

If you prefix a backslash ( $($ ) to an expression argument, the argument passed is the ASCII numeric character string giving the value of the expression.

If you prefix a backslash-apostrophe ( $\backslash$ ') to an expression argument, the argument passed is the string whose value is the SIXBIT string with the integer value of the expression.

If you prefix a backslash-quotemark (\") to an expression argument, the argument passed is the string whose value is the ASCII string with the integer value of the expression.

To show how these work, the following example defines a macro to print the argument passed. Then four different arguments are passed using the various argument interpretations.

LALL
DEFINE LOOKIE (ARG) <
REMAFiK The fassed argument is: AFG $>$
LOOKIE 60́․
FEMARK The fassed arsument is: 60 m
LOOKIE $\backslash 60^{\circ}$
FEMARK The fassed arsimment is: 60 m
LOOKIE \'60~
REMAFK The passed arsiument is: Fi
LOOKIE \"60"
REMAFK The fassed arsumerit is: 0 -
$Z=60$
LOOKIE Z"
REMAFK The fassed arsumerit is: $Z$ -
LOOKIE \Z~
FEMAFK The fassed arsument is: 60 "
LOOKIE \'Z"
REMARK The fassed arsument is: $F=$
LOOKTE \"Z"
FEMAFK The massed arsument is: 0 -
$635170425164 \quad Z Z=$ SIXAIT $^{\prime}$
LOOKIE ZŹ
FEMAFK The Fassed arssument is: ZZ m
LOOKTE: $\backslash Z Z^{\prime \prime}$
FEEMAFK The fassed argumerit is: 63517042w164 -
LOOKIE \'ZZ"
FEMARK The Fisssed arsumerit is: STXBTT -

```
ZZZ="ASCJI"
LOOKTE ZZZ"
FEMAFKK The possed arsumenit is: ZZZ "
LOOKIE \ZZZ"
FEMAFK The rassed arsument js: 20323416231%
LOOKIE \'ZZZ%
FEMAFK The passed arsument is: ASCTI ..
```

CHAPTER 6
ASSEMBLER OUTPUT

MACRO can generate three kinds of output files:

1. A program listing (.LST) file
2. A binary program (.REL) file
3. A UNIVERSAL (.UNV) file

### 6.1 THE PROGRAM LISTING FILE

MACRO outputs the program listing file to the device you specify, usually your terminal or a disk file. You can control the form of the program listing by using the pseudo-ops . DIRECTIVE FLBLST, . DIRECTIVE SFCOND, LIST, XLIST, LALL, XALL and SALL. (See Chapter 3.) All MACRO programs begin with the implicit pseudo-ops LIST and XALL.

The listing has a heading at the top of each page and subpage. The first line gives the program name, the assembler version, the time and date of assembly, and the page number. The second line gives the program filename (including extension), the date and time of creation, and an optional program subtitle:

Example:

| TIMER | MACRO | $\% 53(711)$ | 10:07 | 27-AFF--77 | FAGE 2 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| IMEF' | MAC | 27-A |  | 10:06 | MACDEF |

The listing has up to 55 lines per page. You can change this by using the $L$ switch; /nnL specifies nn lines per page. A formfeed (CTRL/L) in your program begins a new page and increments the page number. If the linecount exceeds lines-per-page before a formfeed is found, a subpage number is formed. For example, the subpages following page 6 are 6-1, 6-2, and so forth. A formfeed would begin page 7 .

The five columns in the program listing give:

1. The CREF line number (if the program was assembled with the CREF switch on).
2. The line sequence number (if the input file is sequenced).
3. The 6-digit octal address of the storage word, usually a sequential location assignment.

$$
\begin{aligned}
& 400066^{\prime} \\
& 400067^{\prime} \\
& 400070^{\prime}
\end{aligned}
$$

An apostrophe (') after the address shows that it is relocatable.

For a PHASE pseudo-op, the phased address is given.
For a BLOCK pseudo-op, only the address of the first word is given.

For a program with PSECTS, the 2 -digit PSECT number of the current PSECT immediately follows the address. For example,
$000100^{\prime} 02$
For a LOC or RELOC pseudo-op, only the address to which the location counter is set is given; the next word of code will be assembled at that address.
4. The assembled binary code (if any) in one of eight formats.

Fullword: all zeros with number sign (000000000000\#), showing that a fullword Polish fixup is required for the word of code.

Halfword: two 18-bit bytes. Each halfword can be followed by an apostrophe (') to indicate that it is relocatable, or by a pound sign (\#) to indicate that a Polish fixup is required for it. When you use the . HWFRMT pseudo-op, all code is listed in halfword format.

Instruction: 9-bit op-code; 4-bit accumulator code; 1-bit indirect code; 4-bit index; 18-bit address.

Input/output: 3-bit I/O code; 7-bit device code; 3-bit operand; l-bit indirect code; 4-bit index; 18-bit address.

Byte pointer: 6-bit byte position; 6-bit byte size; 1 unused bit; l-bit indirect code; 4-bit index; 18-bit address.

ASCII: five 7-bit bytes; one unused bit.
SIXBIT: six 6-bit bytes.

BYTE: binary representation of specified bytes. Bytes appear on the program listing only to the extent that available horizontal space permits. For example, 36 l-bit bytes cannot be represented as individual bytes on the listing. Any halfword byte containing an address can be flagged by an apostrophe (') or by a pound sign (\#). See the halfword format above.

OPDEF or assignment: one or two l8-bit bytes, as needed.

These examples show some code in each format:


An apostrophe (') shows the code as relocatable. The examples show relocatable values in the right half of some words. The left half can also be relocatable.

An asterisk (*) shows a symbol to be EXTERNAL or undefined.
A number sign (\#) shows that a Polish expression is required to resolve the value.
5. Source statements and comments.

If the assembler finds errors in a line of text, it suffixes one or more letters to the sequence number as error codes. These error codes are discussed in Chapter 8. A code is not repeated for multiple errors of the same type in a line.

At the end of the listing, the assembler gives the total number of errors, followed by break addresses. The program break is the largest relocatable address assembled, plus 1. The absolute break is the largest absolute address assembled. The high-segment break is the largest high-segment address assembled. For a program with PSECTs, the break for each PSECT is also given.

The listing gives CPU time in the form mm:ss.sss where mm is minutes and ss.sss is seconds. Core used is given in K's; one $K$ is 1024 words (2000 octal).

In the symbol table at the end of the listing, some symbols may have the following codes:

```
ent result of ENTRY pseudo-op
ext EXTERNAL symbol
int INTERNAL symbol
pol defined in terms of EXTERNAL symbols
sen suppressed result of ENTRY pseudo-op
sex suppressed EXTERNAL symbol
sin suppressed INTERNAL symbol
spd suppressed for debugger
udf undefined symbol
```

If you use the /C switch with MACRO, you can generate three additional tables in the program listing. The /C switch directs MACRO to generate the listing file in a format suitable for input to CREF, the cross-referencing program. This is a . CRF file rather than the usual . LST file.

After assembly, the . CRF file can be used as input to CREF, and the output is the cross-referenced . LST file. This file contains the program listing and symbol table as described above. In addition, it has a cross-referenced symbol table, a table of macros and OPDEFs, and, if you use the / O switch with CREF, a cross-referenced table of opcodes and pseudo-ops.

The cross-referenced symbol table lists each user-defined symbol (except macros, OPDEFs, and SYN symbols), and lists the sequence number of each line containing the symbol.

The table of macros and OPDEFs shows each reference to macros, OPDEFs, and SYN symbols.

The opcode table shows each reference to MACRO-defined opcodes and pseudo-ops, giving the sequence number of each line containing the opcode or pseudo-op.

```
6.2 THE BINARY PROGRAM FILE
MACRO outputs the binary program file to the device you specify,
usually a storage device. The default device is a disk. Most of the
file is the binary expansion of your program instructions. These
instructions are formatted into groups called REL Blocks; each block
is labeled so that LINK can recognize it. Details of this formatting
and labeling are discussed in the LINK Reference Manual.
A relocatable binary program file can be stored on any input/output
device. The output format is not related to either block types or
logical divisions of the device.
```


### 6.3 THE UNIVERSAL FILE

```
THE UNIVERSAL file is output only if the source file contains the UNIVERSAL pseudo-op. (See the discussions at UNIVERSAL in Chapter 3 and in Section 9.2.)
A UNIVERSAL file contains only symbols and definitions. These definitions are available to any program, and can be obtained by using the SEARCH pseudo-op.
```


## CHAPTER 7

## USING THE ASSEMBLER

To assemble a MACRO program, use one of the following:

1. The operating system command COMPILE. (See the Monitor Calls manual for details.)
2. The $\$ M A C R O$ card for the BATCH program. (See the GALAXY Batch Reference Manual.)
3. The MACRO command level.

To assemble a program in the command level of MACRO, type $R$ MACRO to the system. The system then runs MACRO, which responds with an asterisk (*) :

- F MACFO
* 

Then define files for MACRO by typing a command of the form
relfile,listfile=sourcefile,...,sourcefile
where:
relfile is a filespec for the binary program output file.
listfile is a filespec for the program listing output file.
each sourcefile is a filespec for a source program input file;
MACRO assembles source files in the order given.
The default device for each file is DSK:, but you can override this by prefixing devicecode: to any of the files. Default file extensions are. REL for relfile, . LST for listfile (.CRF if you use the /C switch), and .MAC for each sourcefile. You can override these by suffixing a file extension to any of the files.

You can specify a directory for any of these files by suffixing a project-programmer number (PPN) in square brackets.

You can set switches by suffixing /char or (char) to a file, where char is a switch code. Switch codes and their meanings are given in Table 7-l.


NOTE
Many programmers use the following commands to check assembly of short code sequences:
*,TTY: =TTY:
FASS2
This displays the assembled code line by line as you type it in.

Table 7-1
MACRO Switch Options

| Switch | Meaning |
| :---: | :---: |
| /A | Advance magnetic tape reel by one file. The /A switch must immediately follow the device to which the switch refers. |
| /B | Backspace magnetic tape reel by one file. The /B switch must immediately follow the device to which the switch refers. |
| /C | Produce listing file in a format acceptable as input to CREF. Unless the filename is given, CREF.CRF is assumed; if no file extension is given, .CRF is assumed; if no listing device is specified, DSK: is assumed. |
|  | The /C switch can be used only with the file specification for the program listing file; it must appear between the comma and the equal sign. |
| /E | List macro expansions (same as LALL pseudo-op). |
| /F | Output binary listing in multiformat (same as .MFRMT pseudo-op). |
| /G | Output binary listing in halfword format (same as .HWFRMT pseudo-op). |
| /H | Print HELP text (list of switches and explanations). |
| /L | Reinstate listing (same as LIST pseudo-op). |
| /M | List only the call and binary produced in a macro expansion (same as SALL pseudo-op). |
| /N | Suppress error printouts on the terminal. |
| 10 | End literal with CR-LF or right square bracket (same as MLOFF pseudo-op). |
| /P | Increase the size of the pushdown list. This switch can appear as many times as desired. The pushdown list is initially set to a size of 80 (decimal) locations; each /P increases the size by 80 (decimal). /P must appear on the left of the $=$. |
| /Q | Suppress $Q$ (questionable) warning errors on the listing. / Q must appear on the left of the $=$. |

(Continued on next page)

Table 7-1 (Cont.)
MACRO Switch Options

| Switch | Meaning |
| :---: | :---: |
| /S | Suppress listing (same as XLIST pseudo-op). |
| /T | Skip to the logical end of the magnetic tape. The $/ T$ switch must immediately follow the device to which the switch refers. |
| /U | Do not generate a .UNV file on DSK. The /U switch must appear immediately after the specification for the binary program file; that is, it must appear between the file specification and the comma. |
| /W | Rewind the magnetic tape. The /W switch must immediately follow the device to which the switch refers. |
| /X | Suppress listing of macro expansions (same as XALL). |

## CHAPTER 8 <br> ERRORS AND MESSAGES

MACRO has three kinds of messages:

1. Informational messages
2. Single-character error codes
3. MCRxxx messages (where $x x x$ is a 3-letter mnemonic code)

### 8.1 INFORMATIONAL MESSAGES

MACRO's informational messages are printed at the foot of the program listing. These messages and their explanations are given in Table 8-1.

Table 8-1
MACRO Informational Messages

| Message | Explanation |
| :---: | :---: |
| ABSLUTE BREAK | The highest absolute address over 137. |
| CORE USED | The size of the low segment used to assemble the source program. |
| CPU TIME USED | The CPU time for assembly in minutes and seconds. |
| ERRORS DETECTED | The number of errors detected by MACRO during assembly (errors marked on the listing by single-character codes other than Q). |
| HI-SEG. BREAK | The length of the high segment. |
| PROGRAM BREAK | The length of the low segment. |
| PSECT n BREAK | The length of PSECT n . |
| UNASSIGNED DEFINED AS IF EXTERNAL | Undefined symbol; treated as EXTERNAL. |
| WARNINGS GIVEN | The number of $Q$ errors found. Processing is terminated if under BATCH. |

### 8.2 SINGLE-CHARACTER ERROR CODES

Single-character error codes are printed in the program listing near the left margin of the line where the error occurs. If more than one kind of error occurs in the same line, more than one character will be printed; if more than one error of the same kind occurs in the line, the code is printed only once.

Codes for $M, P, V$, and $X$ errors are typed during Pass 1.
If you use CREF to produce a cross-referenced listing file, all the single-character error codes will appear in the cross-reference table as \%....x, where $x$ is the code character.

Table 8-2 gives the single-character error codes and their explanations.

Table 8-2
MACRO Single-Character Error Codes

| Code | Explanation |
| :---: | :---: |
| A | Argument error in pseudo-op. This is a broad class of errors that can be caused by an improper argument in a pseudo-op. The A errors include: |
|  | 1. Symbol used is improperly formed. |
|  | 2. IFIDN comparison string is too large. |
|  | 3. OPDEF of macro or SYN. |
|  | 4. Invalid SIXBIT character. |
|  | 5. Byte size in BYTE more than 36. |
|  | 6. RADIX50 code not absolute. |
|  | 7. End of line of IF reached before < character seen. |
|  | 8. Assignment made in an address field; for example, MOVEI A=10. (However, MOVEI <A=10> is valid.) |
|  | 9. Assignment of a label; for example, TAG: TAG=1. |
|  | 10. Missing symbol in SYN. |
|  | 11. Unknown symbol in SYN. |
|  | 12. Missing right parenthesis in an index. |
|  | 13. Missing left parenthesis in a BYTE statement. |
|  | 14. No comma after repeat count. |
|  | 15. IRP or IRPC not in a macro. |
|  | 16. Argument for IRP or IRPC is not a dummy symbol; for example, DEFINE GO (A) IRP B. |
|  | 17. IRP or IRPC argument is a created symbol. |
|  | 18. STOPI not in IRP or IRPC. |
| D | Multiply defined symbol. The statement contains a tag that refers to a multiply defined symbol. The first definition is used for assembling the statement. |

(Continued on next page)

## ERRORS AND MESSAGES

Table 8-2 (Cont.)
MACRO Single-Character Error Codes

| Code | Explanation |
| :---: | :---: |
| E | Improper use of an EXTERNAL symbol. The E errors include: |
|  | 1. Symbol both EXTERNAL and internal. |
|  | 2. EXTERNAL symbol used as accumulator register address. |
|  | 3. EXTERNAL symbol used with IF. |
|  | 4. EXTERNAL symbol used as address for LOC, RELOC, PHASE, HISEG, or TWOSEG. |
|  | 5. EXTERNAL symbol used for array name or size in ARRAY. |
|  | 6. EXTERNAL symbol used as REPEAT count. |
| L | Literal generates less than 1 or more than 99 words of data. |
| M | Symbol defined more than once; retains its first definition. If a symbol is first defined as a variable and later as a label, it retains the label definition. This error can be caused by multiple appearances of TITLE, or TITLE with UNIVERSAL. |
| N | Number error. The N errors include: |
|  | 1. Number exceeds the permitted range. |
|  | 2. B shift not absolute. |
|  | 3. Digits exceed current radix. If radix is 8, the single character 9 is acceptable but the number 19 is not acceptable. |
|  | 4. Character after up-arrow not B, O, F, L, D, !, or -. |
|  | 5. Illegal expression after E. |
| 0 | Operation code undefined. It is assembled as zeros. |
| P | Phase error. In general, the assembler generates the same number of program locations in Pass 1 and Pass 2. Any discrepancy causes a phase error. |
|  | Phase errors can be caused by incorrect literal allocation. |
|  | If a symbol is used as a macro to generate code in Pass 1 , and is used as a label in Pass 2, a phase error can occur. |
|  | A relocatable label that is defined in a literal and then used in an arbitrary expression; MACRO generates a Polish expression instead of treating the label as EXTERNAL. |

(Continued on next page)

Table 8-2 (Cont.)
MACRO Single-Character Error Codes


### 8.3 MCRxxx MESSAGES

The MCRxxx messages are issued to the terminal during assembly. (The xxx represents a 3-letter code.)

Any MCRxxx message that is preceded by a question mark is normally fatal under batch processing. A few MCRxxx messages are informational; these are issued within square brackets.

Table 8-3 gives all the MCRxxx messages. Each 3-letter code and its message are printed in boldface type. For some messages, an explanation is printed in lightface type.

Table 8-3
MCRxxx Messages

| Code | Message and Explanation |
| :--- | :--- |
| ATS | LINES/PAGE ARGUMENT TOO SMALL |
|  | The argument given must be greater than three to allow space |
|  | for the page heading. |

Table 8-3 (Cont.)
MCRxxx Messages

| Code | Message and Explanation |
| :---: | :---: |
| IBL | INPUT BLOCK TOO LARGE DEVICE |
|  | An input block from the specified device is too large. |
| ICP | INPUT CHECKSUM OR PARITY ERROR DEVICE |
|  | This is a hard-data error. |
| IDE | INPUT DATA ERROR DEVICE |
|  | This is a hard-data error. |
| ISC | ILLEGAL SYNTAX IN CONDITIONAL OR REPEAT |
| ISD | ILLEGAL SYNTAX IN MACRO DEFINITION |
|  | The macro is improperly defined. |
| IS I | ILLEGAL SYNTAX IN [IRP or IRPC] INSIDE MACRO |
| ISR | ILLEGAL SYNTAX IN REPEAT |
| LNF | LOAD THE NEXT FILE |
|  | The command string specifies the next file device as card reader or terminal. Input the file through the appropriate device. |
| LRE | (?) LOOKUP, RENAME, OR ENTER ERROR |
| LRE | (0) FILE WAS NOT FOUND |
| LRE | (1) NO DIRECTORY FOR PROJECT-PROGRAMMER NUMBER |
| LRE | (2) PROTECTION FAILURE |
| LRE | (3) FILE WAS BEING MODIFIED |
| LRE | (4) RENAME FILE NAME ALREADY EXISTS |
| LRE | (5) ILLEGAL SEQUENCE OF UUO'S |
| LRE | (6) BAD UFD OR BAD RIB |
| LRE | (7) NOT A SAV FILE |
| LRE | (10) NOT ENOUGH CORE |
| LRE | (11) DEVICE NOT AVAILABLE |
| LRE | (12) NO SUCH DEVICE |
| LRE | (13) NO TWO RELOC REG. CAPABILITY |

Table 8-3 (Cont.)
MCRxxx Messages

| Code | Message and Explanation |
| :---: | :---: |
| LRE | (14) NO ROOM OR QUOTA EXCEEDED |
| LRE | (15) WRITE LOCK ERROR |
| LRE | (16) NOT ENOUGH MONITOR TABLE SPACE |
| LRE | (17) PARTIAL ALLOCATION ONLY |
| LRE | (20) BLOCK NOT FREE ON ALLOCATION |
| LRE | (21) CAN'T SUPERSEDE (ENTER) AN EXISTING DIRECTORY |
| LRE | (22) CAN'T DELETE (RENAME) A NON-EMPTY DIRECTORY |
| LRE | (23) SFD NOT FOUND |
| LRE | (24) SEARCH LIST EMPTY |
| LRE | (25) SFD NESTED TOO DEEPLY |
| LRE | (26) NO-CREATE ON FOR SPECIFIED PATH |
| LTL | LITERAL TOO LONG |
| MDE | MONITOR DETECTED SOFTWARE INPUT ERROR DEVICE |
|  | The input file is not in a valid mode. |
| MPA | MISSING CLOSE PAREN AROUND ARG LIST |
| NEC | INSUFFICIENT CORE |
|  | Not enough memory is available to assemble the program. |
| NES | NO END StATEMENT ENCOUNTERED ON INPUT FILE |
| NUF | NOT A REAL UNIVERSAL FILE |
|  | No such UNIVERSAL file was found. |
| OBL | OUTPUT BLOCR TOO LARGE DEVICE |
|  | This is a file-status error. |
| OCP | OUTPUT CHECKSUM OR PARITY ERROR DEVICE |
|  | This is a hard-data error. |
| ODE | OUTPUT DATA ERROR DEVICE |
|  | This is a hard-data error. |
| OQE | OUTPUT QUOTA EXCEEDED ON DEVICE |
| OUF | UNIVERSAL FILE DEFAULT ARGUMENTS LOST, REASSEMBLE |

(Continued on next page)

```
Table 8-3 (Cont.)
    MCRxxx Messages
```

| Code | Message and Explanation |
| :---: | :---: |
| PDL | PDP OVERFLOW, TRY / P |
|  | See the /P switch in Table 7-1. |
| PET | INPUT PHYSICAL END OF TAPE DEVICE |
| PGE | PRGEND ERROR |
|  | See the PRGEND pseudo-op for proper use of PRGEND. |
| PTC | POLISH TOO COMPLEX |
|  | A Polish expression is too complex for MACRO to handle. Restructure or split the expression. |
| SOC | STATEMENT OUT OF ORDER . COMMON |
|  | The .COMMON pseudo-op must precede all statements that generate code, and all references to the COMMON block. |
| STO | SEARCH TABLE OVERFLOW, CANNOT SEARCH UNIVERSAL |
| TMU | TOO MANY UNIVERSALS |
|  | Too many UNIVERSAL files are being searched. The number permitted is an assembly parameter; it can be increased by reassembling MACRO. |
| UVS | UNIVERSAL VERSION SKEW, REASSEMBLE UNIVERSAL |
|  | The UNIVERSAL file was assembled with a later version of MACRO than you are using now. Reassemble the UNIVERSAL file. |
| UWU | UNABLE TO WRITE UNIVERSAL FILE |
| WLE | OUTPUT WRITE-LOCK ERROR DEVICE |

## PROGRAMMING CONSIDERATIONS

The previous chapters of this manual define the MACRO language elements. In particular, the pseudo-op definitions in Chapter 3 define many of MACRO's most important features. However, the usefulness of some pseudo-ops can be seen only in the context of a "family" of pseudo-ops.

In this chapter, we discuss three such families of pseudo-ops. The programming features concerned are:

1. Program segmentation
2. UNIVERSAL files
3. Conditional assembly

### 9.1 PROGRAM SEGMENTATION

MACRO's relocation counters can accommodate three types of programs:

1. A single-segment program uses only one relocation counter.
2. A two-segment program also uses one relocation counter, and is characterized by its use of the TWOSEG pseudo-op.
3. A program with PSECTS can use many relocation counters, and is characterized by its use of the .PSECT and .ENDPS pseudo-ops.

### 9.1.1 Single-Segment Programs

A single-segment program uses only one relocation counter. This counter can be used to assign any address from 0 to 777777. The initial setting of the counter is 0 .

As MACRO assembles your program, it places code and data at the address given by the current value of the relocation counter, incrementing the counter's value for each word assembled.

For example, a statement can require assembly of one word of code, incrementing the relocation counter by l. Another statement can require assembly of five words of code, incrementing the relocation counter by 5. Still another statement may not generate code, leaving the relocation counter unchanged.

You can reset the value of the relocation counter by using the pseudo-op RELOC with an argument. For example, using RELOC A sets the value of the relocation counter to the value of $A$.

In the following example, 100 words are allocated for a table, incrementing the relocation counter by 100. Then the table length is calculated as TABLEN. A RELOC TABLE returns to the top of the table, where the first three words are initialized. Finally a RELOC TABLE+TABLEN sets the relocation to the foot of the table to continue assembly.

| $000000^{\circ}$ |  | 000100 | TABLE: | $\begin{aligned} & \text { ELOCK } 100 \\ & \text { TABLEN }=, \text { TABLE } \end{aligned}$ | ;Allocate table <br> prable lensth |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 000000 |  |  |  | FELOC TABLE | ; Tof of table |
| 000000' | 000000 | 000001 |  | EXF 1,2,3 | ; Irit first 3 |
| 000001. | 000000 | 000002 |  |  |  |
| 000002' | 000000 | 000003 |  |  |  |
| $00010{ }^{\prime}$ |  |  |  | FELOC TAELE+TA | ; Contirise |

### 9.1.2 Two-Segment Programs

By using the TWOSEG pseudo-op, you can divide your program into a high segment and a low segment. This pseudo-op must precede any statement in your program that generates code.

The TWOSEG pseudo-op tells MACRO that there will be two segments, and MACRO generates a REL Block Type 3 , which tells LINK to expect two segments for loading.

You can use TWOSEG either with or without an address argument. There are important differences between the two:

1. TWOSEG without an argument specifies that the high segment begins at the address 400000. The initial value of the relocation counter is at the address 0 in the low segment.
2. TWOSEG with an argument specifies that the high segment begins at the given address, and further specifies that the initial value of the relocation counter is that address. (The given address is reduced to the next lower multiple of 2000 octal; if this result is 0 , MACRO treats the TWOSEG as if no argument were given.)

The high-segment starting address divides all code into two segments. MACRO and LINK consider all code at addresses above the high-segment address to be in the high segment, and all other code to be in the low segment.

MACRO always remembers the value the relocation counter had before the last RELOC found. (This stored value is initially 0.)

Therefore in a two-segment program, you can begin in one segment, and then RELOC to the other. From then on, you can switch segments simply by using RELOC with no argument. MACRO will begin assigning addresses at the first unused location in the opposite segment.

## PROGRAMMING CONSIDERATIONS

For example,

| $400000^{\prime}$ | TwOSEG |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 000000' | 000000 | 000001 | EXF 1,2 | 9towses |
| 000001. | 000000 | 000002 |  |  |
| $400000^{\prime}$ |  |  | FELOC 400000 | 9Hiwses |
| $400000^{\prime}$ | 000000 | 000003 | EXF 3.4 |  |
| $400001^{\prime}$ | 000000 | 000004 |  |  |
| 000002' |  |  | FELOC | Ftowses |
| 000002 | 000000 | 000005 | EXF 5.6 |  |
| 0000031 | 000000 | 000006 |  |  |
| 400002' |  |  | FELOC | \%Hi-ses |
| 4000021 | 000000 | 000007 | EXF 7.10 |  |
| 4000031 | 000000 | 000010 |  |  |

### 9.1.3 Programs With PSECTs

You can construct a program having up to 64 segments by using the . PSECT and .ENDPS pseudo-ops. These pseudo-ops control switching among program segments (PSECTs).

Each PSECT has its own relocation counter; each is separately relocated at load time. Therefore a program with two PSECTS is different from a two-segment program in that the PSECTed program has two relocation counters, while the two-segment program has only one.

The pseudo-op . PSECT specifies that code should be assembled for a given PSECT. For example, . PSECT A specifies that code is to be assembled in the program segment (PSECT) called A. The pseudo-op .ENDPS ends assembly in the current PSECT.

PSECTs can be nested up to 16 levels. In a nested PSECT, the .ENDPS pseudo-op begins assembly in the next outer PSECT; in an unnested PSECT, . ENDPS begins assembly in the blank PSECT. (You can think of the blank PSECT as being outside of all your explicitly declared PSECTS.)

Here is an example showing three $\operatorname{PSECTs}(A, B$, and $C$ ):

| $000000^{\prime} 00$ | 000000 | 000001 | EXF 1,2 | ;Blank FGECT |
| :---: | :---: | :---: | :---: | :---: |
| $000001^{\prime} 00$ | 000000 | 000002 |  |  |
| $000000^{\prime} 01$ |  |  | +FSECT A | \% 5 EL FSECT |
| 000000'01 | 000000 | 000003 | EXF 3,4 |  |
| 000001'01 | 000000 | 000004 |  |  |
| $000000^{\prime} 02$ |  |  | . FSECT B | ;2mid FSECT (nested) |
| 000000'02 | 000000 | 000005 | EXF S.6 |  |
| 000001'02 | 000000 | 000006 |  |  |
| $000002^{\prime} 01$ |  |  | - ENIFS B | ; 1 st FSECT |
| $000002^{\prime} 01$ | 000000 | 000007 | EXF 7,10 |  |
| $000003^{\prime} 01$ | 000000 | 000010 |  |  |
| 000002.00 |  |  | -ENHFS A | ; Elark FGECT |
| 000002100 | 000000 | 000011 | EXF 11,12 |  |
| 000003.00 | 000000 | 000012 |  |  |
| $000000^{\prime} 03$ |  |  | -FGECT C | ;3rd FSECT |
| $000000^{\prime} 03$ | 000000 | 000013 | EXF 13,14 |  |
| $000001^{\prime} 03$ | 000000 | 000014 |  |  |
| 000004'00 |  |  | - ENTIFS C | ¢ Elark FPSECT |
| 000002'02 |  |  | -FSECT B | - 2nod FSECT |
| 000002'02 | 000000 | 000015 | EXF 15,16 |  |
| $000003^{\prime 2}$ | 000000 | 000016 |  |  |
| 000004'00 |  |  | - ENIIFS E | ABlark. FSECT |

## PROGRAMMING CONSIDERATIONS

In the example, the blank PSECT surrounds everything. Embedded in the blank PSECT are:

1. PSECT A (which also nests some of PSECT B)
2. PSECT C
3. Another segment of PSECT B

Each PSECT used in a program generates the PSECT name as a global symbol. At load time, this symbol will take the value of the origin specified for the PSECT.

When LINK loads your progran, all the parts of the same PSECT are loaded together. These parts can be in more than one program, or in more than one file. For details of LINK's handling of PSECTs at load time, see the LINK Reference Manual.

### 9.2 UNIVERSAL FILES

A UNIVERSAL file contains direct-assignment symbol definitions. The symbols defined can have any attributes.

A UNIVERSAL file is convenient because it can contain definitions that you want for many programs. Those programs can then obtain the definitions by your use of the SEARCH pseudo-op. This searching adds to the assembly only those definitions that are needed; other definitions in the UNIVERSAL file are not used.

To build a UNIVERSAL file from a MACRO source file, insert the pseudo-op

UNIVERSAL filespec
where the filespec gives the file for output of the UNIVERSAL file. This file will contain all the symbols and definitions given in the program.

Another program can obtain these definitions if it contains the SEARCH pseudo-op:

## SEARCH filespec

where filespec names the UNIVERSAL file. At the end of Pass l assembly, MACRO will search the UNIVERSAL file for any undefined symbols. If a definition is found in the UNIVERSAL file, MACRO moves it into the symbol tables of the current program.

For example, a UNIVERSAL file can contain definitions for register mnemonics:

UNIUEFGAL FEEGS

| 000000 | $\mathrm{FO}=0$ |
| :--- | :--- |
| 000001 | $\mathrm{~F} 1=1$ |
| 000002 | $\mathrm{~F} 2=2$ |
| 000003 | $\mathrm{~F} 3=3$ |
| 000004 | $\mathrm{~T} 1=4$ |
| 000005 | $\mathrm{~T}=5$ |
| 000016 | $\mathrm{SF}=16$ |
| 000017 | $\mathrm{~F}=17$ |
|  | ENII |

## PROGRAMMING CONSIDERATIONS

Then another assembly can obtain these by using the SEARCH REGS pseudo-op:

SEARCH REGS

| $00000{ }^{\prime}$ | 000 | 00 | 0 | 00 | 000000 | Z | FO, |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 000001' | 000 | 01 | O | 00 | 000000 | Z | Fi, |
| 000002. | 000 | 02 | - | 00 | 000000 | Z | F2, |
| 0000031 | 000 | 03 | 0 | 00 | 000000 | Z | F3, |
| $000004^{\prime}$ | 000 | 04 | 0 | 00 | 000000 | Z | T1. |
| 000005' | 000 | 05 | 0 | 00 | 000000 | 2 | T2, |
| 000006 | 000 | 1.6 | 0 | 00 | 000000 | $z$. | SF, |
| $00000{ }^{\prime}$ | 000 | 17 | 0 | 00 | 000000 | Z. | F', |

A UNIVERSAL file can contain definitions for any user-defined symbols. You may find it convenient to build UNIVERSAL files containing macros, OPDEFs, and direct-assignment symbols that you use often in your programs.

An example of a UNIVERSAL program appears in the program examples in Appendix D.

### 9.3 CONDITIONAL ASSEMBLY

Using conditional assembly in your programs can make programming easier, and can make your assembled programs shorter. The pseudo-ops used for conditional assembly are IRP, IRPC, STOPI, .IF, .IFN, and the IFx group. IRP, IRPC, and STOPI are discussed fully in Chapter 3 and Section 5.6.

We will confine the discussion here to a few classic uses of the remaining conditional assembly pseudo-ops.

The first of these is the use of IFNDEF to establish default switch settings for a program. The example here is from the MACRO program itself, and concerns assembly of F 40 -switch-dependent symbols.

Near the beginning of the code, MACRO has the statement:
IFNDEF F40<F40==0>
This statement has effect only if the symbol $F 40$ is not defined, in which case the statement $F 40==0$ is assembled. This sets the $F 40$ switch to "off."

But if a file defining F 40 is assembled with (and before) the MACRO source file, then the statement $F 40==0$ is not assembled, leaving the "outside" definition in force.

Therefore the statement IFNDEF F40 <F40==0> serves as a default definition for $F 40$, and this default is used only if no other definition overrides it.

Another application of conditional assembly is connected with the symbol F40. In MACRO's program segments on symbol searching, some symbols will be defined (and therefore found in the search) only if the F40 switch is "on."

Here is how MACRO's code handles these symbols. There is a code sequence as follows:
; MACRO TO HANDLEE F4O UUOS
IFE F40,
DEFINE XF (SB,CD)
iNULL MACFO
IFN F4O,GYN X,XF: ;USUAL X MACRO
The "usual $x$ macro" is merely a macro to set up symbols to be defined and the code to assemble on finding them. The macro XF will be used to handle definitions for F 40 UUOs.

Now if the F 40 switch is on, the macro XF is made synonymous with the macro $X$, and the $F 40$ UUOs are defined in the same way as other operators. But if the F 40 switch is off, XF is made a null macro so that all the F 40 UUOs are ignored during assembly and are not defined to MACRO.

The assembly of the 540 UUOs depends on the value of the 440 switch, and the value of the switch depends on its definition. If MACRO had no IFNDEF F40 statement, an "outside" file would have to define the switch at every assembly of MACRO. But the default definition allows assembly of MACRO alone, and the outside file is needed only to turn the switch on.

Examples of conditional assembly are shown in the program examples in Appendix D.

APPENDIX A

## MACRO CHARACTER SETS

Table A-l gives the 101 ASCII characters allowed in MACRO and their octal ASCII codes; the 64 SIXBIT characters and their octal SIXBIT codes; and the 40 RADIX50 characters and their octal RADIX50 codes.

| Character | $\begin{aligned} & \text { ASCII } \\ & \text { Code } \end{aligned}$ | SIXBIT <br> Code | $\begin{gathered} \text { RADIX50 } \\ \text { Code } \end{gathered}$ |
| :---: | :---: | :---: | :---: |
| (horizontal tab) | 011 |  |  |
| (linefeed) | 012 |  |  |
| (vertical tab) | 013 |  |  |
| (formfeed) | 014 |  |  |
| (carriage-return) | 015 |  |  |
| ( CTRL/Z) | 032 |  |  |
| (CTRL/-) | 037 |  |  |
| (blank) | 040 | 00 | 00 |
| ! | 041 | 01 |  |
| " | 042 | 02 |  |
| \# | 043 | 03 |  |
| \$ | 044 | 04 | 46 |
| \% | 045 | 05 | 47 |
| \& | 046 | 06 |  |
| , | 047 | 07 |  |
| ( | 050 | 10 |  |
| ) | 051 | 11 |  |
| * | 052 | 12 |  |
| + | 053 | 13 |  |
| , | 054 | 14 |  |
| - | 055 | 15 |  |
|  | 056 | 16 | 45 |
| 1 | 057 | 17 |  |

(Continued on next page)

## MACRO CHARACTER SETS

| Character | ASC I I Code | SIXBIT Code | RADIX50 Code |
| :---: | :---: | :---: | :---: |
| 0 | 060 | 20 | 01 |
| 1 | 061 | 21 | 02 |
| 2 | 062 | 22 | 03 |
| 3 | 063 | 23 | 04 |
| 4 | 064 | 24 | 05 |
| 5 | 065 | 25 | 06 |
| 6 | 066 | 26 | 07 |
| 7 | 067 | 27 | 10 |
| 8 | 070 | 30 | 11 |
| 9 | 071 | 31 | 12 |
| : | 072 | 32 |  |
| ; | 073 | 33 |  |
| $<$ | 074 | 34 |  |
| = | 075 | 35 |  |
| > | 076 | 36 |  |
| ? | 077 | 37 |  |
| @ | 100 | 40 |  |
| A | 101 | 41 | 13 |
| B | 102 | 42 | 14 |
| C | 103 | 43 | 15 |
| D | 104 | 44 | 16 |
| E | 105 | 45 | 17 |
| F | 106 | 46 | 20 |
| G | 107 | 47 | 21 |
| H | 110 | 50 | 22 |
| I | 111 | 51 | 23 |
| J | 112 | 52 | 24 |
| K | 113 | 53 | 25 |
| L | 114 | 54 | 26 |
| M | 115 | 55 | 27 |
| N | 116 | 56 | 30 |
| 0 | 117 | 57 | 31 |
| P | 120 | 60 | 32 |
| Q | 121 | 61 | 33 |
| R | 122 | 62 | 34 |
| S | 123 | 63 | 35 |
| T | 124 | 64 | 36 |
| U | 125 | 65 | 37 |
| V | 126 | 66 | 40 |
| W | 127 | 67 | 41 |

(Continued on next page)

## MACRO CHARACTER SETS

Table A-1 (Cont.)
MACRO Character Sets

| Character | ASCII Code | SIXBIT Code | RADIX50 Code |
| :---: | :---: | :---: | :---: |
| X | 130 | 70 | 42 |
| Y | 131 | 71 | 43 |
| Z | 132 | 72 | 44 |
| [ | 133 | 73 |  |
| 1 | 134 | 74 |  |
| 1 | 135 | 75 |  |
|  | 136 | 76 |  |
| - | 137 | 77 |  |
| a | 141 |  |  |
| b | 142 |  |  |
| c | 143 |  |  |
| d | 144 |  |  |
| e | 145 |  |  |
| f | 146 |  |  |
| 9 | 147 |  |  |
| h | 150 |  |  |
| i | 151 |  |  |
| j | 152 |  |  |
| k | 153 |  |  |
| 1 | 154 |  |  |
| m | 155 |  |  |
| n | 156 |  |  |
| - | 157 |  |  |
| p | 160 |  |  |
| q | 161 |  |  |
| r | 162 |  |  |
| s | 163 |  |  |
| t | 164 |  |  |
| u | 165 |  |  |
| v | 166 |  |  |
| w | 167 |  |  |
| x | 170 |  |  |
| y | 171 |  |  |
| z | 172 |  |  |

A-3

## MACRO SPECIAL CHARACTERS

APPENDIX B

```
Characters and combinations having special interpretations in MACRO
are given in Table B-l. These interpretations apply only in the
contexts described. In particular, they do not apply within text
strings or comment fields.
For each usage of special characters, a cross-reference to a text
discussion is given in the rightmost column of the table. For
references to pseudo-ops, only the pseudo-op name is given; all
pseudo-ops are discussed in alphabetical order in Chapter 3.
```

Table B-1
Interpretations of Special Characters

| Characters | Context | Form | Interpretation | Discussed in Section |
| :---: | :---: | :---: | :---: | :---: |
| B | between two integer expressions | mBn | causes the binary representation of $m$ to be placed with rightmost bit at bit $n$ (decimal). | 2.2 .6 |
| ${ }^{\wedge} \mathrm{B}$ | before integer expression | ${ }^{\wedge} \mathrm{Bn}$ | shows that $n$ is a binary number. | 2.2 .2 |
| * D | before integer expression | ${ }^{\text {a }} \mathrm{Dn}$ | shows that n is a decimal number. | 2.2 .2 |
| E | between floatingpoint decimal number and signed decimal integer | $\mathrm{fE}+\mathrm{n}$ | multiplies $f$ by the +nth power of 10 . | 2.2 .5 |
| ${ }^{\wedge} \mathrm{F}$ | before integer expression | ${ }^{\wedge} \mathrm{Fn}$ | shows that $n$ is a fixedpoint decimal number. | 2.2.4 |
| G | after integer | nG | suffixes nine zeros to $n$. | 2.2.3 |
| K | after integer | nK | suffixes three zeros to $n$. | 2.2 .3 |
| * L | before decimal <br> integer <br> expression | ${ }^{\wedge} \mathrm{Ln}$ | generates the number of leading zeros in the binary representation of $n$. | 2.2.8 |
| M | after integer | nM | suffixes six zeros to n . | 2.2 .3 |

Table B-l (Cont.)
Interpretations of Special Characters

| Characters | Context | Form | Interpretation | Discussed in Section |
| :---: | :---: | :---: | :---: | :---: |
| ${ }^{\circ} \mathrm{O}$ | before integer expression | ${ }^{\circ} \mathrm{On}$ | shows that n is an octal number. | 2.2 .2 |
| : | after symbol | sym: | shows that sym is a label. | 2.4.2.1, 4.1, 4.5 |
| : | after symbol | sym: | shows that sym is a global INTERNAL label. | 2.4.2.1, 4.1, 4.5 |
| : | after symbol | sym: | shows that sym is a label, but not to be output by debugger. | 2.4.2.1, 4.1, 4.5 |
| : : | after symbol | sym: : | shows that sym is a qlobal INTERNAL label, but not to be output by debuqger. | 2.4.2.1, 4.1, 4.5 |
| ; | before end of line | ; text | shows that text is a comment. | 4.4, 4.5 |
| ; | before end of line (usually in a macro) | ; text | shows that text is a comment to be printed in the macro definition but not at call. | 4.4, 4.5 |
| - | as expression |  | generates current value of the location counter. | $2.3,4.6$ |
| - | embedded in numerals | int.fr | shows that int.fr is a floating-point decimal number. | 2.2 .5 |

Table B-1 (Cont.)
Interpretations of Special Characters

| Characters | Context | Form | Interpretation | Discussed in Section |
| :---: | :---: | :---: | :---: | :---: |
| , | among numbers and symbols | , | delimits operands, accumulator, arguments. | $\begin{array}{ll} 4.3, & 4.5 \\ 5.1, & 5.2 \end{array}$ |
| '' | among numbers and symbols | ', | delimits a null macro argument. | 5.2, 5.5 |
| '' | between two expressions | lhw, ,rhw | delimits left halfword (lhw) from right halfword (rhw). | 2.5.4.1 |
| $!$ | between two expressions | A! B | generates the logical <br> inclusive $O R$ of $A$ and $B$. | 2.5 .2 |
| ~! | between two expressions | $A^{\wedge}!B$ | generates the logical exclusive $O R$ of $A$ and $B$. | 2.5.2 |
| \& | between two expressions | $A \& B$ | generates the logical AND of $A$ and $B$. | 2.5 .2 |
| - | before expression | $\bigcirc-\mathrm{A}$ | generates one's complement of value of A (logical NOT). | 2.5.2 |
| * | between two expressions | $A * B$ | generates product of $A$ and $B$. | 2.5.1 |
| 1 | between two expressions | A/B | generates quotient of $A$ by $B$. | 2.5 .1 |
| + | between two expressions | $A+B$ | generates sum of $A$ and $B$. | 2.5 .1 |

MACRO SPECIAL CHARACTERS

| Table B-1 (Cont.) <br> Interpretations of Special Characters |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| Characters | Context | Form | Interpretation | Discussed in Section |
| - | between two expressions | $A-B$ | generates difference of $A$ and $B$. | 2.5 .1 |
| - | before an expression | -A | generates the two's complement of the value of $A$. | $\begin{aligned} & 2.2 .1,2.2 .4, \\ & 2.2 .5 \end{aligned}$ |
| "..." | around text | "text" | shows that text is a 7-bit ASCII string, to be right justified in field of five characters. | ASCII, ASCIZ |
| '...' | around text | 'text' | shows that text is a SIXBIT string, to be right justified in field of six characters. | SIXBIT |
| ' | adjoining dummy argument in macro body | $\begin{gathered} \text { text'darg } \\ \text { or } \\ \text { darg'text } \end{gathered}$ | concatenates passed argument to text at call to macro. | 5.4 |
| \# | after symbol | sym\# | shows that sym is a variable symbol, whose address is usually at the end of the binary program. | 2.4 .3 |
| \#\# | after symbol | sym\#\# | shows that sym is a global EXTERNAL symbol. | 2.4.5.2 |
| 1 | prefixed to expression in macro call | $\backslash \mathrm{expr}$ | directs that the argument passed be the string for the ASCII value of expr in the current radix. | 5.7.1 |

Table B-1 (Cont.)
Interpretations of Special Characters

| Characters | Context | Form | Interpretation | Discussed in Section |
| :---: | :---: | :---: | :---: | :---: |
| '' | prefixed to expression in macro call | \'expr | directs that the argument passed be the string whose SIXBIT code is the value of expr. | 5.7 .3 |
| \" | prefixed to expression in macro call | \"expr | directs that the argument passed be the string whose ASCII code is the value of expr. | 5.7.3 |
| CTRL/-(CONTROLunderscore) | before CR-LF | CTRL/- | continues argument to next line; does not operate across end-of-macro. | 5.2.1 |
| - | between two expressions | A_B | shifts the binary representation of $A$ to the left $B$ positions. (If $B$ is negative, shift is to right.) | 2.2 .6 |
| @ | prefixed to address | @address | sets bit 13 of the instruction word, indicating indirect addressing. | 4.7.1 |
| \% | lst character of dummy argument in macro definition | \%darg | directs that odarg be replaced by a created symbol at macro call; MACRO will substitute a different symbol for it on each use of the macro. | 5.5.2 |
| ( ) |  | (...) | encloses index field; encloses dummy arguments in macro definition; quotes characters for macro argument handling; swaps the two halves of enclosed value. | $\begin{aligned} & 4.7 .1,5.1 \\ & 5.2 .2 \end{aligned}$ |

## MACRO SPECIAL CHARACTERS

Table B-1 (Cont.)
Interpretations of Special Characters

| Characters | Context | Form | Interpretation | Discussed in Section |
| :---: | :---: | :---: | :---: | :---: |
| $<>$ |  | <...> | nests expressions; encloses conditional assembly code; encloses code in REPEAT, IRP, and IRPC pseudo-ops; encloses macrobody in DEFINE pseudo-op; quotes characters for macro argument handling; forces evaluation of symbol. | $2.5 .4$ <br> IFx, .IF, .IFN, REPEAT, IRP, IRPC, DEFINE, $5.1,5.2 .2$ |
| [ ] |  | [...] | ```delimits literals; delimits argument in ARRAY, .COMMON, and OPDEF pseudo-ops; quotes characters for macro argument handling.``` | 2.3, ARRAY, .COMMON, OPDEF, 5.2.2 |
| = | between symbol and expression | sym=exp | assigns value of exp to sym. | 2.4.2.2, 4.5 |
| = | between symbol and expression | sym $=$ exp | assigns value of exp to sym but sym is not output by debugger. | 2.4.2.2, 4.5 |
| = | between symbol and expression | $s y m=: \exp$ | assigns value of exp to sym and declares sym as global INTERNAL. | 2.4.2.2, 4.5 |
| = = | between symbol and expression | sym= $=$ : exp | assigns value of exp to sym and declares sym as global INTERNAL, but sym is not output by debuqqer. | 2.4.2.2, 4.5 |

## APPENDIX C

## MACRO-DEFINED MNEMONICS

This appendix contains tables showing all of MACRO's defined mnemonics and the code they generate. These mnemonics, together with the pseudo-ops and the special characters given in Appendix $B$, make up the entire MACRO language.

NOTE
Throughout this appendix, the following notes apply to the tables:

* Indicates mnemonic defined only if MACRO is assembled with the KLlO switch on.
** Indicates mnemonic defined only if MACRO is assembled with the KIlO switch on.
*** Indicates mnemonic defined only if MACRO is assembled with the F 40 switch on.


## C.l MACHINE INSTRUCTION MNEMONICS

Table C-l shows MACRO's machine instruction mnemonics and the code assembled by each mnemonic. See Section 4.7 for a discussion of machine instructions used in programs.

## MACRO-DEFINED MNEMONICS

Table C-1
Machine Instruction Mnemonics

| 270 | 00 | 0 | 00 | 000000 | ADD | 303 | 00 | 0 | 00 | 000000 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 273 | 00 | 0 | 00 | 000000 | ADDB | 306 | 00 | 0 | 00 | 000000 |
| 271 | 00 | 0 | 00 | 000000 | ADDI | 310 | 00 | 0 | 00 | 000000 | CAILE

(Continued on Next Page)

Table C-l (Cont.)
Machine Instruction Mnemonics

(Continued on Next Page)

## MACRO-DEFINED MNEMONICS

## Table C-1 (Cont.) <br> Machine Instruction Mnemonics

| 435 | 00 | 000 | 000000 | IORI | 471 | 00 | 000 | 000000 | ORCBI |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 436 | 00 | 000 | 000000 | IORM | 472 | 00 | 000 | 000000 | ORCBM |
| 255 | 00 | 000 | 000000 | JFCL | 464 | 00 | 000 | 000000 | ORCM |
| 243 | 00 | 000 | 000000 | JFFO | 467 | 00 | 000 | 000000 | ORCMB |
| 267 | 00 | 000 | 000000 | JRA | 465 | 00 | 000 | 000000 | ORCMI |
| 254 | 00 | 000 | 000000 | JRST | 466 | 00 | 000 | 000000 | ORCMM |
| 266 | 00 | 000 | 000000 | JSA | 435 | 00 | 000 | 000000 | ORI |
| 265 | 00 | 000 | 000000 | JSP | 436 | 00 | 000 | 000000 | ORM |
| 264 | 00 | 000 | 000000 | JSR | 262 | 00 | 000 | 000000 | POP |
| 104 | 00 | 000 | 000000 | JSYS | 263 | 00 | 000 | 000000 | POPJ |
| 320 | 00 | $0 \quad 00$ | 000000 | JUMP | 261 | 00 | 000 | 000000 | PUSH |
| 324 | 00 | 000 | 000000 | JUMPA | 260 | 00 | 000 | 000000 | PUSHJ |
| 322 | 00 | 000 | 000000 | JUMPE | 241 | 00 | 000 | 000000 | ROT |
| 327 | 00 | 000 | 000000 | JUMPG | 245 | 00 | 000 | 000000 | ROTC |
| 325 | 00 | 000 | 000000 | JUMPGE | 424 | 00 | 000 | 000000 | SETA |
| 321 | 00 | 000 | 000000 | JUMPL | 427 | 00 | 000 | 000000 | SETAB |
| 323 | 00 | 000 | 000000 | JUMPLE | 425 | 00 | 000 | 000000 | SETAI |
| 326 | 00 | 000 | 000000 | JUMPN | 426 | 00 | 000 | 000000 | SETAM |
| 135 | 00 | 000 | 000000 | LDB | 450 | 00 | 000 | 000000 | SETCA |
| 242 | 00 | 000 | 000000 | LSH | 453 | 00 | 000 | 000000 | SETCAB |
| 246 | 00 | 000 | 000000 | LSHC | 451 | 00 | 000 | 000000 | SETCAI |
| 257 | 00 | 000 | 000000 | **MAP | 452 | 00 | 000 | 000000 | SETCAM |
| 200 | 00 | 000 | 000000 | MOVE | 460 | 00 | 000 | 000000 | SETCM |
| 201 | 00 | 000 | 000000 | MOVEI | 463 | 00 | 000 | 000000 | SETCMB |
| 202 | 00 | 000 | 000000 | MOVEM | 461 | 00 | 000 | 000000 | SETCMI |
| 203 | 00 | 000 | 000000 | moves | 462 | 00 | 000 | 000000 | SETCMM |
| 214 | 00 | 000 | 000000 | MOVM | 414 | 00 | 000 | 000000 | SETM |
| 215 | 00 | 000 | 000000 | MOVMI | 417 | 00 | 000 | 000000 | SETMB |
| 216 | 00 | 000 | 000000 | MOVMM | 415 | 00 | 000 | 000000 | SETMI |
| 217 | 00 | 000 | 000000 | MOVMS | 416 | 00 | 000 | 000000 | SETMM |
| 210 | 00 | 000 | 000000 | MOVN | 474 | 00 | 000 | 000000 | SETO |
| 211 | 00 | 000 | 000000 | MOVNI | 477 | 00 | 000 | 000000 | SETOB |
| 212 | 00 | 000 | 000000 | MOVNM | 475 | 00 | 000 | 000000 | SETOI |
| 213 | 00 | 000 | 000000 | mOVNS | 476 | 00 | 000 | 000000 | SETOM |
| 204 | 00 | 000 | 000000 | MOVS | 400 | 00 | 000 | 000000 | SETZ |
| 205 | 00 | 000 | 000000 | MOVS I | 403 | 00 | 000 | 000000 | SETZB |
| 206 | 00 | 000 | 000000 | MOVSM | 401 | 00 | 000 | 000000 | SETZI |
| 207 | 00 | 000 | 000000 | mOVSS | 402 | 00 | 000 | 000000 | SETZM |
| 224 | 00 | 000 | 000000 | MUL | 330 | 00 | 000 | 000000 | SKIP |
| 227 | 00 | 000 | 000000 | MULB | 334 | 00 | 000 | 000000 | SKIPA |
| 225 | 00 | 000 | 000000 | MULI | 332 | 00 | 000 | 000000 | SKIPE |
| 226 | 00 | $0 \quad 00$ | 000000 | MULM | 337 | 00 | 000 | 000000 | SKIPG |
| 434 | 00 | 000 | 000000 | OR | 335 | 00 | 000 | 000000 | SKIPGE |
| 437 | 00 | 000 | 000000 | ORB | 331 | 00 | 000 | 000000 | SKIPL |
| 454 | 00 | 000 | 000000 | ORCA | 333 | 00 | 000 | 000000 | SKIPLE |
| 457 | 00 | 000 | 000000 | ORCAB | 336 | 00 | 000 | 000000 | SKIPN |
| 455 | 00 | 000 | 000000 | ORCAI | 360 | 00 | 000 | 000000 | SOJ |
| 456 | 00 | 000 | 000000 | ORCAM | 364 | 00 | 000 | 000000 | SOJA |
| 470 | 00 | 000 | 000000 | ORCB | 362 | 00 | 000 | 000000 | SOJE |
| 473 | 00 | 000 | 000000 | ORCBB | 367 | 00 | 000 | 000000 | SOJG |

(Continued on Next Page)

Table C-l (Cont.)
Machine Instruction Mnemonics

| 365 | 00 | 0 | 00 | 000000 | SOJGE | 667 | 00 | 0 | 00 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

## MACRO-DEFINED MNEMONICS

## C. 2 I/O INSTRUCTION AND DEVICE CODE MNEMONICS

Table C-2 shows MACRO's I/O instruction mnemonics and the code each assembles. Note that $I / O$ machine instructions are executable only in executive mode.

Table C-2
I/O Instruction Mnemonics

| 7 | 000 | 00 | 0 | 00 | 000000 | BLK I | 7 | 000 | 30 | 0 | 00 | 000000 | CONS 2 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | 000 | 10 | 0 | 00 | 000000 | BLKO | 7 | 000 | 04 | 0 | 00 | 000000 | DATAI |
| 7 | 000 | 24 | 0 | 00 | 000000 | CONI | 7 | 000 | 14 | 0 | 00 | 000000 | DATAO |
| 7 | 000 | 20 | 0 | 00 | 000000 | CONO | 7 | 000 | 04 | 0 | 00 | 000000 | RSW |
| 7 | 000 | 34 | 0 | 00 | 000000 | CONSO |  |  |  |  |  |  |  |

Table C-3 shows MACRO's I/O device code mnemonics. Each is assembled with the $I / O$ instruction mnemonic DATAI so that the value of the device code will be in its proper field. In the first table entry, for example, the assembled code is:

$$
702404000000000
$$

where the 7 and 04 are generated by the DATAI instruction, and the 024 by the $A D C$ device code mnemonic.

NOTE
MACRO leaves these device code mnemonics as undefined symbols during Pass l. At the end of Pass 1 , the mnemonics are found in MACRO's tables only if one or more I/O instructions have been found.

Therefore, if a device code mnemonic is not assembled in Pass l, or if no $1 / 0$ instruction mnemonics were found, MACRO will not have defined the device code mnemonic.

## MACRO-DEFINED MNEMONICS

Table C-3
I/O Device Code Mnemonics

| 7 | 024 | 04 | $0 \quad 00$ | 000000 | DATAI ADC, |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | 030 | 04 | 000 | 000000 | DATAI ADC2, |
| 7 | 000 | 04 | 000 | 000000 | DATAI APR, |
| 7 | 014 | 04 | 000 | 000000 | DATAI CCI, |
| 7 | 110 | 04 | 000 | 000000 | DATAI CDP, |
| 7 | 114 | 04 | 000 | 000000 | DATAI CDR, |
| 7 | 070 | 04 | 000 | 000000 | DATAI CLK, |
| 7 | 074 | 04 | 000 | 000000 | DATAI CLK2, |
| 7 | 000 | 04 | 000 | 000000 | DATAI CPA, |
| 7 | 150 | 04 | 000 | 000000 | DATAI CR, |
| 7 | 154 | 04 | 000 | 000000 | DATAI CR2, |
| 7 | 200 | 04 | 000 | 000000 | DATAI DC, |
| 7 | 204 | 04 | 000 | 000000 | DATAI DC2, |
| 7 | 300 | 04 | 000 | 000000 | DATAI DCSA, |
| 7 | 304 | 04 | 000 | 000000 | DATAI DCSB, |
| 7 | 270 | 04 | 000 | 000000 | DATAI DDC, |
| 7 | 274 | 04 | 000 | 000000 | DATAI DDC2, |
| 7 | 270 | 04 | 000 | 000000 | DATAI DF, |
| 7 | 130 | 04 | 000 | 000000 | DATAI DIS, |
| 7 | 134 | 04 | 000 | 000000 | DATAI DIS2, |
| 7 | 060 | 04 | 000 | 000000 | DATAI DLB, |
| 7 | 160 | 04 | 000 | 000000 | DATAI DLB2, |
| 7 | 064 | 04 | 000 | 000000 | DATAI DLC, |
| 7 | 164 | 04 | 000 | 000000 | DATAI DLC2, |
| 7 | 240 | 04 | 000 | 000000 | DATAI DLS, |
| 7 | 244 | 04 | 000 | 000000 | DATAI DLS2, |
| 7 | 250 | 04 | 000 | 000000 | DATAI DPC, |
| 7 | 254 | 04 | 000 | 000000 | DATAI DPC2, |
| 7 | 260 | 04 | 000 | 000000 | DATAI DPC3, |
| 7 | 264 | 04 | 000 | 000000 | DATAI DPC4, |
| 7 | 464 | 04 | 000 | 000000 | DATAI DSI, |
| 7 | 474 | 04 | 000 | 000000 | DATAI DSI2, |
| 7 | 170 | 04 | 000 | 000000 | DATAI DSK, |
| 7 | 174 | 04 | 000 | 000000 | DATAI DSK2, |
| 7 | 460 | 04 | 000 | 000000 | DATAI DSS, |
| 7 | 470 | 04 | 000 | 000000 | DATAI DSS2, |
| 7 | 320 | 04 | 000 | 000000 | DATAI DTC, |
| 7 | 330 | 04 | 000 | 000000 | DATAI DTC2, |
| 7 | 324 | 04 | 000 | 000000 | DATAI DTS, |
| 7 | 334 | 04 | 000 | 000000 | DATAI DTS2, |
| 7 | 124 | 04 | 000 | 000000 | DATAI LPT, |
| 7 | 234 | 04 | 000 | 000000 | DATAI LPT2, |
| 7 | 260 | 04 | 000 | 000000 | DATAI MDF, |
| 7 | 264 | 04 | 000 | 000000 | DATAI MDF2, |
| 7 | 220 | 04 | 000 | 000000 | DATAI MTC, |
| 7 | 230 | 04 | $0 \quad 00$ | 000000 | DATAI MTM, |
| 7 | 224 | 04 | $0 \quad 00$ | 000000 | DATAI MTS, |
| 7 | 010 | 04 | $0 \quad 00$ | 000000 | DATAI PAG, |

(Continued on Next Page)

## MACRO-DEFINED MNEMONICS

Trable C-3 (Cont.)
I/O Device Code Mnemonics

| 7 | 004 | 04 | 0 | 00 | 000000 |  | DATAI PI, |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 7 | 140 | 04 | 0 | 00 | 000000 |  | DATAI PLT, |
| 7 | 144 | 04 | 0 | 00 | 000000 |  | DATAI PLT2, |
| 7 | 100 | 04 | 0 | 00 | 000000 |  | DATAI PTP, |
| 7 | 104 | 04 | 0 | 00 | 000000 | DATAI PTR, |  |
| 7 | 340 | 04 | 0 | 00 | 000000 | DATAI TMC, |  |
| 7 | 350 | 04 | 0 | 00 | 000000 | DATAI TMC2, |  |
| 7 | 344 | 04 | 0 | 00 | 000000 | DATAI TMS, |  |
| 7 | 354 | 04 | 0 | 00 | 000000 | DATAI TMS2, |  |
| 7 | 120 | 04 | 0 | 00 | 000000 | DATAI TTY, |  |
| 7 | 210 | 04 | 0 | 00 | 000000 | DATAI UTC, |  |
| 7 | 214 | 04 | 0 | 00 | 000000 | DATAI UTS, |  |

## MACRO-DEFINED MNEMONICS

## C. 3 KLIO EXTEND INSTRUCTION MNEMONICS

Table $C-4$ shows the KLlo EXTEND instruction mnemonics and the code assembled by each. All of these mnemonics are defined only if MACRO is assembled with the KLlO switch on.

See the Supplement to the Hardware Reference Manual for a discussion of these EXTEND instructions.

Table C-4
KLlO EXTEND Instruction Mnemonics

| 002 | 00 | 0 | 00 | 000000 | *CMPSE | 010 | 00 | 0 | 00 | 000000 | *CVTDBO |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 007 | 00 | 0 | 00 | 000000 | *CMPSG | 011 | 00 | 0 | 00 | 000000 | *CVTDBT |
| 005 | 00 | 0 | 00 | 000000 | *CMPSGE | 004 | 00 | 0 | 00 | 000000 | *EDIT |
| 001 | 00 | 0 | 00 | 000000 | * CMPSL | 016 | 00 | 0 | 00 | 000000 | *MOVSLJ |
| 003 | 00 | 0 | 00 | 000000 | *CMPSLE | 014 | 00 | 0 | 00 | 000000 | *MOVSO |
| 006 | 00 | 0 | 00 | 000000 | *CMPSN | 017 | 00 | 0 | 00 | 000000 | *MOVSRJ |
| 012 | 00 | 0 | 00 | 000000 | *CVTBDO | 015 | 00 | 0 | 00 | 000000 | *MOVST |
| 013 | 00 | 0 | 00 | 000000 | *CVTBDT | 020 | 00 | 0 | 00 | 000000 | * XBLT |

## MACRO-DEFINED MNEMONICS

## C. 4 JRST AND JFCL MNEMONICS

Table C-5 shows mnemonics that assemble both operator and accumulator fields in the machine instruction. The left side of the table shows the mnemonics and the code they generate; the right side shows JRST and JFCL mnemonics with accumulators generating the equivalent code.

Table C-5
JRST and JFCL Mnemonics


## MACRO-DEFINED MNEMONICS

## C. 5 DECsystem-10 MONITOR CALL MNEMONICS

Tables $\mathrm{C}-6$ through $\mathrm{C}-9$ show MACRO's DECsystem-10 monitor call mnemonics and related mnemonics.

Table C-6 shows DECsystem-10 monitor calls and their assembled code.
Tables C-7 through C-9 show DECsystem-10 CALLI, TTCALL, and MTAPE mnemonics and their assembled code. These mnemonics are defined only if MACRO's UUOSYM switch is on when MACRO is assembled.

Table C-6
DECsystem-l0 Monitor Calls

| 040 | 00 | 0 | 00 | 000000 | CALL | 067 | 00 | 0 | 00 | 000000 | OUTPUT |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 047 | 00 | 0 | 00 | 000000 | CALLI | 052 | 00 | 0 | 00 | 000000 | * RDCLK |
| 070 | 00 | 0 | 00 | 000000 | CLOSE | 071 | 00 | 0 | 00 | 000000 | RELEAS |
| 077 | 00 | 0 | 00 | 000000 | ENTER | 055 | 00 | 0 | 00 | 000000 | RENAME |
| 062 | 00 | 0 | 00 | 000000 | GETSTS | 060 | 00 | 0 | 00 | 000000 | SETSTS |
| 056 | 00 | 0 | 00 | 000000 | IN | 061 | 00 | 0 | 00 | 000000 | STATO |
| 064 | 00 | 0 | 00 | 000000 | INBUF | 062 | 00 | 0 | 00 | 000000 | STATUS |
| 041 | 00 | 0 | 00 | 000000 | INIT | 063 | 00 | 0 | 00 | 000000 | STATZ |
| 066 | 00 | 0 | 00 | 000000 | INPUT | 051 | 00 | 0 | 00 | 000000 | TTCALL |
| 076 | 00 | 0 | 00 | 000000 | LOOKUP | 073 | 00 | 0 | 00 | 000000 | UGETF |
| 072 | 00 | 0 | 00 | 000000 | MTAPE | 100 | 00 | 0 | 00 | 000000 | UJEN |
| 050 | 00 | 0 | 00 | 000000 | OPEN | 074 | 00 | 0 | 00 | 000000 | USETI |
| 057 | 00 | 0 | 00 | 000000 | OUT | 075 | 00 | 0 | 00 | 000000 | USETO |
| 065 | 00 | 0 | 00 | 000000 | OUTBUF |  |  |  |  |  |  |

Table C-7 $\quad$ Mnemonics

| ALLOC. | 047 | 000 | 000 | 000144 | IPCFQ. |
| :---: | :---: | :---: | :---: | :---: | :---: |
| APRENB | 047 | 000 | 000 | 000142 | IPCFR. |
| ATTACH | 047 | 000 | 000 | 000143 | IPCFS. |
| CALll. | 047 | 000 | 000 | 000113 | JBSET. |
| CAL78. | 047 | 000 | 000 | 000103 | JOBPEK |
| CHGPPN | 047 | 000 | 000 | 000047 | JOBSTR |
| CHKACC | 047 | 000 | 000 | 000061 | JOBSTS |
| CLRST. | 047 | 000 | 000 | 777777 | LIGHTS |
| CNECT. | 047 | 000 | 000 | 000062 | LOCATE |
| COMPT. | 047 | 000 | 000 | 000060 | LOCK |
| CORE | 047 | 000 | 000 | 000015 | LOGIN |
| CTLJJOB | 047 | 000 | 000 | 000017 | LOGOUT |
| DAEFIN | 047 | 000 | 000 | 000111 | METER. |
| DAEMON | 047 | 000 | $0 \quad 00$ | 000023 | MSTIME |
| DATE | 047 | 000 | 000 | 000126 | MTAID. |
| DDTGT | 047 | 000 | $0 \quad 00$ | 000112 | MTCHR. |
| DDTIN | 047 | 000 | $0 \quad 00$ | 000131 | MVHDR. |
| DDTOUT | 047 | 000 | 000 | 000157 | NODE. |
| DDTRL | 047 | 000 | $0 \quad 00$ | 000077 | OTHUSR |
| DEBRK. | 047 | 000 | 000 | 000145 | PAGE. |
| DEQ. | 047 | 000 | 000 | 000110 | PATH. |
| DEVCHR | 047 | 000 | $0 \quad 00$ | 000033 | PEEK |
| DEVGEN | 047 | 000 | 000 | 000162 | PERF. |
| DEVLNM | 047 | 000 | $0 \quad 00$ | 000135 | PIINI. |
| DEVNAM | 047 | 000 | 000 | 000141 | PIRST. |
| DEVPPN | 047 | 000 | 000 | 000140 | PISAV. |
| DEVSIZ | 047 | 000 | $0 \quad 00$ | 000136 | PISYS. |
| DEVSTS | 047 | 000 | $0 \quad 00$ | 000030 | PJOB |
| DEVTYP | 047 | 000 | $0 \quad 00$ | 000114 | POKE. |
| DISK. | 047 | 000 | $0 \quad 00$ | 000021 | REASSI |
| DSKCHR | 047 | 000 | 000 | 000037 | REMAP |
| DVRST. | 047 | 000 | $0 \quad 00$ | 000117 | RESDV. |
| DVURS. | 047 | 000 | $0 \quad 00$ | 000000 | RESET |
| ENQ. | 047 | 000 | $0 \quad 00$ | 000057 | RTTRP |
| ENQC. | 047 | 000 | 000 | 000035 | RUN |
| ERLST. | 047 | 000 | $0 \quad 00$ | 000027 | RUNTIM |
| ERRPT. | 047 | 000 | 000 | 000150 | SCHED. |
| EXIT | 047 | 000 | 000 | 000056 | SEEK |
| FILOP. | 047 | 000 | $0 \quad 00$ | 000133 | SENSE. |
| FRCUUO | 047 | 000 | 000 | 000002 | SETDDT |
| FRECHN | 047 | 000 | 000 | 000043 | SETNAM |
| GETCHR | 047 | 000 | $0 \quad 00$ | 000032 | SETPOV |
| GETLIN | 047 | 000 | 000 | 000075 | SETUUO |
| GETPPN | 047 | 000 | 000 | 000036 | SETUWP |
| GETSEG | 047 | 000 | 000 | 000031 | SLEEP |
| GETTAB | 047 | 000 | 000 | 000042 | SPY |
| GOBSTR | 047 | 000 | 000 | 000050 | STRUUO |
| HIBER | 047 | 000 | 000 | 000146 | SUSET. |
| HPQ | 047 | 000 | 000 | 000020 | SWITCH |
| IONDX. | 047 | 000 | 000 | 000051 | SYSPHY |

## MACRO-DEFINED MNEMONICS

Table C-7 (Cont.)
DECsystem-10 CALLI Mnemonics

| 047 | 00 | 0 | 00 | 000046 | SYSSTR | 047 | 00 | 0 | 00 | 000025 | TRPSET |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 047 | 00 | 0 | 00 | 000154 | TAPOP. | 047 | 00 | 0 | 00 | 000120 | UNLOK. |
| 047 | 00 | 0 | 00 | 000022 | TIMER | 047 | 00 | 0 | 00 | 000013 | UTPCLR |
| 047 | 00 | 0 | 00 | 000044 | TMPCOR | 047 | 00 | 0 | 00 | 000010 | WAIT |
| 047 | 00 | 0 | 00 | 000115 | TRMNO. | 047 | 00 | 0 | 00 | 000073 | WAKE |
| 047 | 00 | 0 | 00 | 000116 | TRMOP. | 047 | 00 | 0 | 00 | 000063 | WHERE |
| 047 | 00 | 0 | 00 | 000026 | TRPJEN | 047 | 00 | 0 | 00 | 000124 | XTTSK. |

Table C-8
DECsystem-10 TTCALL Mnemonics

| 051 | 11 | 0 | 00 | 000000 | CLRBFI | 051 |  | 0 | 00 | 000000 | IONEOU |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 051 | 12 | 0 | 00 | 000000 | CLRBFO | 051 | 01 | 0 | 00 | 000000 | OUTCHR |
| 051 | 06 | 0 | 00 | 000000 | GETLCH | 051 | 03 | 0 | 00 | 000000 | OUTSTR |
| 051 | 02 | 0 | 00 | 000000 | INCHRS | 051 | 10 | 0 | 00 | 000000 | RESCAN |
| 051 | 00 | 0 | 00 | 000000 | INCHRW | 051 | 07 | 0 | 00 | 000000 | SETLCH |
| 051 | 05 | 0 | 00 | 000000 | INCHSL | 051 | 13 | 0 | 00 | 000000 | SKPINC |
| 051 | 04 | 0 | 00 | 000000 | INCHWL | 051 | 14 | 0 | 00 | 000000 | SKPINL |

Table C-9
DECsystem-10 MTAPE Mnemonics

| 072 | 00 | 0 | 00 | 000013 | MTBLK. | 072 | 00 | 0 | 00 | 000101 | MTIND. |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 072 | 00 | 0 | 00 | 000017 | MTBSF. | 072 | 00 | 0 | 00 | 000001 | MTREW. |
| 072 | 00 | 0 | 00 | 000007 | MTBSR. | 072 | 00 | 0 | 00 | 000016 | MTSKF. |
| 072 | 00 | 0 | 00 | 000100 | MTDEC. | 072 | 00 | 0 | 00 | 000006 | MTSKR. |
| 072 | 00 | 0 | 00 | 000003 | MTEOF. | 072 | 00 | 0 | 00 | 000011 | MTUNL. |
| 072 | 00 | 0 | 00 | 000010 | MTEOT. | 072 | 00 | 0 | 00 | 000000 | MTWAT. |

## MACRO-DEFINED MNEMONICS

## C. 6 F40 UUO MNEMONICS

Table C-lo shows mnemonics that are defined only if MACRO is assembled with the F 40 switch on. These mnemonics generate UUOS, which are handled properly if the program is running under control of the FORSE object-time system.

Table C-10
F40 UUO Mnemonics

| 020 | 00 | 0 | 00 | 000000 | ***DATA. | 017 | 00 | 0 | 00 | 000000 | ***OUT. |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 021 | 00 | 0 | 00 | 000000 | ***FIN. | 027 | 00 | 0 | 00 | 000000 | ***OUTF. |
| 016 | 00 | 0 | 00 | 000000 | ***IN. | 015 | 00 | 0 | 00 | 000000 | ***RESET. |
| 026 | 00 | 0 | 00 | 000000 | ***INF. | 022 | 00 | 0 | 00 | 000000 | ***RTB. |
| 024 | 00 | 0 | 00 | 000000 | ***MTOP. | 025 | 00 | 0 | 00 | 000000 | ***SLIST. |
| 031 | 00 | 0 | 00 | 000000 | ***NLI. | 023 | 00 | 0 | 00 | 000000 | ***WTB. |
| 032 | 00 | 0 | 00 | 000000 | ***NLO. |  |  |  |  |  |  |

## APPENDIX D

## PROGRAM EXAMPLES

The following pages contain examples of MACRO programs. Each program has been assembled with the /C (CREF) switch on; this produces a .CRF file for the program listing (instead of the usual. LST file). The /O switch has been used with the CREF program to produce a . LST file that includes all operators in an operator symbol table.



|  | Second EXAMFL | Example of MACRO Frosram MACRO MAC 2-Mar-78 15:03 | ```%53(1017) 16:12 2-Mar-78 Fase 3 E<amfle Two``` |
| :---: | :---: | :---: | :---: |
|  | 62 |  | ; . |
|  | 63 |  |  |
|  | 64 |  | LEFFINE CLEAR < |
|  | 65 |  | LIEFINE CONCAT (FTXT) < |
|  | 66 |  | DEFINE CONCAT (TEXT) < |
|  | 67 |  | CON1 TEXT>, FTXT\% |
|  | 68 |  |  |
|  | 69 |  | IEFINE EXFANI \&FTXT> |
|  | 70 |  |  |
|  | 71 |  | IIEFINE EXPANI |
|  | 72 |  | $>$ |
|  | 73 |  |  |
|  | 74 |  |  |
|  | 75 |  | DEFINE CONCAT (TEXT) < |
|  | 76 |  | CON1 STEXT>, OTXT'NTXT> $^{\text {a }}$ |
|  | 77 |  |  |
|  | 78 |  |  |
|  | 79 |  | $\rangle$ |
|  | 80 |  |  |
|  | 81 |  | SALL |
|  | 82 |  |  |
|  | 83 |  | ClLEAF |
|  | 84 |  |  |
|  | 85 |  | CONCAT $10 \%$ |
|  | 86 |  | CONCAT <> |
|  | 87 |  | CONCAT <"A"> |
|  | 88 |  | CONCAT <, <-1, 6\%8177\% |
|  | 89 |  |  |
|  | 90 |  | LALL |
|  | $\begin{aligned} & 91 \\ & 92 \end{aligned}$ | $000000^{\prime} 01010100600000$ |  |
|  | 93 |  | SALL |
|  | 94 |  | CLEAF |
|  | 95 |  |  |
|  | 96 |  | CONCAT SIEF\% |
|  | 97 |  | CONCAT SINE FOO (\% |
|  | 98 |  | CONCAT SN) |
|  | 99 100 |  | $\begin{aligned} & \text { CONCAT < } \\ & \text { IEFINE; } \end{aligned}$ |
|  | 101 |  | CONCAT \& BAR ( N ) 3*N\% |
|  | 102 |  | $\rangle$ |
|  | 103 |  |  |
|  | 104 |  | ; . . - |

EXAMFL MAC 2-Mar-78 15:03 E×amFle Two

| 105 |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 106 |  |  |  |  |
| 107 |  |  |  | LALL |
| 108 |  |  |  | EXFANIMMEF |
| 109 |  |  |  | IEFINE BAR |
| 110 |  |  |  | - |
| 111 |  |  |  |  |
| 112 | 000001 | 000000 | 000004 | FOO 2~2*2* |
| 113 | 000002' | 000000 | 000006 | FOO 3~2*3~ |
| 114 | 000003' | 000000 | 000006 | BAR 2^3*2^ |
| 115 | 000004' | 000000 | 000011 | BAF $3^{\sim} 3^{*} 3^{\text {a }}$ |
| 116 |  |  |  |  |
| 117 |  |  |  | FRGENII |

NO ERFORS METECTED
PROGFAM BREAK IS 000005 CFU TIME USEII OO:00.180 1 OF CORE USEI


```
Third Examfle of MACFO Frosram MACFO %S3(1017) 16:12 2-Mar-78 Fase 6
    ; . . .
    SEARICH MACROS
    . IIFECTIUE SFCONI
    HEFINE NUMLST %
        X (FIFST,1)
        X (FOUFTH,4)
        x (SECONI,2)
        X (THIFII,3)
    >
    LIEFINE X (TEXT,JUNN゙) EEXF SIXEIT /TEXT/&
    NAMTBL: NUMLST
            X (FJFST,1)"EXF SIXEIT /FIFST/m
            X (FOUFTTH,4)MEXF SIXEIT /FOURTH/~
            X (SECONH,2)MEXF SIXEIT /SECONI/~
            x (THIRI, 3)`EXF SIXRIT THIRL/m
BLIEEN:==, NAMTEL
    HEFINE X (JUNK,LABL) \$'LABL`
    BLTBL: NUMLST'
    X (FIFST,1)"$1"
    x (FIFST,1)"$1"
    x (FOURTH,4)"$4
    x (SECONII,2)- $2m
    x (THIRI,3)"$3-
    IEFINE X (JUNK,VALU) SIEC VALU.
    VALTEL: NUMLST
        X (FIFST,1) MEC 1'
        X (FOUFTH,4)~ LEEC 4-
        x (FOURTH,4) INEC 4
        X (THIRD,3) MEC 3
. . .
```



NO ERFORS DETECTEN
FROGRAM BREAK IS 000020 CFU TIME USEI OO:00,142

IOF CORE USEA
EXAMFL MAC
MACRO \%53(1017) 16:12 2-Mar-78 15:03
EXIT
EXIT
EBLTBL
$047000 \quad 000012$
$000004^{\prime}$
THLLEN 000000
$\begin{array}{llll}\text { TOFSSO } & 777777 & 777777 & 5 \mathrm{FPD}\end{array}$
UALTEL
$\$ 1$
$\$ 2$
$\$ 2$
$\$ 3$
$\$ 3$
$\$ 4$
000010 $000014^{\prime}$ 000015 000016 000017 ,


E\%amfle Four
; . . .
DEFINE COMMON (COM, VAFS, \%FAL, \%LEN, \%VAL, \%COM, \%FAS) (
. XCFEF \%FAL, \%LEN, \%UAL, \%COM,\%FAS
; Temf macro to strif one pair of ansle brackets from ; a macro arsument and fass it to another macro

LIEFINE \%FAS (A,B) A B
; ; Temp macio to comfute lensth of COMMON
DEFINE \%COM (UAF,LEN 1 ) ( $\%$ LEN $==\% L E N+L E N$
$\begin{array}{ll}\text { \%FAI }==10 & \text {; Save currert radiw, use } 10 \\ \text { FAIIX } 10 & \text {; s so defs read like FORTKAN }\end{array}$
FAIIIX $10 \quad ;$; so def's read like FORTFAN
\#LEN $==0 \quad$; Set to courit leristh of COMMON
IFF UARS \%FAS \%COM,VARS.

- COMMON COM[\%LEN]
; ; Get lensth of this COMMON
; Allocate the whole COMmON
IIEFINE \%COM (UAR,LEN $1 \geqslant$ ) ; ;Set us another temf macro $U A F=\% V A L \quad$; Ilefine COMMON block eritry \%UAL $==\%$ UAL $+L E N \quad$;innement to ne⿻t enitry
$>$
\%LEN $==0$
\%UEN $==0$
\%UAL $==\mathrm{COM}$
IRF UARS $\%$ F'AS \%COM, UARS:
FAIIX \%RAII
IF2, FURGE \%LEN, \%FAAI, \%UAL, \%COM, \%FAS\% ; ; Kees sumbol table clean
$y$

```
Fourth ExamFle of MACFO Frosram MACFO %53(1017) 16:12 2-Mar-78 Fase 10
EXAMFL MAC 2-Mar-78 15:03 E<amFle Four
```

290
291
292
293
294
295
297
298
299
300
300 302000001 , 202 $01000000000 *$ 000001 000002 , 26317000000000

305
306
$307000003^{\prime} 20001000000007$
$308000004 \quad 40200000000000$ $309000005 \quad 25101000000000$
$310-000006^{\prime 26317000000000}$
311
$312000007 \quad 1$
313 000007' 000000* 000000*
314
315
NO ERFRORS IIETECTED
mole Four
; •••
INTEGER SNGLE,ARRAY,MULTI
REAL REAL
IIOUBLE FRECISION IIOUBLE
COMMON /AFEA/SNGLE,REAL, DOUBLE,AFFAY(10),MULTI(5,10)

SSamele routine to do SNGLE=IFIX(FEAL)
IFIX: FIX 1, FEAL
MOUEM 1,SNGLE
FOFJ 17,
;Sample routine to set all elements in ARFAY to 0
ZERD: MOVE $1,[X W D$ ARFAY,ARFAYYI]
SETZM ARFAY

FOFJ 17,
LIT

```
ENI
```

FROGRAM BREAK IS 000010
CFU TIME USED 00:00.228
1OF CORE USED

| Fourth EXAMFL | Example of MACRO MAC 2-Mar- | $\begin{aligned} & \text { Frosr } \\ & 78 \quad 15: \end{aligned}$ |
| :---: | :---: | :---: |
| AREA | 000001' | ext |
| ARRAY | 000000000000\# | fol |
| IOUBLE | 000000000000* | fol |
| IFIX | 000000' |  |
| MULTI | 000000000000 * | fol |
| FEAL | 000000000000* | fol |
| SNGLE | 000000* |  |
| ZEF'O | 000003 |  |


| AREA | 298* | 298307 | 308 | 309 |
| :---: | :---: | :---: | :---: | :---: |
| ARFAY | 298* |  |  |  |
| nOUBLE | 298* |  |  |  |
| IFIX | 301* |  |  |  |
| LELTEL | 181* |  | 204 | 206 |
| MULTI | 298* |  |  |  |
| NAMTEL | 1724 | 177 |  |  |
| REAL | 298* | 301 |  |  |
| SNGLE | 298* | 302 |  |  |
| TELLEN | 177* |  |  |  |
| TOFS10 | 200 | 202 |  |  |
| UALTEL | 189* |  |  |  |
| ZEFO | 307* |  |  |  |
| \$1 | 182 | 199\# |  |  |
| \$2 | 184 | 201* |  |  |
| \$3 | 185 | 203* |  |  |
| \$4 | 183 | 205* |  |  |

■ $I$ -


|  | ELLT | 309 |  |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | BYTE IEC | 91 190 | 191 | 192 | 193 |  |  |  |  |  |  |  |  |  |  |
|  | DEF INE | 10 | 64 | 74 | 83 | 85 | 86 | 87 | 88 | 94 | 96 | 97 | 98 | 100 | 102 |
|  |  | 108 | 109 | 163 | 170 | 179 | 187 | 258 | 298 |  |  |  |  |  |  |
|  | ENI | 315 |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | EXF | 173 | 174 | 175 | 176 |  |  |  |  |  |  |  |  |  |  |
|  | FIX | 301 |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | IF2 | 298 |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | IFE | 200 | 202 | 204 | 206 |  |  |  |  |  |  |  |  |  |  |
|  | IFN | 200 | 202 | 204 | 206 |  |  |  |  |  |  |  |  |  |  |
|  | IFNDEF | 200 | 202 | 204 | 206 |  |  |  |  |  |  |  |  |  |  |
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|  | L.ALL | 90 | 107 |  |  |  |  |  |  |  |  |  |  |  |  |
|  | LIT | 312 |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | move | 307 |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | MOUEM | 302 |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | FOF'J | 303 | 310 |  |  |  |  |  |  |  |  |  |  |  |  |
|  | FRGENI | 21 | 117 | 208 |  |  |  |  |  |  |  |  |  |  |  |
|  | FURGE | 298 |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | FALIIX | 298 |  |  |  |  |  |  |  |  |  |  |  |  |  |
|  | SALL | 81 | 93 |  |  |  |  |  |  |  |  |  |  |  |  |
|  | SEARCH | 160 |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\square$ | SETZM | 308 |  |  |  |  |  |  |  |  |  |  |  |  |  |
| $\stackrel{1}{\square}$ | SIXBIT | 173 | 174 | 175 | 176 |  |  |  |  |  |  |  |  |  |  |
| $\sigma$ | SUBTTL | 1 | 22 | 118 | 209 |  |  |  |  |  |  |  |  |  |  |
|  | TITLE | 23 | 119 | 210 |  |  |  |  |  |  |  |  |  |  |  |
|  | UNIVER | 2 |  |  |  |  |  |  |  |  |  |  |  |  |  |
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## APPENDIX E <br> PSEUDO-OPS FOR PROGRAM COMPATIBILITY

The pseudo-ops in this appendix are included only for compatibility of old programs to be assembled using MACRO Version 53.
HISEG

| FORMAT | HISEG address |
| :--- | :--- |
| address = program high-segment origin address. Must be |  |
| equal to or greater than 400000 and must be a |  |
| multiple of loo . |  |

## RIM

FORMAT
RIM

FUNCTION Specifies a format for absolute binary programs (useful only for PDP-6 systems), and consists of a series of paired words.

The first word of each pair is a paper-tape read instruction giving the memory address of the second word. The last pair of words is a transfer block; the first is an instruction obtained from the END statement and executed when the transfer block is read, and the second is a dummy word to stop the reader.


FUNCTION Causes a program format in which programs are absolute, unblocked, and not checksummed. When the RIMl0 statement follows a LOC statement in a program, the assembler punches out each storage word in the object program, starting at the absolute address specified in the LOC statement. RIM10 writes an arbitary "paper tape"; if it is in the format given below, it can be read by the DECsystem-10 Read-In Mode hardware.

IOWD n,first
where $n$ is the length of the program including the ending word transfer, and first is the first memory location to be occupied. The last location must contain a transfer instruction to begin the program, such as

JFST 4,G0

For example, if a program with RIMlO output has its first location at START and its last location at FINISH, you can write

IOWIFTNISH-START+1, START

NOTE
If the location counter is increased but no binary output occurs (for example, BLOCK, LOC, and VAR pseudo-ops), MACRO inserts a zero word into the binary output file for each location skipped by the location counter.

## RIM10B

RIM10B

If a program is assembled into absolute locations (not relocatable), a RIM1OB statement following the LOC statement at the beginning of the source program causes the assembler to write out the object program in RIMIOB format. This format is designed for use with the DECsystem-10 Read-In Mode hardware.

The program is punched during Pass 2, starting at the location specified in the LOC statement. If the first two statements in the program are

LOC 1000
RIM10E
MACRO assembles the program with absolute addresses starting at 1000 and punches the program in RIM10B format, also starting at location 1000. You can reset the location counter during assembly, but only one RIM10B statement is needed to punch the entire program.

In RIMIOB format, the assembler punches the RIMIOB Loader, followed by the program in l7-word (or less) data blocks, each block separated by blank tape. The assembler inserts an I/O transfer word (IOWD) preceding each data block, and also inserts a 36-bit checksum following each data block. The word count in the IOWD counts only the data words in the block, and the checksum is the 36 -bit added checksum of the IOWD and the data words.

Data blocks can contain less than 17 words. If the assembler assigns a nonconsecutive location, the current data block is terminated, and an IOWD containing the next location is inserted, starting a new data block.

The transfer block consists of two words. The first word of the transfer block is an instruction obtained from the END statement. This first word is executed when the transfer block is read. The second word is a dummy word to stop the reader.

## APPENDIX F

STORAGE ALLOCATION

MACRO allocates storage in two directions:

1. User symbols and macronames are entered in the symbol tables.
2. Macros and literals are entered in free space.

A symbol table entry is two words long. The first word is the symbol name in SIXBIT. The second word has flags in the left half, and either the value or a pointer in the right half. The flags indicate symbol type and attributes.

The following list shows how symbols and values are stored.

Type
18-bit symbol
36-bit symbol (includes OPDEFs and negative numbers)

EXTERNAL symbol

Polish symbol

How Stored
Value in right half of second word.
Value in free storage with a pointer in symbol table.

Pointer in symbol table to a 2-word block in free storage. The first word is the value that is the last reference in a chain of references to the symbol; the second word is the symbol name in SIXBIT.

The symbol table entry points to a 2-word block:
word 1: 0
word 2: negative number, address
Word 1 is the relocation word and is always zero. Word 2 gives the address of a Polish stack in free storage. The Polish stack is of the form:

## STORAGE ALLOCATION

Synonym operator (SYN argument)

Macroname

Words 3 and 4 designate an operand. If the operator is binary, words 5 and 6 designate the second operand; if the operator is unary, the stack contains only four words.

If an operand is EXTERNAL, its two words (3 and 4 , or 5 and 6) are:
word i: pointer to EXTERNAL symbol word $i+1$ : 0

If an operand is itself a Polish symbol, its two words are:

```
    word i: Polish pointer
```

    word i+l: 0
    Polish stack containing:
word 1: 0
word 2: 15
word 3: -2
word 4: referenced PSECT index word 5: relocation constant word 6: address

SIXBIT operator name in free storage with a pointer in the symbol table.

Value in free storage with a pointer to the text string in symbol table.

The text string is stored in a 4-word block of the form:
word 1: link to next block (0 if
last), ,two characters
word 2: five characters
word 3: five characters
word 4: five characters
However, the first such block is special:
word l: link to next block,,link to last block
word 2: pointer to default arg., ,number of args expected + reference count
word 3: five characters
word 4: five characters
The number of args expected is the number of dummy-arguments in the macro definition.

The reference count is incremented when the macro is called and decremented when the macro is exited. When this count goes to zero, the macro is removed from free space.

## STORAGE ALLOCATION

| Macro arguments | Stored in the same linked block, but not in the symbol table. Repeats (two or more times) are also stored in the same way. The text blocks are removed when the macro exits or the repeat exits, since the reference count has gone to zero. <br> The addresses of the actual argument blocks are stored in a pushdown stack in order of generation. <br> Default arguments are stored in the same way, except that the list is in free core. The pointer to the default arg list is stored in the left half of the second word of the first block of the macro definition. |
| :---: | :---: |
| Macros | The macrobody is stored as is, except that dummy-arguments are replaced by special symbols. |
|  | ASCII 177 (RUBOUT) signals that the next character is a special character, as follows: |
|  | 001 ;end of macro <br> 002 ;end of dummy symbol |
|  | 003 ; end of REPEAT |
|  | 004 ;end of IRP or IRPC |
|  | 005 ; RUBOUT |
|  | If the character is more than 5 and less than l00, it is illegal. |
|  | If the character is greater than or equal to l00, it is a dummy symbol; the value of the character is ANDed with 37 to get the dummy symbol number, and the corresponding pointer retrieved from the stack of actual arguments. |
|  | If the symbol was not specified (that is, has no pointer), and if the 40 bit is on, this symbol requires a created symbol, and one is created; otherwise the argument is ignored. |

NOTE

```
Verbose macros can use too much
storage space.
```

```
Four-word block for each word generated
```

    word 1: form word
    word 2: relocation bits
    word 3: code
    word 4: pointer to next block
    Form word is the word used for listing.
    This word is not checked when comparing
    literals, so that different forms
    producing the same code are classed as
    equal.
    Relocation bits are 0,1 , or EXTERNAL
    pointers.
    Pointer is the address of the zero word
    of the next block.
    NOTE
Long literals slow assembly and use storage; they should be written as subroutines or inline code.

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